# PDP-11 MACRO-11

# Language Reference Manual

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		CONTRACTOR		
		CONTENTS		
		· · · · · · · · · · · · · · · · · · ·		
		가지 않는 동네가 가지 않는 것을 하는 것을 하는 것을 하는 것을 가지 않는 것을 하는 것을 하는 것을 하는 것을 가지 않는 것을 하는 것을 가지 않는 것을 가지 않는 것을 가지 않는 것을 가지 않는 이 것은 것은 것은 것은 것은 것은 것을 하는 것을 하는 것을 하는 것을 수 있는 것을 수 있는 것을 하는 것을 수 있는 것을 수 있는 것을 하는 것을 수 있는 것을 수 있는 것을 하는 것을 수 있는		
			19 - C	Page
			• (1) \$ 100 miles	
PREFACE				ix
	0.1	MANUAL OBJECTIVES AND READER ASSUMPTIONS		ix
	0.2	STRUCTURE OF THE DOCUMENT	おまして	ix
	0.3	ASSOCIATED DOCUMENTS		x
	0.3	DOCUMENT CONVENTIONS		x
	0.4	DOCOMENT CONVENTIONS		
		DADE T THERODUCETON EO MACRO 11		
		PART I INTRODUCTION TO MACRO-II		
	-			
CHAPTER	T	MACRO-11 FEATURES		T-T
	1.1	OVERVIEW OF MACRO-11		1-1
	1.1.1	Assembly Pass 1		1-1
	1.1.2	Assembly Pass 2		1-2
		· · · · · · · · · · · · · · · · · · ·		
CHAPTER	2	SOURCE PROGRAM FORMAT		2-1
	2.1	PROGRAMMING STANDARDS AND CONVENTIONS		2-1
	2.2	STATEMENT FORMAT		2-1
	2 2 1	Label Field		$\frac{1}{2}$
	2.2.1	Operator Field		2-4
	2.2.2	Operand Riold		2-4
	2.2.3	Operand Field	S. 8 . 6	2-4
	2.2.4			2-5
	2.3	FORMAT CONTROL		2-0
		TT PROGRAMMENT TH MAGRA 11 ACCOMPTN TANG		
	PART	11 PROGRAMMING IN MACRO-11 ASSEMBLY LANG	UAGE	
	یر رود			~ 1
CHAPTER	3	SYMBOLS AND EXPRESSIONS		3-1
	3.1	CHARACTER SET		3-1
	3.1.1	Separating and Delimiting Characters		3-2
	3.1.2	Illegal Characters		3-3
	3.1.3	Unary and Binary Operators		3-4
	3.2	MACRO-11 SYMBOLS		3-5
	3.2.1	Permanent Symbols		3-5
	3.2.2	User-Defined and Macro Symbols		3-5
	3.3	DIRECT ASSIGNMENT STATEMENTS		3-7
	3 1	DECISTED SYMBOLS		3-9
	2 5	LOCAL SYMBOLS		3-10
	3.5	CURRENT LOCATION COUNTER		3_11
	3.0	CURRENT LOCATION COUNTER		2-12
	3.7	NUMBERS		2 14
	3.8	TERMO		2.1
	3.9	EXPRESSIONS		2-12
CHAPTER	4	RELOCATION AND LINKING		4-⊥
	_			
CHAPTER	5	ADDRESSING MODES		5-1
		3. E. E. M. M. K. S. M. M. K. S. M. M. M. K. M. M. M. K. M.		<b>-</b> -
	5.1	REGISTER MODE		5-1
	5.2	REGISTER DEFERRED MODE		5-2
	5.3	AUTOINCREMENT MODE		5-2

.

-

			Page
	5.4	AUTOINCREMENT DEFERRED MODE	5-3
	5 5	AUTODECDEMENT MODE	5-5
	5.5	AUTODECREMENT MODE	5-3
	5.0	TNDEY MODE	5-3
	5.7	INDEX NODE	5-4
	5.0	INDEX DEFERRED MODE	5-4
	5.9	IMMEDIATE MODE	5-4
	5.10	ABSOLUTE MODE	5-5
	5.11	RELATIVE MODE	5-5
	5.12	RELATIVE DEFERRED MODE	5-6
	5.13	SUMMARY OF ADDRESSING FORMS	5-7
	5.14	BRANCH INSTRUCTION ADDRESSING	5-8
	5.15	USING TRAP INSTRUCTIONS	5-8
		PART III MACRO-II DIRECTIVES	
CHAPTER	6	GENERAL ASSEMBLER DIRECTIVES	6-1
	6 1	LISTING CONTROL DIRECTIVES	6_1
	6 1 1	LIST and NITST Directives	6 1
	6 1 2	Dage Weadings	6 0
	6 1 2	mile Dimontino	0-0
	6.1.3	COMMI Directive	0-11
	0.1.4	.SBTTL Directive	6-11
	6.1.5	.IDENT Directive	6-12
	6.1.6	.PAGE Directive/Page Ejection	6-13
	6.2	FUNCTION DIRECTIVES: .ENABL AND .DSABL	6-13
	6.3	DATA STORAGE DIRECTIVES	6-17
	6.3.1	.BYTE Directive	6-17
	6.3.2	.WORD Directive	6-18
	6.3.3	ASCII Conversion Characters	6-19
	6.3.4	.ASCII Directive	6-20
	6.3.5	.ASCIZ Directive	6-21
	6.3.6	.RAD50 Directive	6-22
	6.3.7	Temporary Radix-50 Control Operator: ^R	6-23
	6.4	RADIX AND NUMERIC CONTROL FACILITIES	6-24
	6.4.1	Radix Control and Unary Control Operators	6-24
	6.4.1.1	.RADIX Directive	6-24
	6.4.1.2	Temporary Radix Control Operators: ^D, ^O,	
		and ^B	6-25
	6.4.2	Numeric Directives and Unary Control Operators	6-26
	6.4.2.1	.FLT2 and .FLT4 - Floating-Point Storage	
		Directives	6-27
	6.4.2.2	Temporary Numeric Control Operators: ^C and ^F	6-27
	6.5	LOCATION COUNTER CONTROL DIRECTIVES	6-29
	6.5.1	.EVEN Directive	6-29
	6.5.2	.ODD Directive	6-29
	6.5.3	.BLKB and .BLKW Directives	6-30
	6.6	TERMINATING DIRECTIVES	6-31
	6.6.1	.END Directive	6-31
	6.6.2	.EOT Directive	6-31
	6.7	PROGRAM BOUNDARIES DIRECTIVE: .LIMIT	6-31
	6.8	PROGRAM SECTIONING DIRECTIVES	6-32
	6.8.1	.PSECT Directive	6-32
	6.8.1.1	Creating Program Sections	6-36
	6.8.1.2	Code or Data Sharing when every second and	6-38
	6.8.1.3	Memory Allocation Considerations	6-38
	6.8.2	.ASECT and .CSECT Directives	6-38
	6.9	SYMBOL CONTROL DIRECTIVE: .GLOBL	6-39
	6.10	CONDITIONAL ASSEMBLY DIRECTIVES	6-41

			Page
	6.10.1	Conditional Assembly Block Directives: .IF,	6-41
	6.10.2	Subconditional Assembly Block Directives:	0-41
	6.10.3	.IFF, .IFT, .IFTF Immediate Conditional Assembly Directive:	6-43
	C 10 4		6-46
	6.10.4	PAL-IIR Conditional Assembly Directives	6-46
CHAPTER	7	MACRO DIRECTIVES	7-1
	7.1	DEFINING MACROS	7-1
	7.1.1	.MACRO Directive	7-1
	7.1.2	.ENDM Directive	7-2
	7.1.3	.MEXIT Directive	7-3
	7.1.4	MACRO Definition Formatting	7-3
	7.2	CALLING MACROS	7-3
	7.3	ARGUMENTS IN MACRO DEFINITIONS AND MACRO CALLS	7-4
	7.3.1	Macro Nesting	7-5
	7.3.2	Special Characters in Macro Arguments	7-6
	7.3.3	Passing Numeric Arguments as Symbols	7-6
	7.3.4	Number of Arguments in Macro Calls	7-7
	7.3.5	Creating Local Symbols Automatically	7-7
	7 3 6	Keyword Arguments	7_9
	737	Concatenation of Macro Arguments	7-10
	7.4	MACRO ATTRIBUTE DIRECTIVES: .NARG, .NCHR, AND	7 11
	7 4 1	NITED DEVICE SELECTION SELECTION CONTRACTOR	7 11
	7.4.1	NARG Directive	7 10
	7 4 2	NUTR Directive	7 12
	7.4.3	.NTYPE Directive	7-13
	7.6	INDEFINITE REPEAT BLOCK DIRECTIVES: .IRP AND	/-14
		•IRPC has a state of the state	7-15
	7.6.1	.IRP Directive	7-15
	7.6.2	.IRPC Directive	7-16
	7.7	REPEAT BLOCK DIRECTIVE: .REPT, .ENDR	7-17
	7.8	MACRO LIBRARY DIRECTIVE: .MCALL	7-18
APPENDIX	A	MACRO-11 CHARACTER SETS	A-1
	ר ב	ASCIT CHARACTER SET	A-1
	A 2	PADIX-50 CHARACTER SET	A-1 A-4
	<b>A</b> • 2	RADIA-JU CHARACIER DEI	A-4
APPENDIX	B	MACRO-11 ASSEMBLY LANGUAGE AND ASSEMBLER DIRECTIVES	B-1
	B.1	SPECIAL CHARACTERS	B-1
	в.2	SUMMARY OF ADDRESS MODE SYNTAX	B-1
	B.3	ASSEMBLER DIRECTIVES	B-2
APPENDIX	0-707 C	PERMANENT SYMBOL TABLE (PST)	C-1
	c.1	OP CODES	c-1
	C.2	MACRO-11 DIRECTIVES	C-4
APPENDIX	D	DIAGNOSTIC ERROR MESSAGE SUMMARY	D-1
	D.1	MACRO-11 ERROR CODES	D-1

v

1

				Page
APPENDIX	E	SAMPLE CODING STANDARD		E-1
	E.1	INTRODUCTION		E-1 E-1
	E.3 7	COMMENTS NAMING STANDARDS		E-2 E-2
	E.4.1	General Purpose Registers		E-2 E-2
	E.4.1.2 E.4.1.3	Hardware Registers		E-2 E-2
	E.4.2	Processor Priority		E-3
	E.4.3 F A A	Uther Symbols Using the Standard Symbolics		E-3
	E.4.5	Symbols		E-3
	E.4.5.1	Global Symbols		E-3
	E.4.5.2	Symbol Examples		E-3 F-4
	E.4.5.3	Program-Local Symbols		E-4
	E.4.5.4	Macro Names		E-5
	E.5	PROGRAM MODULES		E-5
	E.5.1	General Comments on Programs		E-5
	E.5.2	The Module Preface		E-5
	E.5.3	Formatting the Module Preface		E-7
	E.5.4	Modularity	***	E-8
	E.5.4.1	Calling Conventions (Inter-Module)		E-8
	E.5.4.2	S Exiting Statistics Substitution Condition		E-9
	E.5.4.3	Intra-Module Calling Conventions		E-9
	E.5.4.4	Success/Failure Indication		E-9
	E.5.4.5	Module Checking Routines		E-9
	E.6	FORMATTING STANDARDS		E-9
s i se l	E.6.1	Program Flow and the first states of the states.		E-9
	E.6.2	Common Exits		E-11
	E.6.3	Code with Interrupts Inhibited		E-12
	E.7	PROGRAM SOURCE FILES		E-12
	E.8	FORBIDDEN INSTRUCTION USAGE	C. A.	E-12
	E.9	RECOMMENDED CODING PRACTICE		E-13
	E.9.1	Conditional Branches		E-13
	E.10	PDP-11 VERSION NUMBER STANDARD		E-13
	E.10.1 E.10.2	Displaying the Version Identifier Use of the Version Number in the Program	in de la cati Na catilitada	E-14 E-15
APPENDIX	F	ALLOCATING VIRTUAL MEMORY		F-1
	F.1 53	GENERAL HINTS AND SPACE-SAVING GUIDELINES		F-l
	F.2	MACRO DEFINITIONS AND EXPANSIONS		F-2
	F.3	OPERATIONAL TECHNIQUES		F-4
1. ** E	_			-
APPENDIX	G	WRITING POSITION INDEPENDENT CODE		G-1
	G.1	INTRODUCTION TO POSITION INDEPENDENT CODE		G-1
	G.2	EXAMPLES A FOLLAR A TO DEPART REMERSION		G-2
APPENDIX	Н	SAMPLE ASSEMBLY AND CROSS REFERENCE LISTIN	G	H-1

all instruction

.9

## Page

#### FIGURES

FIGURE	3-1	Assembly Listing Showing Local Symbol Block	3-11
	3-2	Sample Assembly Results	3-12
	6-1	Example of Line Printer Assembly Listing	6-6
	6-2	Example of Terminal Assembly Listing	6-7
	6-3	Listing Produced With Listing Control	
		Directives	6-9
	6-4	Assembly Listing Table of Contents	6-12
	6-5	Example of .ENABL and .DSABL Directives	6-16
	6-6	Example of .BLKB and .BLKW Directives	6-30
	7-1	Example of .IRP and .IRPC Directives	7-17

#### TABLES

TABLE	3-1	Special Characters Used in MACRO-11	3-1
	3-2	Legal Separating Characters	3-3
	3-3	Legal Argument Delimiters	3-3
	3-4	Legal Unary Operators	3-4
	3-5	Legal Binary Operators	3-5
	6-1	Symbolic Arguments of Listing Control	
		Directives	6-3
	6-2	Symbolic Arguments of Function Control	
		Directives	6-14
	6-3	Symbolic Arguments of .PSECT Directive	6-33
	6-4	Non-IAS/RSX-11 Program Section Default Values	6-39
	6-5	Legal Condition Tests for Conditional Assembly	
		Directives	6-41
	6-6	Subconditional Assembly Block Directives	6-44

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#### 0.1 MANUAL OBJECTIVES AND READER ASSUMPTIONS

The intent of this manual is to enable users to develop programs coded in the MACRO-11 assembly language. No prior knowledge of the MACRO-11 Relocatable Assembler is assumed.

Although the description of the assembly language is wholly self-contained within this manual, the reader is assumed to be familiar with the PDP-11 processors and related terminology, as presented in the PDP-11 Processor Handbooks. No attempt is made in this document to describe the PDP-11 hardware or the functions of the various PDP-11 instructions.

Since the development of programs necessarily involves linking to create an executable image, the reader is encouraged to become familiar with this process, as presented in the applicable system manual (see Section 0.3).

In presenting MACRO-11, a tutorial bias has been adopted to enlarge upon the reference material. This posture is reflected in the examples and the accompanying commentary describing MACRO-11 language elements in typical applications. Portions of text that are shaded indicate that a particular MACRO-11 feature is not available in the 8K version of MACRO-11.

#### 0.2 STRUCTURE OF THE DOCUMENT

This manual contains three parts. Part I, consisting of two chapters, briefly introduces MACRO-11. Chapter 1 lists the key features of MACRO-11, and Chapter 2 identifies the advantages of following programming standards and conventions. Also described is the format used in coding MACRO-11 source programs.

Part II, consisting of three chapters, presents general information essential to programming with the MACRO-11 assembly language. Chapter 3 describes the symbols, terms, and expressions that form the elements of MACRO-11 instructions. The character set is listed, and the types of programming symbols that may be defined by the user are discussed. Chapter 4 describes the output of MACRO-11 and presents concepts essential to the proper relocation and linking of object modules. Chapter 5 briefly describes how data stored in memory can be accessed and manipulated using the addressing modes recognized by the PDP-11 hardware.

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Part III, consisting of two chapters, describes the MACRO-11 directives that control the processing of source statements during assembly. Chapter 6 discusses directives which accomplish generalized MACRO-11 functions, while Chapter 7 deals with directives used in the definition and expansion of macros.

Finally, several appendixes are provided, supplying additional information of interest to the MACRO-11 programmer.

Appendix A lists the ASCII and Radix-50 character sets that may be used in MACRO-11 programs. Appendix B lists the special characters recognized by MACRO-11, summarizes the syntax of the various addressing modes used in PDP-11 processors, and briefly describes the MACRO-11 directives in alphabetical order. The permanent symbols that have been defined for use with MACRO-11 are listed alphabetically in Appendix C.

The diagnostic error codes produced by MACRO-11 to identify various types of errors detected during the assembly process are listed alphabetically in Appendix D. Appendix E contains a sample coding standard that is recommended practice in preparing MACRO-11 programs. Appendix F discusses several methods of conserving dynamic memory space for users of small systems who may experience difficulty in assembling MACRO-11 programs.

Appendix G is a discussion of position independent code (PIC).

#### 0.3 ASSOCIATED DOCUMENTS

The reader should refer to the applicable documentation directory listed below for descriptions of documents associated with this manual.

IAS Documentation Directory

RSX-11D Documentation Directory

RSX-11M/RSX-11S Documentation Directory

RT-11 Documentation Directory

#### 0.4 DOCUMENT CONVENTIONS

The symbols defined below are used throughout this manual.

Symbol

Definition

[] Brackets indicate that the enclosed argument is optional.

 Wertical bars indicate that a single choice must be made from a list of arguments.

Ellipsis indicates optional continuation of an argument list in the form of the last specified argument. UPPER-CASE Upper-case characters indicate elements of the language that must be used exactly as shown.

lower-case Lower-case characters indicate elements of the language characters that are supplied by the programmer.

(n)

In some instances the symbol (n) is used following a number to indicate the radix. For example, 100(8) indicates that 100 is an octal value, while 100(10) indicates a decimal value. -CASE Upper-dasa characters indicate elements of the language TREAS that must be used exactly as shown.

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INTRODUCTION TO MACRO-11

PART I



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# MACRO-11 FEATURES of anipod box frequencies of anipod box frequencies

The MACRO-11 Assembler provides the following features:

- 1. Source and command string control of assembly functions
- 2. Device and filename specifications for input and output files
- 3. Error listing on command output device
- Alphabetized, formatted symbol table listing; optional cross-reference listing of symbols

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- 5. Relocatable object modules
- 6. Global symbols for linking object modules
- 7. Conditional assembly directives
- 8. Program sectioning directives and an abatration and a section of the section o
  - 9. User-defined macros and macro libraries
  - 10. Comprehensive system macro library
  - 11. Extensive source and command string control of listing functions.

# 1.1 OVERVIEW OF MACRO-11

MACRO-11 is a 2-pass assembler. The functions and operations relevant to each assembly pass are described in the following sections.

# 1.1.1 Assembly Pass 1

The main purpose of assembly pass 1 is to locate and read all required macros from libraries; to build symbol tables and program section tables for the program; while also performing a rudimentary assembly of each source statement.

The first stage of assembly pass l is the initialization of all impure data areas that MACRO-11 uses internally for the assembly process. These areas include all dynamic storage areas and buffer areas used as file storage regions.

After initializing memory areas, MACRO-11 issues a call to a system subroutine which transfers a command line into memory. This command

#### MACRO-11 FEATURES

line contains the specifications of the files to be used during assembly. After scanning the command line for proper syntax, MACRO-11 initializes the specified output files. These files are opened to determine if valid output file specifications have been passed in the command line. They are then closed to minimize requirements for active file space.

As the assembly process begins, MACRO-11 initiates a routine which retrieves source lines from the input file. If no such file is currently open, as is the case at the beginning of assembly, MACRO-11 opens the next input file specified in the command line previously read and begins to assemble the source statements. MACRO-11 determines the length of each instruction and assembles it accordingly as one word, two words, or three words.

At the end of assembly pass 1, MACRO-11 reopens the output files described above and writes out information that is to be used later in linking the object modules. Such information as the object module name, the program version number, and the global symbol directory (GSD) entries for each program section are output to the object file. After writing out the GSD entries for a given program section, MACRO-11 scans through the symbol tables to find all the global symbols that are bound to that particular program section. MACRO-11 then writes out GSD records to the object file for these symbols. This process continues for each program section, bringing to a close assembly pass 1.

#### 1.1.2 Assembly Pass 2

As an integral part of pass 2, MACRO-11 simultaneously writes the object records to the output file and generates the assembly listing, followed by the symbol table listing for the program. A cross-reference listing may also be generated.

Basically, assembly pass 2 consists of the same steps performed in assembly pass 1, except that all source statements containing MACRO-11-detected errors are flagged with an error code as the assembly listing file is created. The object file that is created as the final consequence of pass 2 contains all the object records, together with relocation records containing information necessary for subsequent linking of the object file.

The information thus passed enables the global symbols in the object modules to be associated with absolute or virtual memory addresses, thereby forming an executable body of code.

The user may wish to become familiar with the macro object file format and description. This information is presented in the applicable system manual (see Section 0.3 in the Preface).

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## CHAPTER 2 Mailes Samelie Statistics

#### SOURCE PROGRAM FORMAT

#### 2.1 PROGRAMMING STANDARDS AND CONVENTIONS

Assembly level programming deals directly with the host hardware. Hence, great care must be exercised in establishing programming standards and conventions to enable code written by one group to be interchanged easily with another group. Standards provide a number of advantages. When applied to the program development process, standards make the programming effort easier to:

Plan Comprehend Test Modify Convert

Even though standards must accommodate local requirements, many aspects of the program development process have universal applicability. The standards common to all of DIGITAL's PDP-11 software products are presented in Appendix E as a model for users. Observance of these standards is beneficial to DIGITAL and its users, by simplifying both communications and the continuing task of software maintenance and enhancement.

#### 2.2 STATEMENT FORMAT

A source program is composed of a sequence of source coding lines. Each line contains a single assembly-language statement. MACRO-11 will accept a source line of 132 characters, but 80 characters is the recommended length, because of constraints imposed by listing format and terminal line size.

A MACRO-11 statement may consist of as many as four fields. These fields are identified by their order of appearance within the statement and/or by specified separating characters between fields. The general format of a MACRO-11 statement is:

Label: Operator Operand ;Comment(s)

The label and comment fields are optional. The operator and operand fields are interdependent, i.e., when both fields are present in a source statement, each field is evaluated by MACRO-11 in the context of the other.

A statement may contain an operator field and no operand field, but the reverse is not true. A statement containing an operand with no operator does not conform to established MACRO-11 coding conventions; such a statement is currently interpreted by MACRO-11 during assembly as an implicit .WORD directive (see Section 6.3.2). MACRO-11 interprets and processes source program statements one by one, generating one or more binary instructions or data words, or performing a specified assembly process. Blank lines, although legal, have no significance in the source program.

An assembly-language statement must be completed on one source line; no continuation lines are allowed in MACRO-11.

The tab character can be used in the source statement to format the fields into aligned columns in accordance with DIGITAL's standard source program format, as shown below:

Label - begins in column 1

Operator - begins in column 9

Operand(s) - begin(s) in column 17

eniumsComment(s) - begin(s) in column 33.ed down start to an

For example, the following statement should be formatted in the source program into specific columns, increasing its readability in the assembly listing:

REGTST:BIT#MASK,VALUE;COMPARES BITS IN OPERANDS.

1 9 17	33	(columns)
--------	----	-----------

REGTST: BIT #MASK,VALUE

;COMPARES BITS IN OPERANDS.

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The above formatting conventions are not mandatory in coding MACRO-11 programs (free-field coding is permissible). However, it is recommended that source programs be prepared in accordance with these conventions for consistency and clarity.cord size stouberg erestica

#### 2.2.1 Label Field

A label is a means of symbolically referring to a location in a program.

A label is a user-defined symbol which is assigned the value of the current location counter and entered into the user-defined symbol table. The current location counter is the means by which MACRO-11 assigns memory addresses to the source program statements as they are encountered during the assembly process. The address value of the label is absolute or relocatable, depending on whether the current program section being assembled is absolute correlocatable. (The concept of program sections and the attributes that may be specified for them are discussed in detail in Section 6.8.)

In the case of an absolute program section, the value of the current location counter is likewise absolute, i.e., its value references an absolute virtual memory address (such as location 100). Similarly, the value of the current location counter in a relocatable program section is also relocatable; however, a relocation bias calculated at link time will be added to the apparent value of the current location counter to establish its effective absolute virtual address at execution time.

A statement may contain an operator field and no operand field, but

If present, a label always appears as the first field in a source statement and must be terminated by a colon. For example, if the current location counter value is absolute 100(8), the statement:

ABCD: MOV A,B

assigns the value 100(8) to the label ABCD. Subsequent references to this label would then yield a value of absolute 100(8). In this example, if the location counter value were relocatable, the final value of ABCD would be 100(8)+K, where K represents the relocation bias of the program section, as calculated by the Task Builder at link time.

More than one label may appear within a single label field. Each label so specified is assigned the same address value. For example, if the current location counter value is 100(8), the multiple labels in the following statement:

5.2.2 Operator Field

ABC: \$DD: A7.7: MOV A,B

are each assignet whe value 100(8).

Multiple labels may also appear on successive lines. For example, the statements

ABC: \$DD:

A7.7: MOV A,B

likewise cause the same current location counter value to be assigned to all three labels.

Of the two methods of assigning multiple labels shown above, the second is preferred, because consistency of field positioning within the source program improves readability.

A double colon (::) defines the label as a global symbol. Such a label can be referenced by independently-assembled object modules. References to this label in other modules will be resolved when the modules are linked as a composite executable image. For example, the statement

ABCD:: MOV A,B

establishes the label ABCD as a global symbol. The distinguishing attribute of a global symbol is that it can be referenced from within an object module other than the module in which the symbol is defined (see Section 6.9).

The legal characters for defining labels are:

Α	through	2
0	through	9
	(Period)	
\$	(Dollar	Sign)

By convention, the dollar sign (\$) and period (.) are reserved for use in defining DIGITAL system software symbols. Therefore these characters should not be used in defining labels in MACRO-11 source programs. A label may be any length; however, only the first six characters are significant and, therefore, must be unique among all the labels in the source program. All labels are terminated by a colon (:), which is not considered part of the label. It is a mandatory delimiter. An error code (M) is generated in the assembly listing if the first six characters in two or more labels are the same (see Appendix D).

A symbol used as a label must not be redefined within the source program. If the symbol is redefined, a label with a multiple definition results, causing MACRO-11 to generate an error code (M) in the assembly listing (see Appendix D). Furthermore, any statement in the source program which references a multi-defined label results in an additional diagnostic message; in this case, an error code (D) is generated in the assembly listing (see Appendix D).

#### 2.2.2 Operator Field

The operator field specifies the action to be performed. It may consist of an instruction mnemonic (op code), an assembler directive, or a macro call.

The operator field follows the label field in a source statement. Chapters 6 and 7 describe these three types of operator field entries.

When the operator is an instruction mnemonic, the mnemonic op code specifies the machine instruction to be generated. MACRO-11 then continues with the evaluation of the address(es) of the operand(s) which follow(s). When the operator is a directive, the directive causes MACRO-11 to perform certain control actions or processing operations during the assembly of the source program. When the operator is a macro call, MACRO-11 inserts the code generated by the macro expansion.

The operator field need not be preceded by a label; but it may be preceded by one or more labels and followed by one or more operands and/or a comment. Furthermore, leading and trailing spaces or tabs in the operator field have no significance; such characters serve only to separate the operator field from the preceding and following fields.

An operator is terminated by a space, tab, or any non-RAD50 character, as in the following examples:

A VOM	onys odd	;THE SPACE TERMINATES THE OPERATOR IS ;MOV.	
MOV	A,B	;THE TAB TERMINATES THE OPERATOR MOV.	
MOV@A	,В	;THE @ CHARACTER TERMINATES THE ;OPERATOR MOV.	

Although the statements above are all equivalent in function, the second statement is the recommended form because it conforms to MACRO-11 coding conventions.

By convention, the dollar sign (5) and period (-) are reserved for use in defining DIGITAL system software symbols. Therefore these bills of a set of the set of the

When the operator field contains an instruction mnemonic (op code), the operand field specifies those program variables that are to be evaluated/manipulated by the operator. The operand field may also be used to supply arguments to MACRO-11 directives and macro calls, as described in Chapters 6 and 7, respectively.

Operands may be expressions or symbolic arguments (within the context of the specified operation). Multiple expressions used in the operand field of a MACRO-11 statement must be separated by a comma; multiple symbolic arguments similarly used may be delimited by any legal separator, i.e., a comma, tab, and/or space. An operand should be preceded by an operator field; if it is not, the statement is treated by MACRO-11 as an implicit .WORD directive (see Section 6.3.2).

When the operator field contains an op code, associated operands are always expressions, as shown in the following statement:

MOV R0,A+2(R1)

On the other hand, when the operator field contains a MACRO-11 directive or a macro call, associated operands are normally symbolic arguments, as shown in the following statement:

(32), 十(32);

.MACRO ALPHA ARG1,ARG2

Refer to the description of each MACRO-11 directive to determine the type and number of operands required in issuing the directive.

The operand field is terminated by a semicolon when the field is followed by a comment. For example, in the following statement:

LABEL: MOV A,B ;COMMENT FIELD

the tab between MOV and A terminates the operator field and defines the beginning of the operand field; a comma separates the operands A and B; and a semicolon terminates the operand field and defines the beginning of the comment field. When no comment field follows, the operand field is terminated by the end of the source line.

#### 2.2.4 Comment Field

The comment field normally begins in column 33 and extends through the end of the line. This field is optional and may contain any ASCII characters except null, RUBOUT, carriage-return, line-feed, vertical-tab or form-feed. All other characters appearing in the comment field, even special characters reserved for use in MACRO-11, are checked only for ASCII legality and then included in the assembly listing as they appear in the source text.

All comment fields must begin with the semicolon character(;). When lengthy comments extend beyond the end of the source line (column 80), the comment may be resumed in a following line. Such a line must contain a leading semicolon, and it is suggested that the body of the comment be continued in the same columnar position in which the comment began. A comment line can also be included as an entirely separate line within the code body.

Comments do not affect assembly processing or program execution. However, comments are useful in source listings for later analysis, debugging, or documentation purposes.

# 2.3<sup>02</sup> FORMAT CONTROL METHOD and 10161900 and vd betaluqinam betaulave

Horizontal formatting of the source program is controlled by the space and tab characters. These characters have no effect on the assembly process unless they are embedded within a symbol, number, or ASCII text string, or unless they are used as the operator field terminator. Thus, the space and tab characters can be used to provide an orderly and readable source program, as reflected by the following statements:

LABEL: MOV (SP) +, TAG; POP VALUE OFF STACK.

No spaces or tabs have been used to separate the fields in this statement. Note the difficulty in recognizing where one field ends and the next begins.

LABEL: MOV (SP)+, TAG ; POP VALUE OFF STACK.

This statement conforms to the standard horizontal formatting conventions, i.e., the statement elements are separated into four distinct fields and are therefore easily discernible.

Page formatting and assembly listing considerations are discussed in Chapter 6 in the context of MACRO-11 directives that may be specified to accomplish desired formatting operations. Appendix E describes the coding conventions used in all DIGITAL PDP-11 operating system software.

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#### GARINE TROMMOD (

the tab between MOV and Alterringers the opticator field and definer the beginning of the operand field; a coast apparates the operands A and B; and a semicolog tersinetes the operand field and defines the beginning, of the comment field. When an opment field follows, the operand field is termineted by the option of the source line.

#### blob9 draman0 P.S.S.

The comment field normally bogins is colors -1 and axtends through the end of the line. This field is optimized and may contain any ASOII characters except null. PHOHT, exciption relate, line-feed, vertical-tab or form-face. All other characters appearing in the comment field, even special thereacters recorved for use in MACRO-11, are checked only for ASOII logality and then included in the assembly listing as they appear in the course best.

All comment fields must begin with the secondion character(;). When largely comments extend ocycoid the and of the source line (column 80), the comment may be resuled in a following line. Buch a line must contain a leading semicolon, and it is taggested that the body of the comment be continued in the same columnate position in which the separate line within the code body.

Comments do not affect assembly processing or program execution. Bowever, comments are useful to source listings for later analysis, Gebugging, or documentation purposes. PART II

PROGRAMMING IN MACRO-11 ASSEMBLY LANGUAGE



#### CHAPTER 3

#### SYMBOLS AND EXPRESSIONS

This chapter describes the components of MACRO-11 instructions. The character set, the conventions observed in constructing symbols, and the use of numbers, operators, terms and expressions are discussed as they relate to MACRO-11 programming.

#### 3.1 CHARACTER SET

The following characters are legal in MACRO-11 source programs:

- The letters A through Z. Both upper- and lower-case letters are acceptable, although, upon input, lower-case letters are converted to upper-case (see Section 6.2, .ENABL LC).
- 2. The digits 0 through 9.
- 3. The characters . (period) and \$ (dollar sign). These characters are reserved for use as Digital Equipment Corporation system program symbols.
- 4. The special characters listed in Table 3-1.

Character	Designation	Function
:	Colon	Label terminator.
-bestőhit -tagatos arasad a∎tagatosi tagatosi	Double colon	Label terminator; defines the label as a global label.
- 100 - 100またのを使わせ、現 には成立する	Equal sign	Direct assignment operator; and macro keyword indicator.
== (Herrigin Bilding	Double equal sign	Direct assignment operator; defines the symbol as a global symbol.
8	Percent sign	Register term indicator.
	Tab	Item or field terminator.
	Space	Item or field terminator.

Table 3-1 Special Characters Used in MACRO-11

(Continued on next page)

#### SYMBOLS AND EXPRESSIONS

#### Table 3-1 (Cont.) Special Characters Used in MACRO-11

đ,

-16

Character	Designation	Function	
#	Number sign	Immediate expression indicator.	
e	At sign	Deferred addressing indicator.	
(	Left parenthesis	Initial register indicator.	
) anTacidoutd	Right parenthesis	Terminal register indicator.	
ins elodays poi of bosecciels e	ford Period a croissering for saved	Current location counter	
,	Comma	Operand field separator.	
;	Semicolon	Comment field indicator.	
: ameripord	Left angle bracket	Initial argument or expression indicator.	
aleddol acsorie Starstol sess	Right angle bracket	Terminal argument or expres- sion indicator.	
+ .(35.38)	Plus sign	Arithmetic addition operator or autoincrement indicator.	
aign). These ittal Equipment	Minus sign (Pollar), (Pollar (Pollar), (Pollar) (Pollar systems) (Pollar systems)	Arithmetic subtraction opera- tor or autodecrement indica- stor. and subjects of the subject of	
*	Asterisk .1-0-0000 octoble	Arithmetic multiplication op- erator.	
1	<b>Slash</b>	Arithmetic division operator.	
&	Ampersand	Logical AND operator.	
1	Exclamation point	Logical inclusive OR operator.	
	Double quote	Double ASCII character indica-	
	);BRİMQQC LAAFA		
tr; definds the ladsi is	n diale dial	tor; dued or concatenation indicator.	
sistegs dis .ucistication.	Circumflex	Universal unary operator or argument indicator.	
nt operator; tool as a global	I Backslash We shire the symp icdwys	Macro call numeric argument indicator.	

3.1.1 Separating and Delimiting Characters

Legal separating characters and legal argument delimiters are defined below in Tables 3-2 and 3-3 respectively.

#### SYMBOLS AND EXPRESSIONS

Table 3-2 Legal Separating Characters

Character	Definition	Usage		
Space	One or more spaces and/or tabs	A space is a legal separator between instruction fields and between symbolic arguments within the operand field. Spaces within expressions are ignored (see Section 3.9).		
	Comma	A comma is a legal separator between symbolic arguments within the operand field. Multiple expressions used in the operand field must be separated by a comma.		

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Table 3-3 Legal Argument Delimiters

Character	Definition	n Pil S <b>Usage</b> <sup>M</sup>		
<>	Paired angle brackets	Paired angle brackets may be used anywhere in a program to enclose an expression for treatment as a single term. Paired angle brackets are also used to enclose a macro argument, particularly when that argument contains separ- ating characters (see Section 7.3).		
<b>^xx</b>	Up-arrow (unary oper- ator) construction, where the up-arrow is followed by an argu- ment that is bracketed by any paired printing characters (x).	This construction is equiva- lent in function to the paired angle brackets described above and is generally used only where the argument itself con- tains angle brackets.		

#### 3.1.2 Illegal Characters

A character is determined to be illegal for one of two reasons:

 A character is not an element of the recognized MACRO-11 character set. A character of this kind is replaced in the listing by a question mark, and an error code (I) is printed in the assembly listing (see Appendix D). The exception to this is an embedded null which, when detected, terminates the scan of the current line.  A legal MACRO-11 character is illegal in the context of its usage within the source statement, i.e., its syntax is illegal or questionable. Such a character causes an error code (Q) to be printed in the assembly listing.

#### 3.1.3 Unary and Binary Operators

Legal MACRO-11 unary operators are described in Table 3-4. Unary operators are used in connection with single terms (arguments or operands) to indicate an action to be performed on that term during assembly. A term preceded by a unary operator is considered to contain that operator. The term so specified thus becomes a value which can be used alone or as an element of an expression.

Unary Operator	Explanation	Example	Effect	
+	Plus sign	nieics" +A <sub>leosco</sub> s∧ i	Produces the positive value of A.	
-	Minus sign	-A goid	Produces the negative (2's complement) value of A.	
ed yes a or mary be to mary bo rol and a	Up-arrow, univer- sal unary operator.	^C24	Produces the l's comple- ment value of 24(8).	
ngle term. 6 are also 8 marto	described in detail in Section 6.4.)	^D127	Interprets 127 as a decimal number.	
ing separ- ing separ-	aroument, garticula rhat argument conta abian choractars (se	^F3.0	Interprets 3.0 as a l-word, floating-point number.	
-avispe al	<pre>%_Dil %_Dil %</pre>	^034	Interprets 34 as an octal number.	
svida bedi vico beer	an an anna an an Ailin heackete doson Ailin a generaliy	<b>^B11000111</b>	Interprets 11000111 as a binary number.	
-005 11683 	vor the argument 1 vins angle brackéts	<b>^RABC</b> to be (*x)	Evaluates ABC in Radix-50 form.	

Table 3-4 Legal Unary Operators

Unary operators can be used adjacent to each other or in constructions involving multiple terms, as shown below: areasonad impeting fills

A character is determined to C<^D50) at being the Equivalent to C<^D12 (Equivalent to C<^O12>)

Legal MACRO-11 binary operators are described in Table 3-5. In contrast to unary operators, binary operators specify actions to be performed on multiple items or terms within an expression. Table 3-5 shows the relationships that can be established between expression terms through the use of binary operators.

#### SYMBOLS AND EXPRESSIONS

Table 3-5 Legal Binary Operators

Binary Operator	Explanation	Example	
+ (1 ) (2 ) + (2 ) (3 ) * (3 )	Addition		
<b>-</b> 7.4	Subtraction	<b>А-В</b>	
*	Multiplication	A*B (16-bit product returned)	
	Division	A/B (16-bit quotient returned)	
&	Logical AND	A&B	
	Logical inclusive OR	an an <b>Alb</b> ahar an Albahar an Arabahar an A	

All binary operators have equal priority. Items or terms can be grouped for evaluation within an expression by enclosing them within angle brackets. Terms so enclosed are evaluated first, and remaining operations are performed from left to right, as shown in the examples below:

.WORD	1+2*3	; EQUALS	11(8).
.WORD	1+<2*3>	; EQUALS	7(8).

#### 3.2 MACRO-11 SYMBOLS

Three types of symbols may be defined for use within MACRO-11 source programs: permanent symbols, user-defined symbols, and macro symbols. MACRO-11 maintains three types of symbol tables: the Permanent Symbol Table (PST), the User Symbol Table (UST), and the Macro Symbol Table (MST). The PST contains all the permanent symbols defined within (and thus automatically recognized by) MACRO-11 and is part of the MACRO-11 image. The UST and MST are constructed as the source program is assembled.

#### 3.2.1 Permanent Symbols

Permanent symbols consist of the instruction mnemonics (see Appendix C) and MACRO-11 directives (see Chapters 6 and 7 and Appendix B). These symbols are a permanent part of the MACRO-11 image and need not be defined before being used in the operator field of a MACRO-11 source statement (see Section 2.2.2).

#### 3.2.2 User-Defined and Macro Symbols

User-defined symbols are those symbols treated by the programmer as labels (see Section 2.2.1) or that are equated to a specific value through a direct assignment statement (see Section 3.3) or appear as macro names or dummy arguments. These symbols are added to the User Symbol Table as they are encountered during assembly. Macro symbols are those symbols used as macro names (see Section 7.1). Similarly, these symbols are added to the Macro Symbol Table as they are encountered during assembly. User-defined and macro symbols can be composed of alphanumeric characters, dollar signs (\$), and periods (.) only; any other character is illegal.

#### NOTE

The dollar sign (\$) and period (.) characters are reserved for use in defining Digital Equipment Corporation system software symbols. For example, READ\$ is a file-processing system macro. The user is cautioned not to employ these characters in constructing user-defined symbols or macro symbols in order to avoid possible conflicts with existing or future Digital Equipment Corporation system software symbols.

The following rules govern the creation of user-defined and macro symbols:

- 1. The first character of a symbol must not be a number (except in the case of local symbols; see Section 3.5).
- 2. The first six characters of a symbol must be unique.
- 3. A symbol can be written with more than six legal characters, but the seventh and subsequent characters are checked only for ASCII legality and are not otherwise evaluated or recognized by MACRO-11.
- 4. Spaces, tabs, and illegal characters must not be embedded within a symbol. The legal MACRO-11 character set is defined in Section 3.1.

The value of a symbol depends upon its use in the program. When a symbol appears in the operator field, it may be any one of the three symbol types described above i.e., permanent, user-defined, macro. To determine the value of an operator-field symbol, MACRO-11 searches the symbol tables in the following order:

- 1. Macro Symbol Table
- 2. Permanent Symbol Table
- 3. User-Defined Symbol Table

This search order allows redefinition of Permanent Symbol Table entries as macro symbols. That is, permanent symbols may be used as macro symbols. But the user must keep in mind the sequence in which the search for symbols is performed in order to avoid incorrect interpretation of the symbol's use.

When a symbol appears in the operand field, the User-Defined Symbol Table is searched first, then the Permanent Symbol Table is searched.

Depending on their use in the source program, user-defined symbols have either a local (internal) attribute or a global (external) attribute.

Normally, MACRO-11 treats all user-defined symbols as local, that is, their definition is limited to the module in which they appear. However, symbols can be explicitly declared to be global symbols through one of three methods:

1. Use of the .GLOBL directive (see Section 6.9).

- 2. Use of the double colon (::) in defining a label (see Section 2.2.1).
- 3. Use of the double equal (==) sign in a direct assignment statement (see Section 3.3).

All symbols within a module that remain undefined at the end of assembly are treated as default global references.

#### NOTE

Undefined symbols at the end of assembly are assigned a value of 0 and placed into the user-defined symbol table as undefined default global references. If the .DSABL GBL directive is in effect, however, (see Section 6.2), the automatic global reference default function of MACRO-11 is inhibited, causing the statement containing the undefined symbol to be flagged with an error code (U) in the assembly listing (see Appendix D).

Global symbols provide linkages between independently-assembled object modules within the task image. A global symbol defined as a label, for example, may serve as an entry-point address to another section of code within the image. Such symbols are referenced from other source modules in order to transfer control throughout execution. These global symbols are resolved at link time, ensuring that the resulting image is a logically coherent and complete body of code.

#### 3.3 DIRECT ASSIGNMENT STATEMENTS

A direct assignment statement allows you to equate a symbol to a specific value. When a direct assignment statement is first used to define a symbol, that symbol is entered into the User-Defined Symbol Table. A symbol defined in this manner may be redefined in a subsequent direct assignment statement by assigning a new value to the previously-defined symbol.

The general format for a direct assignment statement is:

symbol=expression

or

symbol==expression

where: expression - can have only one level of forward reference (see 5. below).

- cannot contain an undefined global reference.

A direct assignment statement embodying the double equal (==) sign, as shown above, defines the symbol as global (see Section 6.9).

The following examples illustrate the coding of direct assignment statements:

A=1 ;THE SYMBOL A IS EQUATED TO THE ;VALUE 1.

B=A-1&MASKLOW ;THE SYMBOL B IS EQUATED TO THE ;VALUE OF THE ENTIRE EXPRESSION ;WHICH FOLLOWS.

D=. E: MOV #1,ABLE ;THE SYMBOL D IS EQUATED TO ., AND ;THE LABELS C AND E ARE ASSIGNED A ;VALUE THAT IS EQUAL TO THE LOCATION ;OF THE MOV INSTRUCTION.

The last of the three examples above is provided only to illustrate the performance of MACRO-11 in such situations. See Section 3.6 for a description of the period (.) as the current location counter symbol.

The following conventions apply to the coding of direct assignment statements:

- An equal sign (=) or double equal sign (==) must separate the symbol from the expression defining the symbol's value. Spaces preceding and/or following the direct assignment operators, although permissible, have no significance in the resulting value.
- 2. The symbol being assigned in a direct assignment statement is placed in the label field.
- Only one symbol can be defined in a single direct assignment statement.
- A direct assignment statement may be followed only by a comment field.
- 5. Only one level of forward referencing is allowed, as shown in the following example:
  - X=Y (Illegal forward reference)
  - y=z (Legal forward reference)

Z=1

C:

The above example would result in the generation of an error code (U) in the assembly listing on the line containing the illegal forward reference.

Although one level of forward referencing is allowed for local symbols, a global symbol defined in a direct assignment statement must not contain a forward reference, i.e., the global assignment expression must not itself contain an undefined reference to another symbol. Such a forward reference is illegal, causing an error code (A) to be generated in the assembly listing.

#### SYMBOLS AND EXPRESSIONS

#### 3.4 REGISTER SYMBOLS

The eight general registers of the PDP-11 processor are numbered 0 through 7 and can be expressed in the source program in the following manner:

where % indicates a reference to a register rather than a location. The digit specifying the register can be replaced by any legal, absolute term that can be evaluated during the first assembly pass. Use standard symbolic names for all register references.

The register definitions listed below are automatically assigned by MACRO-11, i.e., these definitions are the normal default values and remain valid for all register references within the source program.

R0=%0	;REGISTER 0 DEFINITION.
R1=%1	;REGISTER 1 DEFINITION.
R2=%2	;REGISTER 2 DEFINITION.
R3=%3	;REGISTER 3 DEFINITION.
R4=84	;REGISTER 4 DEFINITION.
R5=%5	;REGISTER 5 DEFINITION.
SP=%6	STACK POINTER DEFINITION.
PC=%7	; PROGRAM COUNTER DEFINITION.

Note that registers 6 and 7 are given special names because of their unique system functions.

A register symbol may be defined in a direct assignment statement appearing in the program. The defining expression of a register symbol must be a legal, absolute value. Although you can reassign the standard register symbols through the use of the .DSABL REG directive (see Section 6.2), this practice is not recommended. An attempt to redefine a default register symbol without first specifying the .DSABL REG directive to override the normal register definitions causes that assignment statement to be flagged with an error code (R) in the assembly listing. The symbolic default names assigned to the registers, as listed above, are the conventional names used in all DIGITAL-supplied PDP-11 system programs. For this reason, you are well advised to follow these conventions.

All non-standard register symbols must be defined before they are referenced in the source program. A register expression less than 0 or greater than 7 is flagged with an error code (R) in the assembly listing.

The % character may be used with any legal term or expression to specify a register. For example, the statement

CLR %3+1

is equivalent in function to the statement

CLR %4

and clears the contents of register 4.

In contrast, the statement

CLR 4

clears the contents of virtual memory location 4.

#### 3.5 LOCAL SYMBOLS

Local symbols are specially formatted symbols used as labels within a block of coding that has been delimited as a local symbol block. Local symbols are of the form n\$, where n is a decimal integer from 1 to 65535, inclusive. Examples of local symbols are:

1\$ 27\$ 59\$ 104\$

A local symbol block is delimited in one of three ways:

 The range of a local symbol block usually consists of those statements between two normally-constructed symbolic labels (see Figure 3-1). Note that a statement of the form:

ALPHA=expression

is a direct assignment statement (see Section 3.3), but does not create a label and thus does not delimit the range of a local symbol block.

- The range of a local symbol block is normally terminated upon encountering a .PSECT, .CSECT, or .ASECT directive in the source program (see Figure 3-1).
- 3. The range of a local symbol block is delimited through MACRO-11 directives, as follows:

Starting delimiter: .ENABL LSB (see Section 6.2)

Ending delimiter: .ENABL LSB

or

.DSABL LSB (see Section 6.2)

followed by one of: Symbolic label

.PSECT (see Section 6.8.1) .CSECT (see Section 6.8.2) .ASECT (see Section 6.8.2)

Local symbols provide a convenient means of generating labels for branch instructions and other such references within a local symbol block. Using local symbols reduces the possibility of symbols with multiple definitions appearing within a user program. In addition, the use of local symbols differentiates entry-point labels from local labels, since local symbols cannot be referenced from outside their respective local symbol block. Thus, local symbols of the same name can appear in other local symbol blocks without conflict. Local symbols do not appear in cross-reference listings.

#### SYMBOLS AND EXPRESSIONS

Local symbols require less symbol table space than other types of symbols. Their use is recommended. When defining local symbols, use the range from 1\$ to 63\$ first, then the range from 128\$ to 65535\$. Local symbols within the range 64\$ through 127\$, inclusive, can be generated automatically as a feature of MACRO-11. Such local symbols are useful in the expansion of macros during assembly and are described in detail in this context in Section 7.3.5.

Be sure to avoid multiple definitions of local symbols within the same local symbol block. For example, if the local symbol 10\$ is defined two or more times within the same local symbol block, each symbol represents a different address value. Such a multi-defined symbol causes an error code (P) to be generated in the assembly listing.

For examples of local symbols and local symbol blocks as they appear in a source program, see Figure 3-1.

121				2			
122				I PROGRA	M INITI	ALIZATION CO	DE
123				a 🔰 ar bara an gun			
124							
125	000000				,PSECT	XCTPRG, GBL	
126	000000	012700	000000	XCTPRGI	MOV	#IMPURE,RØ	JIMPURE DATA INITIALIZATION
127	000004	005020		15:	CLR	(RØ)+	
128	000006	022700	0000001		CMP	SIMPURT, RU	
129	000012	101374	승규는 승규는 것		8H1	15 0310108	
130						1 . Jacob Maria	
131	000000				PSFCT	YCTPAS. GBL	이 가슴 Nage 이 가슴 영상에 가지 않는 것을 하는 것이 가지 않는 것 같이.
112	0000000	012700	0000001	XCTPASI	MOV	STMPPAS.RO	IPASS INITIAL TRATTON
111	0000004	005020		151	CI R	(80)+	
134	0000004	022700	0000001		CMP	STMPPAT . PA	
134	000000	131374			Aut	15	
135	000012	1013/4			041	1.4	
130					Deset	VETITAL COL	
13/	0000000				.PSELI	ACILINGUE	
138	000000	012700	600600.	XCTLINII	MOV	#IMPLIN, RØ	JLINE INITIALIZATION
139	000004	005020		151	CLR	(RØ)+	
140	000000	022700	0000001		CMP	#IMPLIT, RØ	
141	000012	101374			BHI	15	
142					2125 - EX		

Figure 3-1 Assembly Listing Showing Local Symbol Block

#### 3.6 CURRENT LOCATION COUNTER

The period (.) is the symbol for the current location counter. When used in the operand field of an instruction, it represents the address of the first word of the instruction, as shown in the first example below. When used in the operand field of a MACRO-11 directive, it represents the address of the current byte or word, as shown in the second example below.

A: MOV #.,R0 ;THE PERIOD (.) REFERS TO THE ADDRESS ; OF THE MOV INSTRUCTION.

(The function of the # symbol is explained in Section 5.9.)

SAL=0

.WORD

177535,.+4,SAL ;THE OPERAND .+4 IN THE .WORD ;DIRECTIVE REPRESENTS A VALUE ;THAT IS STORED AS THE SECOND ;OF THREE WORDS DURING ;ASSEMBLY.

Assume that the current value of the location counter is 500. During assembly, MACRO-11 reserves storage in response to the .WORD directive (see Section 6.3.2), beginning with location 500. The operands accompanying the .WORD directive determine the values so stored. The

#### SYMBOLS AND EXPRESSIONS

value 177535 is thus stored in location 500. The value represented by .+4 is stored in location 502; this value is derived as the current value of the location counter (which is now 502), plus the absolute value 4, thereby depositing the value 506 in location 502. Finally, the value of SAL, previously equated to 0, is deposited in location 504.

Figure 3-2 illustrates the result of the example.

LOCATION	CONTENTS
500	177535
502	506
504	0

Figure 3-2 Sample Assembly Results

At the beginning of each assembly pass, MACRO-ll resets the location counter. Normally, consecutive memory locations are assigned to each byte of object data generated. However, the value of the location counter can be changed through a direct assignment statement of the following form:

.=expression

Similar to other MACRO-11 symbols, the current location counter symbol (.) has an attribute of relocatability associated with it: it is either absolute or relocatable, depending on the specific such attribute of the current program section. (A program section and its attributes are defined through the use of the .PSECT directive described in Section 6.8.1.) The existing attribute (or mode) of the current location counter cannot be changed by specifying a defining expression having a different attribute.

Furthermore, such a defining expression must not force the location counter into another program section (.PSECT area), even though the program sections so involved may both be absolute or relocatable. The expression defining the location counter value must not contain a forward reference, i.e., the expression must not contain a reference to a symbol that is not previously defined. Such violations constitute a general assembly error, resulting in an error code (A) in the assembly listing.

Thus, the attribute (or mode) of the current location counter takes on the attribute of the current program section. Therefore, its attribute from program section to program section can be changed only through the program sectioning directives (.PSECT, .ASECT, and .CSECT), as described in Section 6.8.

The following coding illustrates the use of the current location counter:
## SYMBOLS AND EXPRESSIONS

.ASECT ;SET LOCATION COUNTER TO .=500 ;ABSOLUTE 500 (OCTAL). ;THE LABEL "FIRST" HAS THE VALUE FIRST: MOV .+10,COUNT ;500 (OCTAL). ;.+10 EQUALS 510 (OCTAL). THE ;CONTENTS OF THE LOCATION ;510 (OCTAL) WILL BE DEPOSITED ; IN THE LOCATION "COUNT." ;THE ASSEMBLY LOCATION COUNTER .=520 ; NOW HAS A VALUE OF ;ABSOLUTE 520 (OCTAL). ;THE LABEL SECOND HAS THE SECOND: MOV ., INDEX ;VALUE 520 (OCTAL). ;THE CONTENTS OF LOCATION ;520 (OCTAL), THAT IS, THE BINARY ;CODE FOR THE INSTRUCTION ; ITSELF, WILL BE DEPOSITED IN THE ;LOCATION "INDEX." .PSECT

.=.+20			;SET LOCATION COUNTER TO
			RELOCATABLE 20 OF THE
THIRD:	.WORD	0	THE LABEL THIRD HAS THE VALUE OF RELOCATABLE 20.

Storage areas may be reserved in the program by advancing the location counter. For example, if the current value of the location counter is 1000, each of the following statements:

.=.+40

or

.BLKB 40

or

.BLKW 20

reserves 40(8) bytes of storage space in the source program. The .BLKB and .BLKW directives, however, are recommended as the preferred ways to reserve storage space (see Section 6.5.3).

## 3.7 NUMBERS

MACRO-11 assumes that all numbers in the source program are to be interpreted in octal radix, unless otherwise specified. An exception to this is that operands associated with Floating Point Processor instructions and Floating Point Data directives are treated as decimal (see Section 6.4.2). This default radix can be altered with the .RADIX directive (see Section 6.4.1.1). Also, individual numbers can be designated as decimal, binary, or octal numbers through temporary radix control operators (see Section 6.4.1.2).

For every statement in the source program that contains a digit that is not in the current radix, an error code (N) is generated in the assembly listing. However, MACRO-11 continues with the scan of the statement and evaluates each such number encountered as a decimal value.

# SYMBOLS AND EXPRESSIONS

Negative numbers must be preceded by a minus sign; MACRO-11 translates such numbers into two's complement form. Positive numbers may (but need not) be preceded by a plus sign.

A number containing more than 16 significant bits, i.e., greater than 177777(8), is truncated from the left and flagged with an error code (T) in the assembly listing.

Numbers are always considered to be absolute values, i.e., they are not relocatable.

Single-word floating-point numbers may be generated with the <sup>°</sup>F operator (see Section 6.4.2.2) and are stored in the following format:

15	1 <b>14</b> 1619 - 1619 - 1619 1619 - 1619 - 1619 - 1619	7	6	0
Sign	8-bit		7-bi	lt
Bit	Exponent		Mant	issa

Refer to the appropriate PDP-11 Processor Handbook for details of the floating-point number format.

#### 3.8 TERMS

A term is a component of an expression and may be one of the following:

- A number, as defined in Section 3.7, whose 16-bit value is used.
- A symbol, as defined in Section 3.2. Symbols are evaluated as follows:
  - a. A period (.) specified in an expression causes the value of the current location counter to be used.
  - b. A defined symbol is located in the User-Defined Symbol Table (UST) and its value is used.
  - c. A permanent symbol's basic value is used, with zero substituted for the addressing modes. (Appendix C lists all op codes and their values.)
  - d. An undefined symbol is assigned a value of zero and inserted in the User-Defined Symbol Table as an undefined default global reference. If the .DSABL GBL directive (see Section 6.2) is in effect, the automatic global reference default function of MACRO-11 is inhibited, in which case, the statement containing the undefined symbol is flagged with an error code (U) in the assembly listing.
- 3. A single quote followed by a single ASCII character, or a double quote followed by two ASCII characters. This type of expression construction is explained in detail in Section 6.3.3.

- 4. A term may also be an expression enclosed in angle brackets (<>). Any expression so enclosed is evaluated and reduced to a single term before the remainder of the expression in which it appears is evaluated. Angle brackets, for example, may be used to alter the left-to-right evaluation of expressions (as in A\*B+C versus A\*<B+C>), or to apply a unary operator to an entire expression (as in -<A+B>).
  - 5. A unary operator followed by a symbol or number.

## 3.9 EXPRESSIONS

Expressions are combinations of terms joined together by binary operators (see Table 3-5) and which reduce to a 16-bit expression value. The evaluation of an expression includes the determination of its attributes. A resultant expression value may be any one of four types (as described later in this section): absolute, relocatable, external, or complex relocatable.

Expressions are evaluated from left to right with no operator hierarchy rules, except that unary operators take precedence over binary operators. A term preceded by a unary operator is considered to contain that operator. (Terms are evaluated, where necessary, before their use in expressions.) Multiple unary operators are valid and are treated as follows:

-+-A

is equivalent to:

-<+<-A>>

A missing term, expression, or external symbol is interpreted as a zero. A missing or illegal operator terminates the expression analysis, causing an error code (A) or (Q), or both, to be generated in the assembly listing, depending on the context of the expression itself. For example, the expression:

TAG ! LA 177777

is evaluated as

TAG ! LA

because the first non-blank character following the symbol LA is not a legal binary operator, an expression separator (i.e., a comma), or an operand field terminator (i.e., a semicolon or the end of the source line). It should be noted that spaces within expressions are ignored.

The value of an external expression is equal to the value of the absolute part of that expression. For example, the expression EXTERN+A, where "EXTERN" is an external symbol, has a value at assembly-time that is equal to the value of the internal symbol A. This expression, however, when evaluated at link time takes on the resolved value of the symbol EXTERN, plus the value of symbol A.

Expressions, when evaluated by MACRO-11, are determined to be one of four types: absolute, relocatable, external (or global), or complex relocatable. The following distinctions are important:

1. An expression is absolute if its value is fixed. An expression whose terms are numbers and ASCII conversion

characters will reduce to an absolute value. A relocatable expression or term minus a relocatable term, where both elements being evaluated belong to the same program section, are also absolute, since such an expression is reduced to a single term by MACRO-11 upon completion of the expression scan. For example, the expression TAG2-TAG1, where both TAG1 and TAG2 are defined in the same program section, is an absolute expression. Terms that contain labels defined in an absolute section will have an absolute value.

- 2. An expression is relocatable if its value is fixed relative to the base address of the program section in which it appears, but it will have an offset value added at link time. Terms that contain labels defined in relocatable program sections will have a relocatable value; similarly, a period (.) in a relocatable program section, representing the value of the current location counter, will also have a relocatable value.
- .3. An expression is external (or global) if it contains a single global reference (plus or minus an absolute expression value) that is not defined within the current program. Thus, an external expression is only partially defined following assembly and must be resolved at link time.
- 4. An expression is complex relocatable if any of the following conditions applies:
  - It contains a global reference and a relocatable symbol.
  - It contains more than one global reference.
  - It contains relocatable terms belonging to different program sections.
  - The value resulting from the expression has more than one level of relocation. For example, if the relocatable symbols TAG1 and TAG2 associated with the same program section are specified in an expression construction in the form TAG1+TAG2, two levels of relocation would be introduced, since each symbol is evaluated in terms of the relocation bias in effect for the program section.
  - An operation other than addition is specified on an undefined global symbol.
  - An operation other than addition, subtraction, negation, or complementation is specified for a relocatable value.

The evaluation of relocatable, external, and complex relocatable expressions is completed at link time.

# CHAPTER 4

## RELOCATION AND LINKING

The output of MACRO-11 is an object module that must be processed or linked before it can be loaded and executed. Essentially, linking fixes (i.e., makes absolute) the values of external or relocatable symbols in the object module, thus transforming the object module, or several such object modules, into an executable image.

To allow the value of an expression to be fixed at link time, MACRO-11 outputs certain directives in the object file, together with other required parameters. In the case of relocatable expressions in the object module, the base of the associated relocatable program section is added to the value of the relocatable expression provided by MACRO-11. In the case of external expression values, the value of the external term in the expression (since the external symbol must be defined in one of the other object modules being linked together) is determined and then added to the absolute portion of the external expression, as provided by MACRO-11.

All instructions that require modification at link time are flagged in the assembly listing, as illustrated in the example below. The apostrophe (') following the octal expansion of the instruction indicates that simple relocation is required; the letter G indicates that the value of an external symbol must be added to the absolute portion of an expression; and the letter C indicates that complex relocation analysis at link time is required in order to fix the value of the expression.

#### **EXAMPLE:**

005065 000000G	CLR	EXTERN (R5)	;THE VALUE OF THE SYMBOL "EXTERN" IS ;ASSEMBLED AS ZERO AND IS ;RESOLVED AT LINK TIME.
005065 G	CLR	EXTERN+6 (R5)	;THE VALUE OF THE SYMBOL "EXTERN" ;IS RESOLVED AT LINK TIME ;AND ADDED TO THE ABSOLUTE ;PORTION (+6) OF THE EXPRESSION.
005065 000040'	CLR	RELOC (R5)	;ASSUMING THAT THE VALUE OF THE ;SYMBOL "RELOC" IS RELOCATABLE ;40, THE RELOCATION BIAS ;WILL BE ADDED TO THIS VALUE.
005065 000000C	CLR	- <extern+reloc></extern+reloc>	<ul> <li>(R5) ;THIS EXPRESSION IS COMPLEX</li> <li>;RELOCATABLE BECAUSE IT REQUIRES</li> <li>;THE NEGATION OF AN EXPRESSION</li> <li>;THAT CONTAINS A GLOBAL "EXTERN"</li> <li>;REFERENCE AND A RELOCATABLE TERM.</li> </ul>

For a complete description of object records output by MACRO-11, refer to the applicable system manual (see Section 0.3 in the Preface).

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# CHAPTER 5

# ADDRESSING MODES

The program counter (PC) always contains the address of the next word to be fetched, i.e., the address of the next instruction to be executed, or the second or third word of the current instruction.

In order to understand how the address modes operate and how they assemble, the action of the program counter must be understood. The key rule to remember is:

"whenever the processor implicitly uses the program counter (PC) to fetch a word from memory, the program counter is automatically incremented by 2 after the fetch operation is completed."

In the case of 2- or 3-word instructions, the processor uses the PC to fetch the following words as well.

The following symbols are used in describing addressing modes throughout this chapter:

- 1. E is any expression, as defined in Chapter 3.
- 2. R is a register expression, i.e., any expression containing a term preceded by a percent sign (%) or a symbol previously equated to such a term, as shown in the examples below:

R0=%0	;GENERAL	REGISTER	0.
R1=R0+1	;GENERAL	REGISTER	1.
R2=1+%1	;GENERAL	REGISTER	2.

The symbol R may also represent any of the normal default register definitions (see Section 3.4).

- 3. ER is a register expression or an absolute expression in the range 0 to 7, inclusive.
- 4. A is a general addressing specification which produces a 6-bit mode address field, as described in the PDP-11 Processor Handbooks. The addressing specification, A, is described in terms of E, R, and ER, as defined above. Each addressing specification within this section is illustrated using either the single operand instruction CLR or the double operand instruction MOV.

# 5.1 REGISTER MODE

The register itself (R) contains the operand to be manipulated by the instruction.

Format for A: R

Example:

CLR R3

;CLEARS REGISTER 3.

# 5.2 REGISTER DEFERRED MODE

The register (R) contains the address of the operand to be manipulated by the instruction.

Format for A: @R or (ER)

Examples:

CLR	@Rl	;ALL	THESE	INSTRU	JCTIONS	CLEAR
CLR	(Rl)	;THE	WORD A	T THE	ADDRESS	5
CLR	<b>(1)</b>	; CON	FAINED	IN REC	GISTER .	1.

## 5.3 AUTOINCREMENT MODE

The contents of the register (ER) are incremented immediately after being used as the address of the operand (see Note below).

Format for A: (ER)+

Examples:

CLR	(R0)+	; EACH INSTRUCTION CLEARS
CLR	(R4) +	THE WORD AT THE ADDRESS
CLR	(R2)+	;CONTAINED IN THE SPECIFIED
		;REGISTER AND INCREMENTS
		;THAT REGISTER'S CONTENTS
		;BY TWO.

## NOTE

Certain special instruction/address mode combinations, which are rarely or never used, do not operate exactly the same on all PDP-11 processors, as described below.

In the autoincrement mode, both the JMP and JSR instructions autoincrement the register before its use on the PDP-11/40, but not on the PDP-11/45 or 11/10.

In double operand instructions having the addressing form Rn, (Rn) + or Rn, - (Rn), where the source and destination registers are the same, the source operand is evaluated as the autoincremented or autodecremented value, but the destination register, at the time it is used, still contains the originally-intended effective address. In the following example, as executed on the PDP-11/40, Register 0 originally contains 100(8):

# ADDRESSING MODES

MOV	R0, (R0) +	;THE QUANTITY 102 IS MOVED
		;TO LOCATION 100.

MOV R0,-(R0) ;THE QUANTITY 76 IS MOVED ;TO LOCATION 100.

The use of these forms should be avoided, since they are not compatible with the entire family of PDP-11 processors.

An error code (Z) is printed in the assembly listing with each instruction which is not compatible among all members of the PDP-11 family.

# 5.4 AUTOINCREMENT DEFERRED MODE

The register (ER) contains a pointer to the address of the operand. The contents of the register are incremented after being used as a pointer.

Format for A: @(ER)+

Example:

CLR @(R3)+

;THE CONTENTS OF REGISTER 3 POINT ;TO THE ADDRESS OF A WORD TO BE ;CLEARED BEFORE THE CONTENTS OF THE ;REGISTER ARE INCREMENTED BY TWO.

## 5.5 AUTODECREMENT MODE

The contents of the register (ER) are decremented before being used as the address of the operand (see Note above in Section 5.3).

Format for A: -(ER)

Examples:

CLR	- (RÛ)	;DECREMENT THE CONTENTS OF THE SPECI-
		;FIED REGISTER (0, 3, OR 2) BY TWO
CLR	- (R3)	;BEFORE USING ITS CONTENTS
CLR	- (R2)	;AS THE ADDRESS OF THE WORD TO BE
	•	;CLEARED.

#### 5.6 AUTODECREMENT DEFERRED MODE

The contents of the register (ER) are decremented before being used as a pointer to the address of the operand.

Format for A: @-(ER)

Example:

CLR @-(R2)

;DECREMENT THE CONTENTS OF ;REGISTER 2 BY TWO BEFORE ;USING ITS CONTENTS AS A POINTER ;TO THE ADDRESS OF THE WORD TO BE ;CLEARED.

## 5.7 INDEX MODE

The value of an expression (E) is stored as the second or third word of the instruction. The effective address of the operand is calculated as the value of E, plus the contents of register ER. The value E is the offset of the instruction, and the contents of register ER form the base.

Format for A: E(ER)

Examples:

	+2(R1)	;THE EFFECTIVE ADDRESS OF THE WORD ;TO BE CLEARED IS X+2, PLUS THE
MOV R	0,-2(R3)	;CONTENTS OF REGISTER 1. ;THE EFFECTIVE ADDRESS OF THE ;DESTINATION LOCATION IS -2, PLUS

# 5.8 INDEX DEFERRED MODE

An expression (E), plus the contents of a register (ER), yields a pointer to the address of the operand. As in index mode above, the value E is the offset of the instruction, and the contents of register ER form the base.

Format for A: @E(ER)

Example:

CLR @114(R4)

;IF REGISTER 4 CONTAINS 100, THIS ;VALUE, PLUS THE OFFSET 114, YIELDS ;THE POINTER 214. IF LOCATION 214 ;CONTAINS THE ADDRESS 2000, LOCATION ;2000 WOULD BE CLEARED.

# 5.9 IMMEDIATE MODE

Immediate mode allows the operand itself (E) to be stored as the second or third word of the instruction. This mode is assembled as an autoincrement of the PC.

Format for A: #E

Examples:

MOV	#100,R0	; MOVE THE VALUE 100 INTO REGISTER 0.
MOV	<b>#X</b> ,R0	; MOVE THE VALUE OF SYMBOL X INTO
		;REGISTER 0.

The number sign (#) in the MACRO-ll character set has special significance as an addressing mode indicator. When this character appears in the operand field, as shown above, it specifies the immediate addressing mode, indicating to MACRO-ll that the operand itself immediately follows the instruction word.

The operation of this mode can be shown through the first example, MOV #100,R0, which assembles as two words:

Location 20: 0 1 2 7 0 0 Location 22: 0 0 0 1 0 0

Location 24: Next instruction

Note that the source operand (the value 100) is assembled immediately following the instruction word, i.e., as the second word in the instruction. Upon execution of the instruction, the processor fetches the first word (MOV) and increments the PC by 2 so that it points to location 22 (which contains the source operand).

After the next fetch and increment cycle, the source operand (100) is moved into register 0, leaving the PC pointing to location 24 (the next instruction).

## 5.10 ABSOLUTE MODE

Absolute mode is the equivalent of immediate mode deferred. The address expression @#E specifies an absolute address which is stored as the second or third word of the instruction. In other words, the value immediately following the instruction word is taken as the absolute address of the operand. Absolute mode is assembled as an autoincrement deferred of the PC.

Format for A: @#E

Examples:

MOV	@#100,R0	; MOVE THE CONTENTS OF ABSOLUTE
		;LOCATION 100 INTO REGISTER RO.
CLR	@#X	CLEAR THE CONTENTS OF THE LOCATION; WHOSE ADDRESS IS SPECIFIED BY THE SYMBOL X.

The operation of this mode can be shown through the first example, MOV @#100,R0, which assembles as two words:

Location 20: 0 1 3 7 0 0

Location 22: 0 0 0 1 0 0

Location 24: Next instruction

Note that the absolute address 100 is assembled immediately following the instruction word, i.e., as the second word in the instruction. Upon execution of the instruction, the processor fetches the first word (MOV) and increments the PC by 2 so that it points to location 22 (which contains the absolute address of the source operand). After the next fetch and increment cycle, the contents of absolute address 100 (the source operand) are moved into register 0, leaving the PC pointing to location 24 (the next instruction).

## 5.11 RELATIVE MODE

Relative mode is the normal mode for memory references within your program. It is assembled as index mode, using the PC as the index register.

Format for A: E

Examples:

CLR	100	;CLEAI	R ABS	SOLUTE	LOCA	TION	100	
MOV	R0,Y	; MOVE	THE	CONTEN	ITS O	F RE	GISTER	0
		;TO LO	DCAT	ION Y				

In relative mode, the offset for the address calculation is assembled as the second or third word of the instruction. This value is added to the contents of the PC (the base register) to yield the address of the source operand.

The operation of relative mode can be shown with the statement MOV 100,R3, which assembles as two words:

Location 20: 0 1 6 7 0 3

Location 22: 0 0 0 0 5 4

Location 24: Next instruction

Note that the constant 54 is assembled immediately following the instruction word, i.e., as the second word in the instruction. Upon execution of the instruction, the processor fetches the first word (MOV) and increments the PC by 2 so that it points to location 22 (containing the value 54). After the next fetch and increment cycle, the processor calculates the effective address of the source operand by taking the contents of location 22 (the offset) and adding it to the current value of the PC, which now points to location 24 (the next instruction). Thus, the source operand address is the result of the calculation OFFSET+PC = 54+24 = 100(8), causing the contents of location 100 to be moved into register 3.

Since MACRO-11 considers the contents of the current location counter (.) as the address of the first word of the instruction, an equivalent index mode statement is shown below:

MOV 100-.-4(PC),R3

This instruction has a relative addressing mode because the operand address is calculated relative to the current value of the location counter. The offset is the distance (in bytes) between the operand and the current value of the location counter.

# 5.12 RELATIVE DEFERRED MODE

The relative deferred mode is similar in operation to the relative mode above, except that the expression E is used as a pointer to the address of the operand. In other words, the operand following the instruction word is added to the contents of the PC to yield a pointer to the address of the operand.

Format for A: @E

Example:

MOV @X,R0

;RELATIVE TO THE CURRENT VALUE OF ;THE PC, MOVE THE CONTENTS OF THE ;LOCATION WHOSE ADDRESS IS POINTED ;TO BY LOCATION X INTO REGISTER 0.

# 5.13 SUMMARY OF ADDRESSING FORMS

Each PDP-11 instruction takes at least one word. Operands of the form listed below do not increase the length of an instruction.

Form	Meaning
R	Register mode
@R or (ER)	Register deferred mode (see Note below)
(ER) +	Autoincrement mode
@(ER)+	Autoincrement deferred mode
-(ER)	Autodecrement mode
@-(ER)	Autodecrement deferred mode

Operands of the following forms add one word to the instruction length for each occurrence of an operand of that form:

Form	Meaning
E(ER)	Index mode
@E(ER)	Index deferred mode
#E	Immediate mode
@#E	Absolute mode (see Note below)
E	Relative mode
@E	Relative deferred mode

The syntax of the addressing modes is summarized in Appendix B. Additional discussion of addressing modes is provided in the applicable PDP-11 Processor Handbook.

#### NOTE

An alternate form for @R is (ER). However, the form @(ER) is only logically, but not physically equivalent to the expression @0(ER). The addressing form @#E differs from form E in that the second or third word of the instruction contains the absolute address of the operand, rather than the relative distance between the operand and the PC. Thus, the instruction CLR @#100 clears absolute location 100, even if the instruction is moved from the point at which it was assembled. See the description of the .ENABL AMA function in Section 6.2, which causes all relative mode addresses to be assembled as absolute mode addresses.

# 5.14 BRANCH INSTRUCTION ADDRESSING

The branch instructions are 1-word instructions. The high-order byte contains the operator, and the low-order byte contains an 8-bit signed offset (seven bits, plus sign), which specifies the branch address relative to the current value of the PC. The hardware calculates the branch address as follows:

- 1. Extends the sign of the offset through bits 8-15.
- Multiplies the result by 2, creating a byte offset rather than a word offset.
- 3. Adds the result to the current value of the PC to form the effective branch address.

MACRO-11 performs the reverse operation to form the word offset from the specified address. Remember that when the offset is added to the current value of the PC, the PC is pointing to the word following the branch instruction; hence, the factor -2 in the following calculation:

Word offset = (E-PC)/2 truncated to eight bits.

Since the value of the PC = .+2, we have:

Word offset = (E-.-2)/2 truncated to eight bits.

In using branch instructions, you must exercise care to avoid the following error conditions:

- 1. Branching from one program section to another;
- Branching to a location that is defined as an external (global) symbol; or
- 3. Specifying a branch address that is out of range, i.e., the branch offset is a value that does not lie within the range -128(10) to +127(10).

The above conditions cause an error code (A) to be generated in the assembly listing for the statement in error.

# 5.15 USING TRAP INSTRUCTIONS

The EMT and TRAP instructions do not use the low-order byte of the instruction word, allowing information to be transferred to the trap handlers in the low-order byte. If the EMT or TRAP instruction is followed by an expression, the value of the expression is stored in the low-order byte of the word. However, if the expression is greater than 377(8), it is truncated to eight bits and an error code (T) is generated in the assembly listing.

# PART III

# MACRO-11 DIRECTIVES

Chapters 6 and 7 describe all the directives used with MACRO-11. Directives are statements that cause MACRO-11 to perform certain operations during assembly. Chapter 6 describes several types of directives, including those which control symbol interpretation, listing header material, program sections, data storage formats, and assembly listings. Chapter 7 describes those directives concerning macros, macro arguments, and repetitive coding sequences.

MACRO-11 directives can be preceded by a label (subject to any restrictions associated with specific directives) and followed by a comment. A MACRO-11 directive occupies the operator field of a source statement. Only one directive can be included in any given source line. The operand field may be occupied by one or more operands or left blank; legal operands differ with each directive specified.



# CHAPTER 6

# GENERAL ASSEMBLER DIRECTIVES

This category of directives includes:

- 1. Listing control
- 2. Function control
- 3. Data storage
- 4. Radix and numeric control
- 5. Location counter control
- 6. Terminators
- 7. Program boundaries
- 8. Program sectioning
- 9. Symbol control
- 10. Conditional assembly
- 11. PAL-11R conditional assembly.

Each is described in its own section of this chapter.

## 6.1 LISTING CONTROL DIRECTIVES

Listing control directives control the content, format, and pagination of all line printer and teleprinter listing output generated during assembly. Facilities also exist for creating object module names and other identification information in the listing output.

# 6.1.1 .LIST and .NLIST Directives

Listing control options can be specified in the text of a MACRO-11 program through the .LIST and .NLIST directives. These directives are of the form:

.LIST .LIST arg .NLIST .NLIST arg where: arg represents one or more of the optional symbolic arguments defined in Table 6-1.

As indicated above, the listing control directives may be used without arguments, in which case the listing directives alter the listing level count. The listing level count is initialized to zero. At each occurrence of a .LIST directive, the listing level count is incremented; at each occurrence of an .NLIST directive, the listing level count is decremented. When the listing level count is negative, the listing is suppressed (unless the line contains an error). Conversely, when the listing level count is greater than zero, the listing is always generated. Finally, when the count is zero, the line is either listed or suppressed, contingent upon the other listing controls currently in effect for the program. For example, the following macro definition employs the .LIST and .NLIST directives to selectively list portions of the macro body when the macro is expanded:

	MACDO	TURCU		יודכת שבט	2 TP					
	.MACKO	11091		10101 100		~ ~ ~ ~ ~ ~ ~ ~		A		
;	A-THIS LINE SI	HOULD I	LIST	;LISTING	LEVEL	COUNT	$\mathbf{1S}$	0.		
	.NLIST			;LISTING	LEVEL	COUNT	IS	-1.		
;	B-THIS LINE SI	HOULD N	NOT LIST							
	.NLIST			;LISTING	LEVEL	COUNT	IS	-2.		
;	C-THIS LINE SH	HOULD N	NOT LIST							
	.LIST			;LISTING	LEVEL	COUNT	IS	-1.		
;	D-THIS LINE SI	HOULD N	NOT LIST							
	.LIST			;LISTING	LEVEL	COUNT	IS	0.		
;	E-THIS LINE SH	HOULD I	LIST	;LISTING	LEVEL	COUNT	IS	BACK	то	0.
	.ENDM									
	•									
	•									
	•									
	.LIST	ME		;LIST MAG	CRO EXE	PANSION	1.			
	LTEST			;CALL THE	E MACRO	<b>)</b>				
;	A-THIS LINE SI	HOULD I	LIST	;LISTING	LEVEL	COUNT	IS	0.		

An important purpose of the level count is to allow macro expansions to be listed selectively and yet exit with the listing level count restored to the value existing prior to the macro call.

When used with arguments, the listing directives do not alter the listing level count; however, the .LIST and .NLIST directives can be used to override current listing control, as shown in the example below:

.MACRO	XX
• · · ·	
.LIST	
.NLIST	
•	

; E-THIS LINE SHOULD LIST

;LIST NEXT LINE.

;DO NOT LIST REMAINDER OF MACRO ;EXPANSION.

;LISTING LEVEL COUNT IS BACK TO 0.

. ENDM

.NLIST ME XX ;DO NOT LIST MACRO EXPANSIONS.

X=.

X = .

The symbolic arguments allowed for use with the listing directives are described in Table 6-1. These arguments can be used singly or in combination with each other. If multiple arguments are specified in a listing directive, each argument must be separated by a comma, tab, or

space. For any argument not specifically included in a listing control statement, the associated default assumption (List or No list) is applicable throughout the source program. The default assumptions for the listing control directives also appear in Table 6-1.

Table 6-1

Argument	Default	Function
SEQ*	List	Controls the listing of source line sequence numbers. MACRO-11 assigns sequence number 1 to the first source line in a file, and increments the sequence number for each additional line in the file. If this field is suppressed through an .NLIST SEQ
		directive, MACRO-11 generates a tab, effectively allocating space for the field, but fills the field with blanks. Thus, the inter-positional relationships of subsequent fields in the listing remain undisturbed. During the assembly process, MACRO-11 examines each source line for possible error conditions. For any line in error, an appropriate error flag is printed preceding the line sequence number field (see Appendix D). MACRO-11 does not assign sequence numbers for files that have had sequence numbers assigned by other programs, such
LOC*	List	as an editor.
		location counter field. Normally, this field is not suppressed. However, if it is suppressed through the .NLIST LOC directive, MACRO-11 does not generate a tab, nor does it allocate space for the field, as is the case with the source line sequence number field (SEQ) described above. Thus, the suppression of the current location counter (LOC) field effectively left-justifies all subsequent fields (while preserving inter-positional relationships) to that position otherwise normally occupied by this field.
BIN*	List	Controls the listing of generated binary code. If this field is suppressed through an .NLIST BIN directive, left-justification of the source code field occurs in the same manner described above for the current location counter (LOC) field.
BEX	List	Controls the listing of binary extensions, i.e., the locations and binary contents beyond those that will fit on the source statement line. This is a subset of the BIN argument.
SRC*	List	Controls the listing of source lines.

# Symbolic Arguments of Listing Control Directives

(Continued on next page)

# Table 6-1 (Cont.) Symbolic Arguments of Listing Control Directives

Argument	Default	Function
СОМ	List	Controls the listing of comments. This is a subset of the SRC argument. The .NLIST COM directive reduces listing time and space when comments are not desired.
MD	List	Controls the listing of macro definitions and repeat range expansions.
MC	List	Controls the listing of macro calls and repeat range expansions.
ME	No list	Controls the listing of macro expansions.
MEB	No list	Controls the listing of macro expansion binary code. A .LIST MEB directive causes only those macro expansion statements that generate binary code to be listed. This is a subset of the ME argument.
CND	List	Controls the listing of unsatisfied conditional coding and associated.IF and .ENDC directives in the source program. This argument permits conditional assemblies to be listed without including unsatisfied conditional coding.
LD	No list	Controls the listing of all listing directives having no arguments, i.e., those listing directives that alter the listing level count.
TOC	List	Controls the listing of the table of contents during assembly pass 1 (see Section 6.1.4 describing the .SBTTL directive). This argument does not affect the printing of the full assembly listing during assembly pass 2.
SYM	List	Controls the listing of the symbol table resulting from the assembly of the source program.
ТТМ	List	Controls the listing output format. The default can be set by the system manager. If the system manager does not set a default, it is set to line printer format. Figure 6-1 illustrates the line printer output format. Figure 6-2 illustrates the teleprinter output format.

\* If the .NLIST arguments SEQ, LOC, BIN, and SRC are in effect at the same time, i.e., if all four significant fields in the listing are to be suppressed, the printing of the resulting blank line is inhibited.

An example of an assembly listing, as sent to a 132-column line printer, is shown in Figure 6-1. Note that binary extensions for statements generating more than one word are formatted horizontally on the source line.

An example of an assembly listing, as sent to a teleprinter (in the same format as for an 80-column line printer), is shown in Figure 6-2. Notice that binary extensions for statements generating more than one word are printed on subsequent lines. There is no explicit truncation of output to 80 characters by the assembler.

Any argument specified in a .LIST/.NLIST directive other than those listed in Table 6-1 causes the directive to be flagged with an error code (A) in the assembly listing.

The listing control options can also be specified at assembly time through switches included in the command string to MACRO-11 (see the appropriate system manual). The use of these switches overrides all corresponding listing control (.LIST or .NLIST) directives specified in the source program.

PAGE	
15:47	
09-JUL-74	
M0707	
MACRO	
ND CSI2	INES
OF CSI1 A	COMMAND L
- TEST	PARSE
CSITST -	READ ANC

ß	RSE COMMAN	
PAGE	AND PA	X
15:47	READ	
	, SBTTL	S MIS
50 20204		GETINE
ES ES		
UNAMMI NAMAND		
50		-
R SE		230

19 10 10 10 10 10 10 10 10 10 10 10 10 10	103903 103964 103964 016946	0 0 0 0 0 0 0 0		GETLN: 15:	• SHTTL GCMLS ECCT EXITS CST51 MOV	READ AND PARSE COMMAND LINES #GCLBLK JGET LINE VIA GCML 15 JSKIP IF NO ERROR JELSE, EXIT G.CMLD+2(R0),G.CMLD(R0),#10 JSEND OUT THE INPUT L #GSIBLK,GCLBLK+G.CMLD+2,GCLBLK+G.CMLD 25 JSL JO+2(R0),-(SP) JPUT STRING EAROR ADDR IN STK
001332 0013350 0013560 001464 001412 001412 001459 001454	166416 066668 162668 166868 166868 08855	000004 000016 000002 00002	୪୦୪୦୦ ୦୦୦୦୦ ୦୦୦୦୦ ୦୦		84440004400 9740007440 9760007440 9760007440 97400007440 97400000000000000000	C.CMLD+2(H0), (3P) ; CALCULATE LENGTH OF FIRST PAR C.CMLD+2(H0), (3P) ; SEND OUT FIRST PART OF STRIN C.FLLD+2(H0), (SP), #15 ; SEND OUT FIRST PART OF ST C.FLLD+2(H0), C.FLLD+2(H0) ; FIS ; SEND OUT SECOND PART C.FLLD(H0), C.FLLD+2(H0) ; CALC ADDR OF LAST PART OF ST C.FLLD(H0), C.CMLD(H0) ; CALC LENGTH OF LAST PART C.FLLD(H0), C.CMLD(H0) ; CALC LENGTH OF LAST PART C.FLLD(H0), C.CMLD(H0) ; CALC LENGTH OF LAST PART C.FLLD+2(H0), C.CMLD(H0) ; SEND OUT LAST PART C.FLLD+2(R0), C.CMLD(H0), #40 ; SEND OUT LAST PART STLLD+2(R0), C.CMLD(H0), #40 ; SEND OUT LAST PART
9 001475 9 001502 9 001504 9 001528 9 001528 9 001528 9 001528 9 001528 9 001528 9 001528 9 001528 9 001528	005760 001652 1127667 1327667 001402 105267 105267 103441	000002 000050 100040 176416	176432 000001	25: 105: 0Parse:	142 142 142 144 144 144 144 144	C_CMLD(R0) ICHECK LENGTH DF LINE GETLN IF NULL, SKIP BACK FOR NEXT LINE HID, EQUBIT JASUME EQUAL SIGN NOT FOUND HID, EQUBIT JASUME EQUAL SIGN NOT SEN SKIP IF EQUAL SIGN NOT SEN ISS JSKIP IF EQUAL SIGN NOT SEN SKIP IF EQUAL SIGN NOT SEN EQUBIT JSKIP IF EQUAL SIGN FOUND EQUBIT JSKIP IF EQUAL SIGN FOUND EQUAL JASTUS MESSAGE OFT-40 JSEND OUTPUT SCAN MESSAGE INIT2 JINIT LOCNS FOR CSI2 CALL/TEST JUNIT2 JSKIP ON ERROR CS2ERR JSKIP ON ERROR
991622 991626 991634 991663 991665 991796 991796 991712 991712	132760 801360 103477 132760 801360 801360	070020 000020 000020	000001 000001	I PARSE :	CALL BB14L BB14L CALTE BB14L CACTE BB14L B	EVALUB JEVALUATE RESULTS OF SEMANTIC PARSE MCS,MOR,C.STAT(R0) JADDITIONAL OUTPUT SPECS7 OPARSE JYES, CONTINUE WITH OUTPUT SCAN IT 40 JEND INUT SCAN MESSAGE INIZ JINIT LOCKS FOR CSIZ CALL/TEST JINIT JONUT SPEC INUT PARSE INPUT SPEC JINPUT,#SWTBL JPARSE INPUT SPEC CSZERR JSKIP ON EHOOR EVALUG JEVALUG JEADOTTIONAL INPUT SPECS7 IPARSE JYES, CONTINUE WITH INPUT SCAN JMPGET JGET ANOTHER COMMAND LINE

Figure 6-1 Example of Line Printer Assembly Listing

GENERAL ASSEMBLER DIRECTIVES

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CSITST -- TEST OF CSI1 AND CSI2 MACRO M0707 09-JUL-74 15:59 PAGE 5 READ AND PARSE COMMAND LINES

209				SBITL	READ AND PARSE	COMMAND LINES
211	001230		GETLNI	GCMLS	#GCLBLK	JGET LINE VIA GCML
212	001244	103003		BCC	15	ISKIP IF NO ERROR
213	001246			EXITSS	-	IFLSE, FXIT
214	001254		151	TYPE	G.CHLD+2(RR).G.	CMLD(R0).#10 JSEND OUT THE INPUT LINE
214	001200		1	CSISI	#CSTBLK.GCLBLK+	G.CMI D+2.GEL BI K+G.CMI D
216	001324	103064		нсс	25	IBRANCH IF NO FRROR DETECTED
217	001324	016046		MOV	C. FTLD+2(RA) (	SP) IPUT STRING FROM ADDR IN STK
£1/	001320	000020		0.00		
218	001332	166016		SUB	C.CMLD+2(RØ),(S	P) JCALCULATE LENGTH OF FIRST PART
		000004				an 网络哈雷尔马马马马马尔马马利亚尔马利亚尔马马马马马马
219	001336			TYPE	$C_{CMLD+2(RØ)}$ (S	P),#'S JSEND OUT FIRST PART OF STRING
220	001360			TYPE	C.FILD+2(R0),C.	FILD(RØ),#'S ISEND OUT SECOND PART
221	001404	066060		ADD	C.FILD(R0),C.FI	LD+2(RØ) JCALC ADDR OF LAST PART OF STRING
		000016				
222	001412	162660		SUB	(SP)+.C.CMIDCRO	DEDUCT LENGTH OF FIRST PART
	001-12	888882				
223	001416	166060		SUB	C. FTLD(RA) . C. CM	U.D.(RO) ICALC LENGTH OF LAST PART
665	001-10	000016		000		
		0000000				
224	001424	000002		TYPE	C. FTI D+2(PA) .C.	CHID (RA) #40 ISEND OUT LAST PART
224	001424			TYPEN	STY 40	ISEND SYNTAX FREDR MESSAGE
225	001450	000655		80	GETIN	TRY FOR MORE
220	0014/4	0000000		<b>D</b> A	02121	
221	001476	005760		T 9 T	C CHIDIPAN	CHECK LENGTH OF LINE
220	0014/0	005700	231	131	C.C.C.C.(KD)	
		000002		850	CETI N	TE NULL SKTP BACK FOR NEYT LINE
229	001502	001002		NOVE	HIA FOURTT	TASSIME FOULL STON NOT FOUND
230	001504	112/0/		HUYD	- DIEGODII	TROSDIC EGORE SIGN NOT FOUND
		176430				
		170432		DTTD	TO FOU C STAT	DAN ACHECK STATUS
231	001215	132/00		0110	#C3, L40, C, 31 A1(	RUJ JCHECK STRIDS
		000040				
		000001		NEO.	1.0.0	PRID TE EDUAL STON NOT SEEN
232	001520	001402		THEO	103	FLOE TNOTCATE FOUND STON FOUND
233	001522	10526/		INCO	EGUBII	JELSE, INDICATE EQUAL SIGN FOUND
		1/0410			5.011 4.0	
234	001526		1051	TYPEM	E00,40	JSEND EQUAL SIGN STATUS MESSAGE
235	001552			TYPEM	UP1,40	JSENU UUIPUI SLAN MESSAGE
236	001576		OPARSES	LALL	INITZ	FINIT LULNS FUR USIZ LALL/IEST
237	001002			LS132	, UUTPUT, #SWIBL	JPARSE DUIPUT SPEL
238	001620	103441		BCS	CSZERR	JSKIP UN ERRUR
239	001622			CALL	EVALUB	JEVALUATE RESULTS OF SEMANTIC PARSE
240	001626	132760		BITB	#CS.MOR,C.STATO	(RØ) JADDITIONAL OUTPUT SPECS?
		000020				
		000001				
241	001634	001360		BNE	OPARSE	JYES, CONTINUE WITH OUTPUT SCAN
242	001636			TYPEM	IPT, 40	ISEND INPUT SCAN MESSAGE
243	001662		IPARSE:	CALL	INIT2	INIT LOCNS FOR CSI2 CALL/TEST
244	001666			CSI\$2	, INPUT, #SWTBL	IPARSE INPUT SPEC
245	001704	103407		BCS	CS2ERR	ISKIP ON ERROR
246	001706			CALL	EVALUB	JEVALUATE RESULTS OF SEMANTIC PARSE
247	001712	132760		8178	#CS. MOR, C. STATO	(RØ) JADDITIONAL INPUT SPECS?
		000020				
		000001				
248	001720	001360		BNE	IPARSE	IYES, CONTINUE WITH INPUT SCAN

Figure 6-2 Example of Terminal Assembly Listing

6-7

Figure 6-3 shows a listing, produced in line printer format, reflecting the use of the .LIST and .NLIST directives in the source program and the effects such directives have on the assembly listing output.

# 6.1.2 Page Headings

MACRO-11 prints each assembly page in the format shown in either Figure 6-1 or Figure 6-2, depending on the listing mode (see TTM, Table 6-1). On the first line of each page, MACRO-11 prints the following (from left to right):

1. Title of the object module, as established through the .TITLE directive (see next section).

*m* 

- 2. Assembler version identification.
- 3. Date.
- 4. Time-of-day.
- 5. Page number.

The second line of each assembly listing page contains the subtitle text specified in the last-encountered .SBTTL directive (see Section 6.1.4).

OM JCOMMENT LINES TEST Om .2,3,4,5 Om Om	COM, BEX> JCOMMENT LINES AND EXTENDED BINARY TEST OM, BEX 22.3.4.5 DM, BEX			VAPROW LISTING MODE IS IN EFFECT	SEQUENCE NUMBERS TEST	THIS IS A COMMENT		EXTENDED BINARY TEST HIS IS A COMMENT		oduced With Listing Control Directives
LSTMAC NLIST . NLIST . NDRD 1 . LIST	LSTMAC .NLIST .NCIST .LIST				-					ting Pr
		46E 1-1		TTM	SEO	1, 2, <b>3</b> , 4, 9	SEG	BEX BEX 1,2,3,4,5	BEX	i-3 List
00003	00003	16129 PJ		LIST.	LSTMAC NI 1ST	20	LIST	LSTMAC .NLIST . WORD	LIST.	igure 6
ଡ଼ଡ଼ଡ଼ଡ଼ ଡ଼େଡ଼ଡ଼ଡ ଡ଼େଡ଼ଡ଼େ	000005									Ч
000001 00000	000001	-60 202				000001 0000001 0000001 0000001 0000001		000001	000001	
888862 888862 898852	000074 900074	NACRO M			000100	000106 000110 000112 000112		898128 888128		
N 8 6 9 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8	32		242	96	80		99	4	44	

瘠

1

3

6-9

- (1 m)						. NLIST . LIST	X T T	JWIDE LISTING MODE IS IN EFFECT JLIST MACRO EXPANSIONS
4 10 4					LISTI	NG CONTR	OL TEST MACRO	
N= 0 0 0 0 0					2000 1000 1000 1000 1000 1000 1000 1000	М А С И А С И А С И А С И А С И И А С И И И И	L 9 T M A C A R G A R G 1,2,3,4,5 A R G	ITHIS IS A COMMENT
m 7 8 4								
	000012 000001	000005	000003		. NL 191	L STMAC L OC 1,2,3,4	LOC 11HIS 1	JLOCATION COUNTER TEST S a comment
8		-				.L1ST	Loc	
50	0,0024					LSTMAC	NIB	JGENERATED BINARY TEST
10	000024			1.2.3.4	1,5	I THIS I LIST	S A COMMENT Bin	
212	000036	8888	0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1		LSTMAC . NLIST MODD	86× 86× • • • • • •	JEXTENDED BINARY TEST 
44							BEX	
9 G V V	000050					LSTMAC	SRC	JSOURCE LINES TEST
	898858 88855 88855	000001 000004	<i><b><i><b>6</b>6666666666666</i></b></i>	00003		,LIST	U S S	
			Figure	6–3 ((	Cont.)	Listi	ng Produced	With Listing Control Directives

.MAIN. MACRO M0707 09-JUL-74 16129 PAGE 1

GENERAL ASSEMBLER DIRECTIVES

6-10

# 6.1.3 .TITLE Directive

The .TITLE directive is used to assign a name to the object module as the first entry in the header of each page in the assembly listing. The name so assigned is the first six non-blank characters following the .TITLE directive. This name should be six Radix-50 characters or less in length; any characters beyond the first six are checked for ASCII legality, but they are not used as part of the object module name. For example, the directive:

# .TITLE PROGRAM TO PERFORM DAILY ACCOUNTING

causes the assembled object module to be named PROGRA. Note that this 6-character name bears no relationship to the filename of the object module, as specified in the command string to MACRO-11. The name of an object module (specified in the .TITLE directive) appears in the load map produced at link time. This is also the module name which the Librarian will recognize.

If the .TITLE directive is not specified, MACRO-11 assigns the default name .MAIN. to the object module. If more than one .TITLE directive is specified in the source program, the last .TITLE directive encountered establishes the name for the entire object module.

All spaces and/or tabs up to the first non-space/non-tab character following the .TITLE directive are ignored by MACRO-11 when evaluating the text string.

If the .TITLE directive is specified without an object module name, or if the first non-space/non-tab character in the object module name is not a Radix-50 character, the directive is flagged with an error code (A) in the assembly listing.

Section A.2 of Appendix A contains a table of Radix-50 characters.

## 6.1.4 .SBTTL Directive

The .SBTTL directive is used to produce a table of contents immediately preceding the assembly listing and to further identify each page in the listing. In the latter case, the text following the .SBTTL directive is printed as the second line of the header of each page in the listing, continuing until altered by a subsequent .SBTTL directive in the program. For example, the directive:

## .SBTTL CONDITIONAL ASSEMBLIES

causes the text

#### CONDITIONAL ASSEMBLIES

to be printed as the second line in the header of the assembly listing.

During assembly pass 1, a table of contents is printed for the assembly listing, containing the line sequence number, the page number, and the text accompanying each .SBTTL directive. The listing of the table of contents is suppressed whenever an .NLIST TOC directive is encountered in the source program (see Table 6-1). An example of a table of contents listing is shown in Figure 6-4.

CSITST -- TEST OF CSI1 AND CSI2 MACRO M0707 09-JUL-74 15:47 TABLE OF CONTENTS

2- 55	MACRO DEFINITIONS
3- 74	MESSAGE STRINGS
4-153	MISCELLANEOUS DATA
5-209	READ AND PARSE COMMAND LINES
6-255	EVALUATE THE SEMANTIC ANALYSIS
7-345	SUBROUTINES

Figure 6-4 Assembly Listing Table of Contents

# 6.1.5 .IDENT Directive

The .IDENT directive provides an additional means of labeling the object module produced by MACRO-11. In addition to the name assigned to the object module with the .TITLE directive (see Section 6.1.3), a character string up to six Radix-50 characters can be specified between paired printing delimiters to label the object module with the program version number. This directive takes the following form:

.IDENT /string/

where:

- string represents six or fewer legal Radix-50 characters which establish the program identification or version number. This number is included in the global symbol directory of the object module; the first four characters are printed in the load map and librarian listing.
  - / represent delimiting characters. These delimiters may be any paired printing characters, other than the equal sign (=), the left angle bracket (<), or the semicolon (;), as long as the delimiting character is not contained within the text string itself. If the delimiting characters do not match, or if an illegal delimiting character is used, the .IDENT directive is flagged with an error code (A) in the assembly listing.

An example of the .IDENT directive is shown below:

## .IDENT /V05A/

The character string V05A is converted to Radix-50 representation and included in the global symbol directory of the object module. This character string also appears in the load map produced at link time and the Librarian directory listings.

When more than one .IDENT directive is encountered in a given program, the last such directive encountered establishes the character string which forms part of the object module identification.

# 6.1.6 .PAGE Directive/Page Ejection

Page ejection is accomplished in one of four ways:

- 1. After reaching a count of 58 lines in the listing, MACRO-11 automatically performs a page eject to skip over page perforations on line printer paper and to formulate teleprinter output into pages. The page number is not changed.
- In addition, the .PAGE directive is used within the source program to perform a page eject at desired points in the listing. The format of this directive is:

. PAGE

This directive takes no arguments and causes a skip to the top of the next page when encountered. It also causes the page number to be incremented and the line sequence counter to be cleared. The .PAGE directive does not appear in the listing.

When used within a macro definition, the .PAGE directive is ignored during the assembly of the macro definition. Rather, the page eject operation is performed as the macro itself is expanded. In this case, the page number is also incremented.

- 3. A page eject is performed when a form-feed character is encountered. If the form-feed character appears within a macro definition, a page eject occurs during the assembly of the macro definition, but not during the expansion of the macro itself. A page eject resulting from the use of the form-feed character likewise causes the page number to be incremented and the line sequence counter to be cleared.
- 4. Encountering a new source file causes the page number to be incremented and the line sequence count to be reset.

# 6.2 FUNCTION DIRECTIVES: .ENABL AND .DSABL

Several function control options are provided by MACRO-11 through the .ENABL and .DSABL directives. These directives are included in a source program to invoke or inhibit certain MACRO-11 functions and operations incidental to the assembly process itself. These directives take the following form:

.ENABL arg .DSABL arg

where: arg represents one or more of the optional symbolic arguments defined in Table 6-2.

Specifying any argument in an .ENABL/.DSABL directive other than those listed in Table 6-2 causes that directive to be flagged with an error code (A) in the assembly listing.

# Table 6-2 Symbolic Arguments of Function Control Directives

Argument	Default	Function
ABS	Disable	Enabling this function produces absolute binary output in FILES-11 format. To convert this output to Formatted Binary format (as required by the Absolute Loader), use the FLX utility.
АМА	Disable	Enabling this function causes all relative addresses (address mode 67) to be assembled as absolute addresses (address mode 37). This function is useful during the debugging phase of program development.
CDR	Disable	Enabling this function causes source columns 73 and greater, i.e., to the end of the line, to be treated as a comment. The most common use of this feature is to permit sequence numbers in card columns 73-80.
CRF	Enable	Disabling this function inhibits the generation of cross-reference output. This function only has meaning if cross-reference output generation is specified in the command string.
FPT	Disable	Enabling this function causes floating- point truncation; disabling this function causes floating-point rounding.
LC	Disable	Enabling this function causes MACRO-ll to accept lower-case ASCII input instead of converting it to upper-case. If this function is not enabled, all text is converted to upper-case.
LSB	Disable	This argument permits the enabling or disabling of a local symbol block. Although a local symbol block is normally established by encountering a new symbolic label or a .PSECT directive in the source program, an .ENABL LSB directive establishes a new local symbol block which is not terminated until (1) another .ENABL LSB is encountered, or (2) another symbolic label or .PSECT directive is encountered following a paired .DSABL LSB directive.
		Although the .ENABL LSB directive permits a local symbol block to cross .PSECT boundaries, local symbols cannot be defined in a program section other than the one that was in effect when the block was entered. The basic function of this directive with regard to .PSECT's is limited to those instances

(Continued on next page)

# Table 6-2 (Cont.) Symbolic Arguments of Function Control Directives

24

2

Argument	Default	Function
LSB (Cont.)	Disable	where it is desirable to leave a program section temporarily to store data, followed by a return to the original program section. Attempts to define local symbols in an alternate program section are flagged with an error code (P) in the assembly listing.
		An example of the .ENABL LSB and .DSABL LSB directives, as typically used in a source program, is shown in Figure 6-5.
PNC	Enable	Disabling this function inhibits binary output until an .ENABL PNC statement is encountered within the same module.
REG	Enable	When specified, the .DSABL REG directive inhibits the normal MACRO-11 default register definitions; if not disabled, the default definitions listed below remain in effect.
		R0=%0 R1=%1 R2=%2 R3=%3 R4=%4 R5=%5 SP=%6 PC=%7
		The .ENABL REG statement may be used as the logical complement of the .DSABL REG directive. The use of these directives, however, is not recommended. For logical consistency, use the normal default register definitions listed above.
GBL	Enable*	When the .ENABL GBL directive is specified, MACRO-11 treats all symbol references that are undefined at the end of assembly pass 1 as default global references; when the .DSABL GBL directive is specified, MACRO-11 treats all such references as undefined symbols. In assembly pass 2, if the .DSABL GBL function is still in effect, these undefined symbols are flagged with an error code (U) in the assembly listing; otherwise, they continue to be regarded by MACRO-11 as global references.

\* The default is Disable for RT-11 MACRO programs.

SQUEEZE MACRO M0707 09-JUL-74 15:13 PAGE 4

		JPUT ADDR OF LINE IN H3	JPOINT R3 PAST LAST CHAR IN LINE	JODES R3 POINT TO START OF LINE?	JIF SO, LEAVE INDICATING FAILURE	JIS THE LAST CHARACTER SEMICOLON?	JND, CONTINUE LOOKING	JYES, POINT R2 PAST NEW END-OF-LINE	JLEAVE VIA COMMON SUCCESS CODE		POINT R2 PAST END-OF-LINE	JODES R2 POINT TO START OF LINE?	JIF SO, LEAVE WITH FAILURE	JIS THE LAST CHARACTER A TAB?	JIF SO, IGNORE IT	JIS IT A BLANK?	JIF SO, IGNORE IT	NON-BLANK CHARACTER-POINT PAST IT	PRE-COMPUTE LINE LENGTH	JINDICATE SUCCESS	JBRANCH TO LEAVE	JINDICATE FAILURE				
	LSB	R1, R3	42,R3	R3, R1	305	-(R3),#SEMIC	15	R3, R2	205		R1, R2	R2, R1	305	-(H2),#TAB	105	(R2),#8LANK	105	R2	R1,R2		405			LSB		
	.ENABL	> 0 H	ADD	CND	BEQ	CMPB	BNE	×04	8 B		A D D	CMP	BEQ	CMPB	BEG	CMPB	BEQ	INC	SUB	CLC CLC	8R	SEC	RETURN	<b>DSABL</b>		
		FNDSMI		151							SKPBLKI	1051							2051			3051	4051			
						000073								110000		000040										
		010103	060203	020301	001422	124327	001373	010302	000412		060102	020201	001412	124227	001773	121227	091770	005202	160102	000241	000401	000261				
		903142	903144	993146	103150	303152	903156	99160	393162		903164	303166	993170	303172	303176	003200	303204	303206	003210	903212	003214	003216	903220			
272	275	278 0	279 8	280	281 0	282	283 6	284 6	285	286	287 6	288	289 6	290	291 6	292 6	293 6	294 6	295 (	296	297	298	299	300	100	302

Example of .ENABL and .DSABL Directives Figure 6-5 \$

## 6.3 DATA STORAGE DIRECTIVES

A wide range of data and data types can be generated with the following directives, ASCII conversion characters, and radix-control operators:



These MACRO-11 facilities are described in the following sections.

## 6.3.1 .BYTE Directive

The .BYTE directive is used to generate successive bytes of binary data in the object module. The directive is of the form:

.BYTE	exp	;STORES THE	BINARY	VALUE	OF THE	
	anda nacan kara	;EXPRESSION	"EXP"	IN THE	NEXT BYTE	•
BYTE	expl,exp2,expn	;STORES THE	BINARY	VALUES	S OF THE I	IST
		;OF EXPRESS	IONS IN	SUCCES	SSIVE BYTE	s.

A legal expression must reduce to eight bits of data or less. The operands of a .BYTE directive are evaluated as word expressions before being truncated to the low-order eight bits. The 16-bit value of the specified expression must have a high-order byte (which is truncated) that is either all zeros (0) or all ones (1). Each expression value is stored in the next byte of the object module. Multiple expressions, which must be separated by commas, are stored in successive bytes, as described below:

SAM=5 .=410

.BYTE ^D48,SAM

;THE VALUE 060 (OCTAL EQUIVALENT OF 48 ;DECIMAL) IS STORED IN LOCATION 410. ;THE VALUE 005 IS STORED IN LOCATION ;411.

If the high-order byte of the expression reduces to a value other than 0 or -1, the value is truncated to the low-order eight bits and flagged with an error code (T) in the assembly listing.

The construction <sup>^</sup>D in the first operand of the .BYTE directive above illustrates the use of a temporary radix-control operator. The function of such special unary operators is described in Section 6.4.1.2.

At link time, it is likely that a relocatable expression will result in a value having more than eight bits, in which case the linker issues a truncation diagnostic for the object module in question. For example, the following statements create such a possibility:

.BYTE 23 ;STORES OCTAL 23 IN NEXT BYTE. A: .BYTE A ;RELOCATABLE VALUE A WILL PROBABLY ;CAUSE TRUNCATION ;DIAGNOSTIC.

If an expression following the .BYTE directive is null, it is interpreted as a zero, as described below:

.=420

.BYTE ,,,

;ZEROS ARE STORED IN BYTES 420, 421, ;422, AND 423.

Note that in the above example, four bytes of storage result from the .BYTE directive. The three commas in the operand field represent an implicit declaration of four null values, each separated from the other by a comma. Hence, four bytes, each containing a value of zero (0), are reserved in the object module.

# 6.3.2 .WORD Directive

The .WORD directive is used to generate successive words of data in the object module. The directive is of the form:

.WORD exp	;STORES THE BINARY EQUIVALENT OF THE
	;EXPRESSION EXP IN THE NEXT WORD.
.WORD expl,exp2,expn	;STORES THE BINARY EQUIVALENTS OF THE
	;LIST OF EXPRESSIONS IN SUCCESSIVE
	; WORDS.

A legal expression must result in 16 bits of data or less. Each expression is stored in the next word of the object program. Multiple expressions must be separated by commas and stored in successive words, as shown in the following example:

SAL=0 .=500

.WORD 177535,.+4,SAL ;STORES THE VALUES 177535, 506, AND ;0 IN WORDS 500, 502, AND 504, ;RESPECTIVELY.

If an expression following the .WORD directive contains a null value, it is interpreted as a zero, as shown in the following example:

.=500

.WORD ,5,

;STORES THE VALUES 0, 5, AND 0 IN ;LOCATION 500, 502, AND 504, ;RESPECTIVELY.

A statement containing a blank operator field, i.e., a symbol that is not recognized by MACRO-11 as a macro call, an instruction mnemonic, a MACRO-11 directive, or a semicolon is interpreted during assembly as an implicit .WORD directive, as shown in the example below:

.=440 LABEL: 100,LABEL

;STORES THE VALUE 100 IN LOCATION 440 ;AND THE VALUE 440 IN LOCATION 442.

#### CAUTION

You should not use this technique to generate .WORD directives because it may not be included in future PDP-11 assemblers.

# 6.3.3 ASCII Conversion Characters

The single quote (') and the double quote (") characters are unary operators that can appear in any MACRO-11 expression. When so used, these characters cause a 16-bit expression value to be generated.

When the single quote is used, MACRO-11 takes the next character in the expression and converts it from its 7-bit ASCII value to a 16-bit expression value. The 16-bit value is then used as an absolute term within the expression. For example, the statement:

## MOV #'A,RO

results in the following 16-bit expression value being moved into register 0:



-Binary Value of ASCII A

Thus, in the example above, the expression 'A results in a value of 101(8). Note that the high-order byte is always zero (0) in the resulting expression value when the single quote unary operator is used.

The ' character must not be followed by a carriage-return, null, RUBOUT, line-feed, or form-feed character; if it is, an error code (A) is generated in the assembly listing.

When the double quote is used, MACRO-11 takes the next two characters in the expression and converts them to a 16-bit binary expression value from their 7-bit ASCII values. This 16-bit value is then used as an absolute term within the expression. For example, the statement:

#### MOV #"AB,R0

results in the following 16-bit expression value being moved into register 0:



Thus, in the example above, the expression "AB results in a value of 041101(8).

The " character also must not be followed by a carriage-return, null, RUBOUT, line-feed, or form-feed character; if it is, an error code (A) is likewise generated in the assembly listing.

The ASCII character set is listed in Section A.l, Appendix A.

# 6.3.4 .ASCII Directive

The .ASCII directive translates character strings into their 7-bit ASCII equivalents and stores them in the object module. The format of the .ASCII directive is as follows:

.ASCII /string l/.../string n/

- is a string of printable ASCII characters. printable ASCII characters are legal. where: string A11 The vertical-tab, null, line-feed, RUBOUT, and all other non-printable ASCII characters, except carriage-return and form-feed, are illegal illegal characters. Such an non-printing character is flagged with an error code (I) in the assembly listing. The carriage-return and form-feed characters terminate the scan of the source line. This premature termination of the .ASCII statement results in the generation of an error code (A) in the assembly listing, because MACRO-11 is unable to complete the scan of the matching delimiter at the end of the character string.
  - / / represent delimiting characters. These delimiters
    may be any paired printing characters, other than
    the equal sign (=), the left angle bracket (<), or
    the semicolon (;), as long as the delimiting
    character is not contained within the text string
    itself. If the delimiting characters do not
    match, or if an illegal delimiting character is
    used, the .ASCII directive is flagged with an
    error code (A) in the assembly listing.</pre>

A non-printing character can be expressed in an .ASCII statement only by enclosing its equivalent octal value within angle brackets. Each set of angle brackets so used represents a single character. For example, in the following statement:

## .ASCII <15>/ABC/<A+2>/DEF/<5><4>

the expressions <15>, <A+2>, <5>, and <4> represent the values of non-printing characters. Furthermore, the expressions must reduce to eight bits of absolute data or less, subject to the same rules for generating data as with the .BYTE directive (see Section 6.3.1).

Angle brackets can be embedded between delimiting characters in the character string, but angle brackets so used do not take on their usual significance as delimiters for non-printing characters. For example, the statement:

# .ASCII /ABC<expression>DEF/

contains a single ASCII character string, and performs no evaluation of the embedded, bracketed expression. This use of the angle brackets is shown in the third example of the .ASCII directive below:
.ASCII	/HELLO/	;STORES THE BINARY REPRESENTATION ;OF THE LETTERS HELLO IN FIVE ;CONSECUTIVE BYTES.
.ASCII	/ABC/<15><12>/DEF/	;STORES THE BINARY REPRESENTATION ;OF THE CHARACTERS A,B,C,CARRIAGE ;RETURN,LINE FEED,D,E,F IN EIGHT ;CONSECUTIVE BYTES.
.ASCII	/A<15>B/	;STORES THE BINARY REPRESENTATION ;OF THE CHARACTERS A, <, 1, 5, >, :AND B IN SIX CONSECUTIVE BYTES.

The semicolon (;) and equal sign (=) can be used as delimiting characters in an ASCII string, but care must be exercised in so doing because of their significance as a comment indicator and assignment operator, respectively, as illustrated in the examples below:

.ASCII	;ABC;/DEF/	;STORES THE BINARY REPRESENTATION OF ;THE CHARACTERS A, B, C, D, E, AND F ;IN SIX CONSECUTIVE BYTES; NOT ;RECOMMENDED PRACTICE.
.ASCII	/ABC/;DEF;	;STORES THE BINARY REPRESENTATIONS OF ;THE CHARACTERS A, B, AND C IN THREE ;CONSECUTIVE BYTES; THE CHARACTERS D, ;E, F, AND ; ARE TREATED AS A COMMENT.
.ASCII	/ABC/=DEF=	;STORES THE BINARY REPRESENTATION ;OF THE CHARACTERS A, B, C, D, E, AND ;F IN SIX CONSECUTIVE BYTES; NOT ;RECOMMENDED PRACTICE.
		RECOMMENDED PRACTICE.

An equal sign is treated as an assignment operator when it appears as the first character in the ASCII string, as illustrated by the following example:

.ASCII	=DEF=		;THE	DIRECT	ASS	IGNMENT	OPER	ATION
			; .ASC	CII=DEF	IS	PERFORM	ED, A	NDAQ
			; (SYN	NTAX) E	RROR	IS GEN	ERATE	D UPON
			;ENCC	DUNTERI	NG T	HE SECO	ND =	SIGN.

# 6.3.5 .ASCIZ Directive

The .ASCIZ directive is equivalent to the .ASCII directive described above, except that a zero byte is automatically inserted as the final character of the string. Thus, when a list or text string has been created with an .ASCIZ directive, a search for the null character in the last byte can effectively determine the end of the string, as reflected by the coding below:

CR=15						
LF=12 HELLO:	.ASCIZ .EVEN	<cr><lf>/MACRC</lf></cr>	0-11 V01A/ <cr><lf> ;INTRODUCTORY MESSAGE</lf></cr>			
	•					
10\$:	MOV MOVB BNE	#HELLO,R1 #LINBUF,R2 (R1)+,(R2)+ 10\$	;GET ADDRESS OF MESSAGE. ;GET ADDRESS OF OUTPUT BUFFER. ;MOVE A BYTE TO OUTPUT BUFFER. ;IF NOT NULL, MOVE ANOTHER BYTE.			
	•					

The .ASCIZ directive is subject to the same checks for character legality and proper character string construction as described above for the .ASCII directive.

# 6.3.6 .RAD50 Directive

1

The .RAD50 directive allows the user to generate data in Radix-50 packed format. Radix-50 form allows three characters to be packed into sixteen bits (one word); therefore, any 6-character symbol can be stored in two consecutive words. The form of the directive is:

.RAD50 /string 1/.../string n/

where: string represents a series of characters to be packed (three characters per word). The string must consist of the characters A through Z, 0 through 9, dollar sign (\$), period (.) and space (). An illegal printing character causes an error flag (Q) to be printed in the assembly listing.

> If fewer than three characters are to be packed, the string is packed left-justified within the word, and trailing spaces are assumed.

> As with the .ASCII directive described in Section 6.3.4, the vertical-tab, null, line-feed, RUBOUT, and all other non-printing characters, except carriage-return and form-feed, are illegal characters, resulting in an error code (I) in the assembly listing. Similarly, the carriage-return and form-feed characters result in an error code (A) because these characters end the scan of the line, preventing MACRO-11 from detecting the terminating matching delimiter.

/ represent delimiting characters. These delimiters
may be any paired printing characters, other than
the equal sign (=), the left angle bracket (<), or
the semicolon (;), provided that the delimiting
character is not contained within the text string
itself. If the delimiting characters do not
match, or if an illegal delimiting character is
used, the .RAD50 directive is flagged with an
error code (A) in the assembly listing.</pre>

Examples of .RAD50 directives are shown below:

.RAD50	/ABC/	;PACKS ABC INTO ONE WORD.
.RAD50	/AB/	;PACKS AB (SPACE) INTO ONE WORD.
.RAD50	/ABCD/	; PACKS ABC INTO FIRST WORD AND
		;D (SPACE) (SPACE) INTO SECOND WORD.
.RAD50	/ABCDEF/	; PACKS ABC INTO FIRST WORD, DEF INTO
		;SECOND WORD.

Each character is translated into its Radix-50 equivalent, as indicated in the following table:

Character	Radix-50	Octal	Equivalent
(space)		0	
À-Z		1-32	non hagenañ Seo a carrege
· \$ ** 한 일 · · · · · · · · · · · · · · · · · ·		33 34	
(undefined)		35	
0-9		36-47	

The Radix-50 equivalents for characters 1 through 3 (C1,C2,C3) are combined as follows:

Radix-50 Value = ((C1\*50)+C2)\*50+C3

For example:

Radix-50 Value of ABC = ((1\*50)+2)\*50+3 = 3223

Refer to Section A.2 in Appendix A for a table of Radix-50 equivalents.

Angle brackets (<>) must be used in the .RAD50 directive whenever special codes are to be inserted in the text string, as shown in the example below:

.RAD50 /AB/<35> ;STORES 3255 IN ONE WORD.

CHR1=1 CHR2=2

CHR3=3

.RAD50 <CHR1><CHR2><CHR3> ;EQUIVALENT TO .RAD50 /ABC/.

#### 6.3.7 Temporary Radix-50 Control Operator: R

The <sup>°</sup>R operator specifies that an argument is to be converted to Radix-50 format. This allows up to three characters to be stored in one word. The <sup>°</sup>R operator is coded as follows:

#### ^Rccc

where ccc represents a maximum of three characters to be converted to a 16-bit Radix-50 value. If more than three characters are specified, any following the third character are ignored. If fewer than 3 are specified, it is assumed that the trailing characters are blanks. The following example shows how the <sup>R</sup> operator might be used to pack a 3-character file type specifier (MAC) into a single 16-bit word.

MOV # RMAC, FILEXT ;STORE RAD50 MAC AS FILE EXTENSION

The number sign (#) is used to indicate immediate data, i.e., data to be assembled directly into object code. ^R specifies that the characters MAC are to be converted to Radix-50. This value is then stored in location FILEXT.

# 6.4 RADIX AND NUMERIC CONTROL FACILITIES

# 6.4.1 Radix Control and Unary Control Operators

The normal default assumption for numeric values or expression values appearing in a MACRO-11 source program is octal. However, numerous instances may occur where an alternate radix is useful for portions of a program or for variables within a given statement. It may be useful, for example, to declare a given radix for applicability throughout a program or to specify a numeric value or expression value in a manner that causes it to be interpreted as a binary, octal, or decimal value during assembly. In other such instances, it may be useful to complement numeric values or expression values. These MACRO-11 facilities are described in the following sections.

#### NOTE

When two or more unary operators appear together, modifying the same term, the operators are applied, from right to left, to the term.

6.4.1.1 .RADIX Directive - Numbers used in a MACRO-11 source program are initially considered to be octal values; however, you can declare any one of the following radices for applicability throughout the source program or within specific portions of the program:

2, 8, 10

This is accomplished via a .RADIX directive of the form:

.RADIX n

n

where:

represents one of the three acceptable radices listed above. If the argument n is not specified, the octal default radix is assumed.

The argument in the .RADIX directive is always interpreted as a decimal value. Any alternate radix declared in the source program through the .RADIX directive remains in effect until altered by the occurrence of another such directive, i.e., a given radix declaration is valid throughout a program until changed. For example, the statement:

.RADIX 10 ;BEGINS A SECTION OF CODE HAVING A ;DECIMAL RADIX.

.RADIX

; REVERTS TO OCTAL RADIX.

Any value other than null, 2, 8, or 10 specified as an argument in the .RADIX directive causes an error code (A) to be generated in the assembly listing.

In general, macro definitions should not contain or rely on radix settings established with the .RADIX directive. Rather, temporary radix control operators should be used within a macro definition. Where a possible radix conflict exists within a macro definition or in possible future uses of that code, it is recommended that the user

specify numeric or expression values using the temporary radix control operators described below.

6.4.1.2 Temporary Radix Control Operators: D, O, and B - Once the user has specified a given radix for a section of code or has decided to use the default octal radix, he may discover a number of cases where an alternate radix is more convenient or desirable (particularly within macro definitions). The creation of a mask word, for example, might best be accomplished through the use of a binary radix.

MACRO-11 has three unary operators that allow the user to establish an alternate radix, as shown below:

D"number"	("number"	is	evaluated	as	a decimal	l number)
^O"number"	("number"	is	evaluated	as	an octal	number)
B"number"	("number"	is	evaluated	as	a binary	number)

Thus, an alternate radix can be declared temporarily to meet a localized requirement in the source program. Such a declaration can be made at any time, regardless of the existence of the default octal radix or another specific radix declaration elsewhere in the program. In other words, the effect of a temporary radix control operator is limited to the term or expression immediately following the operator. Any value specified in connection with a temporary radix control operator is evaluated during assembly as a 16-bit entity. Temporary radix control declarations can be included in the source program anywhere a numeric value is legal.

The expressions below are representative of the methods of specifying temporary radix control operators:

^D123	Decimal radix
^O 47	Octal Radix
^в 00001101	Binary Radix
^O <a+13></a+13>	Octal Radix

Note that the up-arrow and the radix control operator may not be separated, but the radix control operator and the following term or expression can be physically separated by spaces or tabs for legibility or formatting purposes. A multi-element term or expression that is to be interpreted in an alternate radix should be enclosed within angle brackets, as shown in the last of the four temporary radix control expressions above.

The following example also illustrates the use of angle brackets to delimit an expression that is to be interpreted in an alternate radix:

.RADIX 10

A=10

.WORD ^O<A+10>\*10

When the temporary radix expression in the .WORD directive above is evaluated, it effectively yields the following equivalent statement:

.WORD 180.

MACRO-11 also allows a temporary radix change to decimal by specifying a number, immediately followed by a decimal point (.), as shown below:

100.	Equivalent	to	144(8)
1376.	Equivalent	to	2540(8)
128.	Equivalent	to	200(8)

The above expression forms are equivalent in function to those listed below:

^D100 ^D1376 ^D128

# 6.4.2 Numeric Directives and Unary Control Operators

Two storage directives and two numeric control operators are available to simplify the use of the floating-point hardware on the PDP-11. These facilities allow floating-point data to be created in the program, and numeric values to be complemented or treated as floating-point numbers.

1

A floating-point number is represented by a string of decimal digits. The string (which can be a single digit in length) may optionally contain a decimal point, and may be followed by an optional exponent indicator in the form of the letter E and a signed decimal integer exponent. The number may not contain embedded blanks, tabs or angle brackets and may not be an expression. Such a string will result in one or more errors (A or Q) in the assembly listing.

The list of numeric representations below contains seven distinct, valid representations of the same floating-point number:

3 3.0 3.0E0 3E0 .3E1 300E-2

As can be inferred, the list could be extended indefinitely (e.g., 3000E-3, .03E2, etc.). A leading plus sign is optional (e.g., 3.0 is considered to be +3.0). A leading minus sign complements the sign bit. No other operators are allowed (e.g., 3.0+N is illegal).

All floating-point numbers are evaluated as 64 bits in the following format:

 64
 63
 56
 55
 0

S EEEEEEEE MMM....MMM

Mantissa	(55 bits)
Exponent	(8 bits)
Sign	(1 bit)

MACRO-11 returns a value of the appropriate size and precision via one of the floating-point directives. The values returned may be truncated or rounded (see Section 6.2).

Floating-point numbers are normally rounded. That is, when a floating-point number exceeds the limits of the field in which it is to be stored, the high-order bit of the unretained word is added to the low-order bit of the retained word, as shown below. For example, if the number is to be stored in a 2-word field, but more than 32 bits are needed to express its exact value, the highest bit (32) of the

unretained field is added to the least significant bit (0) of the retained field (see illustration below). The .ENABL FPT directive is used to enable floating-point truncation; .DSABL FPT is used to return to floating-point rounding (see Table 6-2).

Bit		Bit Bit	Bit
32		0 32 31	0
Ref	etained ield	Unretain field	ned

Note that all numeric operands associated with Floating Point Processor instructions are automatically evaluated as single-word, decimal, floating-point values unless a temporary radix control operator is specified. For example, to add (floating) the octal constant 41040 to the contents of floating accumulator zero, the following instruction must be used:

ADDF #^041040,F0

where: F0 is assumed to represent floating accumulator zero.

Floating-point numbers are described in greater detail in the applicable PDP-11 Processor Handbook.

6.4.2.1 .FLT2 and .FLT4 - Floating-Point Storage Directives - MACRO-11 supports two directives that evaluate successive floating-point numbers and store the results in the object module. These directives are similar to the .WORD directive and are of the form:

.FLT2 argl,arg2,... .FLT4 argl,arg2,...

where: argl,arg2,... represent one or more floating point numbers as described in Section 6.4.2. Multiple arguments must be separated by commas.

.FLT2 causes two words of storage to be generated for each argument, while .FLT4 generates four words of storage for each argument.

6.4.2.2 Temporary Numeric Control Operators:  $^{\circ}C$  and  $^{\circ}F$  - The  $^{\circ}C$  unary operator allows you to specify an argument that is to be complemented as it is evaluated during assembly. The  $^{\circ}F$  unary operator allows you to specify an argument consisting of a l-word floating-point number.

As with the radix control operators described above, the numeric control operator (°C) can be used anywhere in the source program that an expression value is legal. Such a construction is evaluated by MACRO-11 as a 16-bit binary value before being complemented. For example, the following statement:

TAG4: .WORD C151

causes the l's complement of the value 151 (octal) to be stored as a 16-bit value in the program. The resulting value expressed in octal form is 177626(8).

Because the <sup>°</sup>C construction is a unary operator, the operator and its argument are regarded as a term. Thus, more than one unary operator may be applied to a single term. For example, the following construction:

^C^D25

causes the decimal value 25 to be complemented during assembly. The resulting binary value, when expressed in octal form, reduces to 177746(octal).

The term created through the use of the temporary numeric control operator thus becomes an entity that can be used alone or in combination with other expression elements. For example, the following construction:

.B.

^C2+6

is equivalent in function to:

<^C2>+6

This expression is evaluated during assembly as the 1's complement of 2, plus the absolute value of 6. When these terms are combined, the resulting expression value generates a carry beyond the most significant bit, leaving 000003(8) as the reduced value.

As shown above, when the temporary numeric control operator and its argument are coded as a term within an expression, angle brackets should be used as delimiters to ensure precise evaluation and readability.

MACRO-11 also supports a unary operator for numeric control which allows you to specify an argument consisting of a 1-word floating-point number. For example, the following statement:

A: MOV #^F3.7,R0

creates a 1-word floating-point number at location A+2 containing the value 3.7 formatted as shown below.

0

BIT 15

14 7 6

S EEEEEEE MMMMMMM

Sign (bit 15) Exponent (bits 14-7) Mantissa (bits 6-0)

The importation below.	ance of ordering with respect to unary operators is shown
	<pre>'F1.0 = 040200 'F-1.0 = 140200 -^F1.0 = 137600 -^F-1.0 = 037600</pre>
The value of term that of	created by the ^F unary operator and its argument is then a can be used by itself or in an expression. For example:
°C^F6.	.2
is equivale	ent to:
C<^F	5.2> and and a state of the sta

For this reason, the use of angle brackets is advised. Expressions used as terms or arguments of a unary operator must be explicitly grouped.

# 6.5 LOCATION COUNTER CONTROL DIRECTIVES

The directives used in controlling the value of the current location counter and in reserving storage space in the object program are described in the following sections.

In this connection, it should be noted that several MACRO-11 statements may cause an odd number of bytes to be allocated, as listed below:

- 1. .BYTE directive
- 2. .BLKB directive
- 3. .ASCII or .ASCIZ directive
- 4. .ODD directive
- A direct assignment statement of the form .=.+expression, which results in the assignment of an odd address value.

In cases that yield an odd address value, the next word-boundaried instruction automatically forces the location counter to an even value, but that instruction is flagged with an error code (B) in the assembly listing.

# 6.5.1 .EVEN Directive

The .EVEN directive ensures that the current location counter contains an even value by adding 1 if the current value is odd. If the current location counter is already even, no action is taken. Any operands following an .EVEN directive are flagged with an error code (Q) in the assembly listing.

The .EVEN directive is used as follows:

.ASCIZ /THIS IS A TEST/ .EVEN ;ENSURES THAT THE NEXT STATEMENT WILL ;BEGIN ON A WORD BOUNDARY. .WORD XYZ

#### 6.5.2 .ODD Directive

The .ODD directive ensures that the current location counter contains an odd value by adding 1 if the current value is even. If the current location counter is already odd, no action is taken. Any operands following an .ODD directive are also flagged with an error code (Q) in the assembly listing.

# 6.5.3 .BLKB and .BLKW Directives

Blocks of storage can be reserved in the object program by means of the .BLKB and .BLKW directives. The .BLKB directive is used to reserve byte blocks; similarly, the .BLKW directive reserves word blocks. The two directives are of the form:

```
.BLKB exp
.BLKW exp
```

where: exp represents the specified number of bytes or words to be reserved in the object program. If no argument is present, a default value of 1 is assumed. These directives should not be used without arguments. Any expression that is completely defined at assembly-time and that reduces to an absolute value is legal. If the expression specified in either of these directives is not an absolute value, the statement is flagged with an error code (A) in the assembly listing.

Figure 6-6 illustrates the use of the .BLKB and .BLKW directives.

166	600000		PSECT	IMPURE, D	
167	000000		PASSII BLKW	1	IPASS FLAG
168			•-•	-	INFYT GROUP MUST STAY TOGETHER
169	000000		,PSECT	IMPPAS, D, GBL	
170	000000		SYMBOL::.BLKW	2	SYMBOL ACCUMULATOR
171	000004		MODESS		IMODE/FLAGS BYTE
172	000004		FLAGSII ,BLKB	1	1
173	000005		SECTOR:: BLKB	1	SYMBOL/EXPRESSION TYPE
174	000006		VALUE:: BLKW	i	JEXPRESSION VALUE
175	000010		RELLVL BLKW	1	IRFI OCATION I EVEL
176		000003	REPT	MAXXMT-<<,-SYMB	0L>/2>
177			BLKW	1	
178			ENDR		
179			• - • • •		
180	000020		CLCNAM:: BLKW	2	CURRENT LOCATION COUNTER NAME
181	000024		CLCFGS::.BLKB	1	<b>)</b>
182	000025		CLCSEC:: BLKB	1	j - Alexandra de la composición de la c
183	000026		CLCLOC:: BLKW	1	
184	000030		CLCMAXEE BLKW	Ĩ	JEND OF GROUPED DATA
185	000032		CHRPNTSS BLKW	Ĩ	ICHARACTER POINTER
186	000034		SYMBEG BLKW	i	POINTER TO START OF SYMBOL
187	000036		ENDFLGII BLKW	ī	
188	000000		PSECT		

Figure 6-6 Example of .BLKB and .BLKW Directives

The .BLKB directive in a source program has the same effect as the following statement:

# .=.+expression

which causes the value of the expression to be added to the current value of the location counter. The .BLKB directive, however, is easier to interpret in the context of the source code in which it appears and is therefore recommended.

# 6.6 TERMINATING DIRECTIVES

## 6.6.1 .END Directive

exp

The .END directive indicates the logical end of the source input, and takes the following form:

.END exp

where:

represents an optional expression value which, if present, indicates the program-entry point, i.e., the transfer address at which program execution is to begin.

When MACRO-ll encounters a valid occurrence of the .END directive, it terminates the current assembly pass. Any additional text beyond this point in the current source file, as well as in additional source files identified in the command line, will be ignored.

When creating an image consisting of several object modules, only one object module may be terminated with an .END exp statement specifying the starting address. All other object modules must be terminated with an .END statement without an address argument; otherwise, a diagnostic message will be issued at link time. If no starting address is specified in any of the object modules, image execution will begin at location 1 of the image and immediately fault because of an odd addressing error.

The .END statement must not be used within a macro expansion or a conditional assembly block; if it is so used, it is flagged with an error code (O) in the assembly listing. The .END statement may be used, however, in an immediate conditional statement (see Section 6.10.2).

If the source program input is not terminated with an .END directive, an error code (E) results in the assembly listing.

#### 6.6.2 .EOT Directive

Under RSX-11, RT-11, and IAS operating systems, the MACRO-11 .EOT directive is ignored and simply treated as a directive without effect, i.e., as a no-op.

#### 6.7 PROGRAM BOUNDARIES DIRECTIVE: .LIMIT

It is often desirable to know the upper and lower address boundaries of the image. When the .LIMIT directive is specified in the source program, MACRO-11 effectively generates the following instruction:

.BLKW

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causing two storage words to be reserved in the object module. Later, at link time, the lowest address in the load image is inserted into the first reserved word, and the address of the first free word following the image is inserted into the second reserved word.

During linking, the size of the image is rounded upward to the nearest 2-word boundary.

For a discussion of memory allocation and mapping, refer to the applicable system manual (see Section 0.3 in the Preface).

# 6.8 PROGRAM SECTIONING DIRECTIVES

The MACRO-11 program sectioning directives are used to declare names for program sections and to establish certain program section attributes essential to the linking processing.

# 6.8.1 .PSECT Directive

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The .PSECT directive allows absolute control over the memory allocation of a program at link time, because any program attributes established through this directive are passed to the linker.

For example, if you are writing programs for a multi-user environment, a program section containing pure code (instructions only) or a program section containing impure code (data and instructions) may be explicitly declared through the .PSECT directive. Furthermore, these program sections may be explicitly declared as read-only code, qualifying them for use as protected, reentrant programs.

In addition, program sections exhibiting the global (GBL) attribute can be explicitly allocated in a program's overlay structure at link time.

The advantages gained through sectioning programs in this manner therefore relate primarily to control of memory allocation, program modularity, and more effective partitioning of memory. Refer to the applicable system manual for a discussion of memory allocation (see Section 0.3 in the Preface).

The .PSECT directive is formatted as follows:

.PSECT name,argl,arg2,...argn

where: name represents the symbolic name of the program section, as described in Table 6-3.

represents any legal separator (comma, tab and/or space).

more of the legal symbolic argl, represent one or .PSECT arguments defined use with the for arg2,... directive, as described in Table 6-3. The slash argn separating each pair of symbolic arguments listed in the table indicates that these optional arguments are mutually exclusive, i.e., one or the other, but not both, may be specified. Multiple arguments must be separated by a legal separating character. Any symbolic argument specified in the .PSECT directive other than those listed in Table 6-3 will cause that statement to be flagged with an error code (A) in the assembly listing.

# Table 6-3 Symbolic Arguments of .PSECT Directive

Argument	Default	Meaning
NAME	Blank	Establishes the program section name, which is specified as one to six Radix-50 characters. If this argument is omitted, a comma must appear in place of the name parameter. The Radix-50 character set is listed in Section A.2 of Appendix A.
RO/RW	RW	Defines which type of access is permitted to the program section:
		RO=Read-Only Access RW=Read/Write Access
		IAS and RSX-llD set hardware protection for RO program sections. RSX-llM and RT-ll do not provide such protection.
I/D	I	Defines the program section as containing either instructions (I) or data (D). These attributes allow the linker to differentiate global symbols that are program entry-point instructions (I) from those that are data values (D).
GBL/LCL	LCL	Defines the scope of the program section, as subsequently interpreted at link time.
		In building single-segment nonoverlaid programs, the GBL/LCL arguments have no meaning, because the total memory allocation for the program will go into the root segment of the image. The GBL/LCL arguments apply only in the case of overlays.
		If an object module contains a local program section, then the storage allocation for that module will occur within the segment in which the module resides. Many modules can reference this same program section, and the memory allocation for each module is either concatenated or overlaid within the segment, depending on the argument of the program section (.PSECT) defining

(continued on next page)

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# Table 6-3 (Cont.) Symbolic Arguments of .PSECT Directive

Argument Default		Meaning		
GBL/LCL (cont'd)	LCL	its allocation requirements (see CON/OVR below). If an object module contains a global program section, the contributions to this program section are collected across segment boundaries, and the allocation of memory for that section will go into the segment nearest the root in which the first contribution to this program section appeared. (The term contribution implies an allocation of memory to the program section.)		
ABS/REL	REL	Defines the relocatability attribute of the program section:		
		ABS=Absolute (non-relocatable). When the ABS argument is specified, the program section is regarded at link time as an absolute module, thus requiring no relocation. The program section is assembled and loaded, starting at absolute virtual address 0.		
		The location of data in absolute program sections must fall within the virtual memory limits of the segment containing the program section; otherwise, an error results at link time. For example, the following code, although valid at during assembly, may generate an error message if virtual location 100000 is outside the segment's virtual address space:		
		.PSECT ALPHA,ABS .=.+100000 .WORD X		
		The above coding assembles properly, but the resulting load address may be outside the respective segment's boundaries. In such cases, the linker recognizes this as an attempt to load data outside the image and responds with an error message.		
		REL=Relocatable. When the REL argument is specified, the linker calculates a relocation bias and adds it to all references to locations within the program section, i.e., all references to the program section must have a relocation bias added to them to make them absolute.		

(Continued on next page)

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	Table 6-3	3 (Cont.)	
Symbolic	Arguments	of .PSECT	Directive

Argument	Default	Meaning
CON/OVR	CON	Defines the allocation requirements of the program section:
		CON=Concatenated. All program section contributions are to be concatenated with other references to this same program section in order to determine the total memory allocation requirement for this program section.
		OVR=Overlaid. All program section contributions are to be overlaid. Thus, the total allocation requirement for the program section is equal to the largest allocation request made by any individual contribution to this program section.

The only argument in the .PSECT directive that is position-dependent is NAME. If it is omitted, a comma must be used in its place. For example, the directive:

.PSECT ,GBL

shows a .PSECT directive with a blank name argument and the GBL argument. Default values (see Table 6-3) are assumed for all other unspecified arguments.

Once the attributes of a program section are declared through a .PSECT directive, MACRO-11 assumes that these attributes remain in effect for all subsequent .PSECT directives of the same name that are encountered within the module.

MACRO-11 provides for 256(10) program sections, as listed below:

- 1. One default absolute program section (. ABS.)
- 2. One default unnamed relocatable program section
- 3. Two-hundred-fifty-four named program sections.

The .PSECT directive enables the user to:

- 1. Create program sections (see Section 6.8.1.1)
- 2. Share code and data among program sections (see Section 6.8.1.2).

For each program section specified or implied, MACRO-11 maintains the following information:

- 1. Program section name
- 2. Contents of the current location counter
- 3. Maximum location counter value encountered
- 4. Program section attributes, i.e., the .PSECT arguments described in Table 6-3 above.

6.8.1.1 Creating Program Sections - MACRO-11 automatically begins assembling source statements at relocatable zero of the unnamed program section, i.e., the first statement of a source program is always an implied .PSECT directive.

The first occurrence of a .PSECT directive with a given name assumes that the current location counter is set at relocatable zero. The scope of this directive then extends until a directive declaring a different program section is specified. Further occurrences of a program section name in subsequent .PSECT statements cause the resumption of assembly where that section previously ended. For example:

	.PSECT		;DECLARES UNNAMED RELOCATABLE PROGRAM
Α:	.WORD	0	;SECTION ASSEMBLED AT RELOCATABLE
В:	.WORD	0	;ADDRESSES 0, 2, AND 4.
C:	.WORD	0	
	.PSECT	ALPHA	;DECLARES RELOCATABLE PROGRAM SECTION
Х:	.WORD	0	;NAMED ALPHA ASSEMBLED AT RELOCATABLE
Υ:	.WORD	0	;ADDRESSES 0 AND 2.
	.PSECT		;RETURNS TO UNNAMED RELOCATABLE
D:	.WORD	0	; PROGRAM SECTION AND CONTINUES ASSEM-
			BLY AT RELOCATABLE ADDRESS 6.

A given program section may be defined completely upon encountering its first .PSECT directive. Thereafter, the section can be referenced by specifying its name only, or by completely respecifying its attributes. For example, a program section can be declared through the directive:

#### .PSECT ALPHA, ABS, OVR

and later referenced through the equivalent directive:

#### .PSECT ALPHA

which requires no arguments.

By maintaining separate location counters for each program section, MACRO-11 allows the user to write statements that are not physically contiguous within the program, but that can be loaded contiguously following assembly, as shown in the following example.

SEC1,REL,RO 0 0 0	;START A RELOCATABLE PROGRAM SECTION ;NAMED SEC1 ASSEMBLED AT RELOCATABLE ;ADDRESSES 0, 2, AND 4.
Α	ASSEMBLE CODE AT RELOCATABLE
В	;ADDRESSES 6 THROUGH 12.
C	
SECA, ABS	;START AN ABSOLUTE PROGRAM SECTION ;NAMED SECA. ASSEMBLE CODE AT
.+2.A	;ABSOLUTE ADDRESSES 0 AND 2.
SEC1	RESUME RELOCATABLE PROGRAM SECTION
A ST	SEC1. ASSEMBLE CODE AT RELOCATABLE ;ADDRESSES 14 AND 16.
	SEC1,REL,RO 0 0 A B C SECA,ABS .+2,A SEC1 A ST

All labels in an absolute program section are absolute; likewise, all labels in a relocatable section are relocatable. The current location counter symbol (.) is also relocatable or absolute when referenced in a relocatable or absolute program section, respectively.

Any labels appearing on a line containing a .PSECT (or .ASECT or .CSECT) directive are assigned the value of the current location counter before the .PSECT (or other) directive takes effect. Thus, if the first statement of a program is:

#### A: .PSECT ALT, REL

the label A is assigned to relocatable address zero of the unnamed (or blank) program section.

It is not known during assembly where relocatable program sections will be loaded, therefore all references between relocatable sections in a single assembly are translated by MACRO-11 to references relative to the base of the referenced section. Thus, MACRO-11 provides the linker with the necessary information to resolve the linkages between various program sections. Such information is not necessary, however, when referencing an absolute program section, because all instructions in an absolute program section are associated with an absolute virtual address.

In the following example, references to the symbols X and Y are translated into references relative to the base of the relocatable program section named SEN.

.=.+10	.PSECT	ENT, ABS	
A:	CLR	X	;ASSEMBLED AS CLR BASE OF ;RELOCATABLE SECTION + 10.
	JMP	Y	;ASSEMBLED AS JMP BASE OF ;RELOCATABLE SECTION + 6.
	.PSECT MOV	SEN,REL R0,R1	
Y:	JMP HALT	А	;ASSEMBLED AS JMP 1000.
X:	WORD	0	

#### NOTE

In the preceding example, using a constant in conjunction with the current location counter symbol (.) in the form .=1000 would result in an error, because constants are always absolute and are always associated with the program's .ASECT (. ABS.). If the form .=1000 were used, a program section incompatibility would be detected. See Section 3.6 for a discussion of the current location counter.

6.8.1.2 Code or Data Sharing - Named relocatable program sections with the arguments GBL and OVR operate in the same manner as FORTRAN COMMON, i.e., program sections of the same name with the arguments GBL and OVR from different assemblies are all loaded at the same location at link time. All other program sections, i.e., those with the argument CON, are concatenated.

Note that no conflict exists between internal symbolic names and program section names, i.e., it is legal to use the same symbolic name for both purposes. Considering FORTRAN again, using the same symbolic name is necessary to accommodate the following statement:

COMMON /X/A,B,C,X

where the symbol X represents the base of the program section and also the fourth element of that section.

6.8.1.3 Memory Allocation Considerations - The assembler does not generate an error when a module ends at an odd location. This allows you to place odd length data at the end of a module. However, when several modules contain object code contributions to the same program section having the concatenate attribute (see Table 6-3), odd length modules (except the last) may cause succeeding modules to be linked starting at odd locations, thereby making the linked program unexecutable. To avoid this problem, code and data should be separated from each other and be placed in separately named program sections. This permits the linker to automatically begin each program section on an even address. Refer to the applicable system manual for further information on memory allocation of tasks (see Section 0.3 in the Preface).

# 6.8.2 .ASECT and .CSECT Directives

IAS and RSX-11 assembly-language programs use the .PSECT and .ASECT directives exclusively, since the .PSECT directive provides all the capabilities of the .CSECT directive defined for other PDP-11 assemblers. MACRO-11 will accept both .ASECT and .CSECT directives, but assembles them as though they were .PSECT directives with the default attributes listed in Table 6-4. Also, compatibility exists between other MACRO-11 programs and the IAS/RSX-11 Task Builders, since the respective Task Builders recognize the .ASECT and .CSECT directives that appear in such programs and likewise assign the default values listed in Table 6-4.

	e se e se de la composición La composición de la	Default Value	
Attribute	ASECT	.CSECT (named)	.CSECT (unnamed)
Name	. ABS.	name	Blank
Access	RW	RW	RW
Туре	Ι	an an <b>T</b> aran an Araba an Araba. An Taran an Araba	I and a second sec
Scope	GBL	GBL	LCL
Relocation	ABS	REL	REL
Allocation	OVR	OVR	CON

Table 6-4 Non-IAS/RSX-ll Program Section Default Values

The allowable syntactical forms of the .ASECT and .CSECT directives are:

.ASECT .CSECT .CSECT symbol

Note that the statement:

.CSECT JIM

is identical to the statement:

.PSECT JIM,GBL,OVR

because the .CSECT default values GBL and OVR are assumed for the named program section.

# 6.9 SYMBOL CONTROL DIRECTIVE: .GLOBL

MACRO-11 produces a relocatable object module and a listing file containing the assembly listing and symbol table. The linker joins separately-assembled object modules into a single executable image. During linking, object modules are relocated as a program function of the specified base of the module. The object modules are then linked via global symbols, such that a global symbol in one module, defined either by a global assignment operator (==), a global label operator (::), or the .GLOBL directive can be referenced from another module. Thus, all symbols which will be referenced by other program modules must be singled out as global symbols in the defining modules.

The .GLOBL directive is provided to define (and thus provide linkage to) symbols not otherwise defined as global symbols within a module. For example, if the .DSABL GBL directive is in effect (see Section 6.2), .GLOBL directives might be included in a source program to effect linkage to library routines. For a global symbol definition, the directive .GLOBL A,B,C is equivalent to:

> A==expression (or A::) B==expression (or B::) C==expression (or C::)

Thus, the general form of the .GLOBL directive is:

.GLOBL syml, sym2,...symn

where: syml, represent legal symbolic names. When multiple sym2,... symbols are specified, they are separated by any symn legal separator (comma, space, and/or tab).

A .GLOBL directive may also embody a label field and/or a comment field.

3

At the end of assembly pass 1, MACRO-11 determines whether a given global symbol is defined within the current program module or whether it is to be treated as an external symbol. All internal symbols appearing within a given program must be defined at the end of assembly pass 1 or they will be assumed to be default global references. Refer to Section 6.2 for a description of enabling/disabling of global references.

In the example below, A and B are entry-point symbols. The symbol A has been explicitly defined as a global symbol by means of the .GLOBL directive, and the symbol B has been explicitly defined as a global label by means of the double colon (::). Since the symbol C is not defined as a label within the current assembly, it is an external (global) reference if .ENABL GBL is in effect.

•			
; DE	FINE A SUBR	OUTINE WITH	2 ENTRY POINTS WHICH CALLS AN
; EX	TERNAL SUBR	OUTINE	
;			
	.PSECT		;DECLARE THE UNNAMED PROGRAM SECTION.
	.GLOBL	Α	;DEFINE A AS A GLOBAL SYMBOL.
Α:	MOV	@(R5)+,R0	;DEFINE ENTRY POINT A.
	MOV	#X,Rl	
X:	JSR	PC,C	;CALL EXTERNAL SUBROUTINE C.
	RTS	R5	;EXIT.
B::	MOV	(R5)+,R1	;DEFINE ENTRY POINT B.
	CLR	R2	
	BR	Х	

External symbols can appear in the operand field of an instruction or MACRO-11 directive as a direct reference, as shown in the examples below:

CLR EXT .WORD EXT CLR @EXT

External symbols may also appear as a term within an expression, as shown below:

CLR EXT+A .WORD EXT-2 CLR @EXT+A(R1)

It should be noted that an undefined external symbol cannot be used in the evaluation of a direct assignment statement or as an argument in a conditional assembly directive (see Sections 6.10.1 and 6.10.3).

#### 6.10 CONDITIONAL ASSEMBLY DIRECTIVES

Conditional assembly directives allow you to include or exclude blocks of source code during the assembly process, based on the evaluation of stated condition tests within the body of the program. This capability allows several variations of a program to be generated from the same source code.

# 6.10.1 Conditional Assembly Block Directives: .IF, .ENDC

The general form of a conditional assembly block is as follows:

cond, argument(s) ;START CONDITIONAL ASSEMBLY BLOCK. .IF ;RANGE OF CONDITIONAL ASSEMBLY BLOCK. range ;END OF CONDITIONAL ASSEMBLY BLOCK. .ENDC represents a specified condition that must be met where: cond if the block is to be included in the assembly. by the The conditions that may be tested conditional assembly directives are defined in Table 6-5. represents any legal separator (comma, space, and/or tab). symbolic argument(s) or represent(s) the argument(s) expression(s) of the specified conditional test. These arguments are thus a function of the specified condition to be tested (see Table 6-5). represents the body of code that is either range included in the assembly or excluded, depending upon whether the specified condition is met. terminates the conditional assembly block. This .ENDC directive must be present to end the conditional assembly block.

A condition test other than those listed in Table 6-5, an illegal argument, or a null argument specified in an .IF directive causes that line to be flagged with an error code (A) in the assembly listing.

Table 6-5						
Legal	Condition	Tests	for	Conditional	Assembly	Directives

Cond	itions	and the second sec	n an thair an air an
Positive	Complement	Arguments	Assemble Block If:
EQ	NE	Expression	Expression is equal to 0 (or not equal to 0).
GT	LE	Expression	Expression is greater than 0 (or less than or equal to 0).

		Τa	able	6-5 (Cont.)		
Legal	Condition	Tests	for	Conditional	Assembly	Directives

Conditions		en e	
Positive	Complement	Arguments	Assemble Block If:
LT	GE	Expression	Expression is less than 0 (or greater than or equal to 0).
DF	NDF	Symbolic argument	Symbol is defined (or not defined).
В	NB	Macro-type argument	Argument is blank (or non-blank).
IDN	DIF	Two macro-type arguments	Arguments are identical (or different).
Z	NZ	Expression	Same as EQ/NE.
G	L	Expression	Same as GT/LT.

#### NOTE

A macro-type argument (which is a form of symbolic argument), as shown below, is enclosed within angle brackets or denoted with an up-arrow construction (as described in Section 7.3.1).

# <A,B,C> ^/124/

An example of a conditional assembly directive follows:

	IF E	EQ	ALPHA+1	;ASSEMBLE	BLOCK	IF	ALPHA+1=	:(	)		
--	------	----	---------	-----------	-------	----	----------	----	---	--	--

.ENDC

&

1

The two operators & and ! have special meaning within DF and NDF conditions, in that they are allowed in grouping symbolic arguments.

Logical AND operator

Logical inclusive OR operator

For example, the conditional assembly statement:

.IF DF SYM1 & SYM2

.ENDC

results in the assembly of the conditional block if the symbols SYML and SYM2 are both defined.

6-42

Nested conditional directives take the form:

Conditional Assembly Directive Conditional Assembly Directive

. ENDC

For example, the following conditional directives:

.IF DF SYM1 .IF DF SYM2 . . . ENDC .ENDC

can govern whether assembly is to occur. In the example above, if the outermost condition is unsatisfied, no deeper level of evaluation of nested conditional statements within the program occurs.

Each conditional assembly block must be terminated with an .ENDC directive. An .ENDC directive encountered outside a conditional assembly block is flagged with an error code (O) in the assembly listing.

MACRO-ll permits a nesting depth of 16(10) conditional assembly levels. Any statement that attempts to exceed this nesting level depth is flagged with an error code (0) in the assembly listing.

# 6.10.2 Subconditional Assembly Block Directives: .IFF, .IFT, .IFTF

Subconditional directives may be placed within conditional assembly blocks to indicate:

- 1. The assembly of an alternate body of code when the condition of the block tests false.
- 2. The assembly of a non-contiguous body of code within the conditional assembly block, depending upon the result of the conditional test in entering the block.
- 3. The unconditional assembly of a body of code within a conditional assembly block.

The subconditional directives are described in detail in Table 6-6. If a subconditional directive appears outside a conditional assembly block, an error code (O) is generated in the assembly listing.

# Table 6-6Subconditional Assembly Block Directives

Subconditional Directive	Function
.IFF	If the condition tested upon entering the conditional assembly block is false, the code following this directive, and continuing up to the next occurrence of a subconditional directive or to the end of the conditional assembly block, is to be included in the program.
.IFT	If the condition tested upon entering the conditional assembly block is true, the code following this directive, and continuing up to the next occurrence of a subconditional directive or to the end of the conditional assembly block, is to be included in the program.
.IFTF	The code following this directive, and continuing up to the next occurrence of a subconditional directive or to the end of the conditional assembly block, is to be included in the program, regardless of the result of the condition tested upon entering the conditional assembly block.

The implied argument of a subconditional directive is the condition test specified upon entering the conditional assembly block, as reflected by the initial directive in the conditional coding examples below. Conditional or subconditional directives in nested conditional assembly blocks are not evaluated if the previous (or outer) condition in the block is not satisfied. Examples 3 and 4 below illustrate nested directives that are not evaluated because of previous unsatisfied conditional coding.

EXAMPLE 1: Assume that symbol SYM is defined.

.IF DF SYM	;TESTS TRUE, SYM IS DEFINED. ASSEMBLE ;THE FOLLOWING CODE.
.IFF	;TESTS FALSE. SYM IS DEFINED. DO NOT ;ASSEMBLE THE FOLLOWING CODE.
.IFT	;TESTS TRUE. SYM IS DEFINED. ASSEM- ;BLE THE FOLLOWING CODE.
.IFTF	;ASSEMBLE FOLLOWING CODE UNCONDITION- ;ALLY.
.IFT 	;TESTS TRUE. SYM IS DEFINED. ASSEM- ;BLE REMAINDER OF CONDITIONAL ASSEM- ;BLY BLOCK.

6-44

EXAMPLE	2:	Assu defi	me th ned.	at	symbol	Xi	s	def	ine	đ	and	1	that	t sy	ymbol	Y	is	n	ot
	.IF .IF .IFF	DF DF	X Y			; T ; T ; T ; A	'ES 'ES 'SS	TS TS TS EMB	TRU FAL TRU LE	E, SE E, TH	SYI SYI E F(	MBC YME MBC OLI	DL 3 BOL DL 3 LOW]	K IS Y I Y IS ING	5 DEF 5 NO 5 NOT CODE	INE T C DE	D. EF	INI NEI	ED. D,
	.IFI					;1 ;D	'ES OO	TS NOI	FAL AS	SE SEI	, S MBL	YME E 1	BOL PHE	Y I FOI	S NO LLOWI	T E NG	EF CO	INI DE	ED.
	. ENC . ENC	DC DC																	
EXAMPLE	3:	Assu defi	ned.	at	symbol	Ai	s	def	ine	d i	and	l tł	nat	syn	nbol	Ві	s	no	t
	.IF	DF	A			;1 ;A	res Ass	TS EME	TRU	E. TH	A E F	I S	S DI	EFIN ING	NED. CODE	•			
	MOV •		A,Rl																
	•																		
	.IFE	P				;1 ;A	TES ASS	TS Eme	FAL BLE	SE	E F	A :	IS I LOW:	DEFI ING	INED. CODE		00	NO	r
	MOV		R1,R0																
	.IF	NDF	В			; N ; E	NES EVA	TEC LUA	) CC	)ND	ITI	ON	AL I	DIRI	ECTIV	Έ	ſS	NO	Т
	. ENI	DC																	

.ENDC

EXAMPLE 4: Assume that symbol X is not defined and that symbol Y is defined.

.IF DF	X	;TESTS FALSE. SYMBOL X IS NOT DEFINED ;DO NOT ASSEMBLE THE FOLLOWING CODE.
.IF DF	Y	;NESTED CONDITIONAL DIRECTIVE IS NOT ;EVALUATED.
era 🔸 raa		
.IFF		;NESTED SUBCONDITIONAL DIRECTIVE IS ;NOT EVALUATED.
.IFT		;NESTED SUBCONDITIONAL DIRECTIVE IS ;NOT EVALUATED.
•		
ENDC		

# 6.10.3 Immediate Conditional Assembly Directive: .IIF

An immediate conditional assembly directive provides a means for writing a 1-line conditional assembly block. In using this directive, no terminating .ENDC statement is required, and the condition to be tested is completely expressed within the line containing the directive. Immediate conditional assembly directives are of the form:

.IIF cond, arg, statement

where: cond represents one of the legal condition tests defined for conditional assembly blocks in Table 6-5.

represents any legal separator (comma, space, and/or tab).

arg represents the argument associated with the immediate conditional directive, i.e., an expression, symbolic argument, or macro-type argument, as described in Table 6-5.

represents the separator between the conditional argument and the statement field. If the preceding argument is an expression, then a comma must be used; otherwise, a comma, space, and/or tab may be used.

statement represents the specified statement to be assembled if the condition is satisfied.

For example, the immediate conditional statement:

.IIF DF FOO, BEQ ALPHA

generates the code

BEQ ALPHA

if the symbol FOO is defined within the source program.

As with the .IF directive, a condition test other than those listed in Table 6-5, an illegal argument, or a null argument specified in an .IIF directive results in an error code (A) in the assembly listing.

# 6.10.4 PAL-11R Conditional Assembly Directives

In order to maintain compatibility with programs developed under PAL-11R, the following conditionals remain permissible under MACRO-11. It is advisable, however, to develop future programs using the format for MACRO-11 conditional assembly directives.

Directive	Arguments	Assemble Block if				
.IFZ or .IFEQ .IFNZ or .IFNE .IFL or .IFLT .IFG or .IFGT .IFLE .IFDF .IFDF	expression expression expression expression symbolic argument	expression=0 expression not equal 0 expression<0 expression>0 expression is < or =0 symbol is defined symbol is updofined				
	by moorre argumente	DIMPOT TO ANACITICA				

The rules governing these directives are the same as for the MACRO-11 conditional assembly directives previously described.

# CHAPTER 7

## MACRO DIRECTIVES

# 7.1 DEFINING MACROS

In assembly-language programming, it is often convenient and desirable to generate a recurring coding sequence by invoking a single statement within the program. In order to do this, the desired coding sequence is first established with dummy arguments as a macro definition. Once a macro has been defined, a single statement calling the macro by name with a list of real arguments (replacing the corresponding dummy arguments in the macro definition) generates the desired coding sequence. This sequence is called the macro expansion.

#### 7.1.1 .MACRO Directive

label

The first statement of a macro definition must be a .MACRO directive. This directive takes the form:

label: .MACRO name, dummy argument list

where:

represents the programmer-assigned symbolic name name of the macro. This name may be any legal symbol and may be used as a label elsewhere in the program.

represents an optional statement label.

represents any legal separator (comma, space, and/or tab).

represents a number of legal symbols (see 3.2.2) dummy that may appear anywhere in the body of the macro argument definition, even as a label. These dummy symbols list can be used elsewhere in the program with no conflict of definition. Multiple dummy arguments specified in this directive may be separated by any legal separator. The detection of a duplicate an illegal symbol in a dummy argument list or terminates the scan and causes an error code to be generated.

A comment may follow the dummy argument list in a .MACRO directive, as shown below:

> .MACRO ABS A,B ; DEFINES MACRO ABS WITH TWO ARGUMENTS.

# NOTE

Although it is legal for a label to appear on a .MACRO directive, this practice is discouraged, especially in the case of nested macro definitions, because invalid labels or labels constructed with the concatenation character will cause the macro directive to be ignored. This may result in improper termination of the macro definition. This NOTE also applied to .IRP, .IRPC, and .REPT.

# 7.1.2 .ENDM Directive

The final statement of every macro definition must be an .ENDM directive of the form:

.ENDM name

where: name represents an optional argument specifying the symbolic name of the macro being terminated by the directive, as shown in the following example:

.ENDM

# ;TERMINATES THE CURRENT ;MACRO DEFINITION.

.ENDM ABS ;TERMINATES THE CURRENT ;MACRO DEFINITION NAMED ABS.

If specified, the symbolic name in the .ENDM statement must match the name specified in the corresponding .MACRO directive. Otherwise, the statement is flagged with an error code (A) in the assembly listing (see Appendix D). In either case, the current macro definition is terminated. Specifying the macro name in the .ENDM statement thus permits MACRO-11 to detect missing .ENDM statements or improperly-nested macro definitions.

The .ENDM directive may be followed by a comment field, but must not contain a label, as shown below:

.MACRO TYPMSG MESSGE ;TYPE A MESSAGE. JSR R5,TYPMSG .WORD MESSGE .ENDM ;END OF TYPMSG MACRO.

An .ENDM statement encountered by MACRO-11 outside a macro definition is flagged with an error code (O) in the assembly listing (see Appendix D).

#### NOTES

1. Labels on .ENDM directives are ignored.

2. Illegal labels will cause the directive to be bypassed.

# 7.1.3 .MEXIT Directive

The .MEXIT directive may be used to terminate a macro expansion before the end of the macro is encountered. This directive is also legal within repeat blocks (see Sections 7.6 and 7.7). It is most useful in the context of nested macros. The .MEXIT directive terminates the current macro as though an .ENDM directive had been encountered. Using the .MEXIT directive bypasses the complexities of nested conditional directives and alternate assembly paths, as shown in the following example:

.ENDM

## ;NORMAL END OF MACRO.

Considering the above macro, in an assembly where the real argument for the dummy symbol N is equal to zero (see Table 6-5), the conditional block would be assembled, and the macro expansion would be terminated by the .MEXIT directive. When macros are nested, a .MEXIT directive causes an exit to the next higher level of macro expansion.

A .MEXIT directive encountered outside a macro definition is flagged with an error code (O) in the assembly listing.

#### 7.1.4 MACRO Definition Formatting

A form-feed character used within a macro definition causes a page eject during the assembly of the macro definition. A page eject, however, is not performed when the macro is expanded.

Conversely, when the .PAGE directive is specified within a macro definition, it is ignored during the assembly of the macro definition, but a page eject is performed when that macro is expanded.

#### 7.2 CALLING MACROS

A macro definition must be established by means of the .MACRO directive (see Section 7.1.1) before the macro can be expanded within the source program. Macro calls are of the general form:

label: name real arguments

where: label represents an optional statement label.

name represents the name of the macro, as specified in the .MACRO directive (see Section 7.1.1).

real represent symbolic arguments which replace the dummy arguments specified in the .MACRO arguments directive. When multiple arguments are specified, they are separated by any legal separator. Arguments to the macro call are treated as character strings whose usage is determined by the macro definition. Note that MACRO-11 accepts the ASCII value of lower-case alphabetic characters when .ENABL LC has been specified.

When a macro name is the same as a user label, the appearance of the symbol in the operator field designates the symbol as a macro call; the appearance of the symbol in the operand field designates it as a label, as shown below:

ABS:	MOV	(R0),R1	;ABS IS DEFINED AS A LABEL.	
	•			
	•			
	BR	ABS	;ABS IS CONSIDERED TO BE A LABEL.	
	•			
	•			
	ABS	#4,ENT,LAR	;ABS IS A MACRO CALL.	

# 7.3 ARGUMENTS IN MACRO DEFINITIONS AND MACRO CALLS

Arguments within a macro definition or macro call are separated from other arguments by any of the legal separating characters described in Section 3.1.1.

Macro definition arguments (dummy) and macro call arguments (real) normally maintain a strict positional relationship. That is, the first real argument in a macro call corresponds with the first dummy argument in a macro definition. Only the use of keyword arguments in a macro call can override this correspondence (see Section 7.3.6).

For example, the following macro definition and its associated macro expansion contain multiple arguments:

.MACRO REN A, B, C

.

REN ALPHA, BETA, <C1, C2>

Arguments which themselves contain separating characters must be enclosed in paired angle brackets, as shown above. For example, the macro call:

REN <MOV X,Y>,#44,WEV

causes the entire expression

MOV X,Y

to replace all occurrences of the symbol A in the macro definition. Real arguments within a macro call are considered to be character strings and are treated as a single entity during the macro expansion. The up-arrow (<sup>^</sup>) construction is provided to allow angle brackets to be passed as part of the argument. This construction, for example, could have been used in the above macro call, as follows:

REN ^/<MOV X,Y>/,#44,WEV

causing the entire character string <MOV X,Y> to be passed as an argument.

The following macro call:

REN #44,WEV^/MOV X,Y/

however, contains only two arguments (#44 and WEV^/MOV X,Y/), because the up-arrow is a unary operator (see Section 3.1.3) and it is not preceded by an argument separator.

As shown in the examples above, spaces can be used within bracketed argument constructions to increase the legibility of such expressions.

# 7.3.1 Macro Nesting

The nesting of macros, where the expansion of one macro includes a call to another, causes one set of angle brackets in the macro definition to be removed from an argument with each nested call. The depth of nesting allowed is dependent upon the amount of dynamic memory used by the source program being assembled.

To pass an argument containing legal argument delimiters to nested macros, the argument in the macro definition should be enclosed within one set of angle brackets for each level of nesting, as shown in the coding sequence below. It should be noted that this extra set of angle brackets for each level of nesting is required in the macro definition, not in the macro call.

.MACRO	LEVEL1 DI <dum1></dum1>	UM1,DUM2
LEVELZ	(DOM2)	
• • • • • • • • •		
.MACRO	LEVEL2 D	UM 3
DUM 3		
ADD	#10,RU	
MOV	R0, (R1) +	
.ENDM		

A call to the LEVEL1 macro, as shown below, for example:

LEVEL1 <MOV X,R0>,<MOV R2,R0>

causes the following macro expansion to occur:

VOM	X,RO
ADD	#10,R0
MOV	R0, (R1) +
MOV	R2,R0
ADD	#10,R0
MOV	R0, (R1) +

When macro definitions are nested, i.e., when a macro definition is contained entirely within the definition of another macro, the inner definition is not a callable macro until the outer macro has been called and expanded. For example, in the following coding:

```
.MACRO LV1 A,B
...
.MACRO LV2 C
...
.ENDM
.ENDM
```

the LV2 macro cannot be called and expanded until the LV1 macro has been so invoked. Likewise, any macro defined within the LV2 macro definition cannot be called and expanded until LV2 has also been invoked.

# 7.3.2 Special Characters in Macro Arguments

An argument may include special characters without enclosing them in a bracketed construction if that argument does not contain spaces, tabs, semicolons, or commas. For example, the macro definition:

```
.MACRO PUSH ARG
MOV ARG,-(SP)
.ENDM
.
PUSH X+3(%2)
```

causes the following code to be generated:

MOV X+3(%2),-(SP)

# 7.3.3 Passing Numeric Arguments as Symbols

When macro arguments are passed, an absolute symbol value can be passed which is treated by the macro as a numeric string. An argument preceded by the unary operator backslash (\) is treated as a numeric value in the current program radix. The ASCII characters representing this value are inserted in the macro expansion, and their function is defined in the context of the resulting code, as shown in the following example:

	.MACRO CON	INC A,B A,\B	B IS TREATED AS A NUMBER IN CURRENT
B=B+1	.ENDM		PROGRAM RADIX.
A'B:	.MACRO .WORD .ENDM	CON A,B 4	;A'B IS DESCRIBED IN SECTION 7.3.6.
	•		
C=0	INC	X,C	

The above macro call (INC) would thus expand to:

X0: .WORD 4

Note in this expanded code that the label X0: is the result of the concatenation of two real arguments. The single quote (') character in the label A'B: causes the real arguments X and 0 to be concatenated as they are passed during the expansion of the macro. This type of argument construction is described in further detail in Section 7.3.6.

A subsequent call to the same macro would generate the following code:

X1: .WORD 4

and so on, for later calls. The two macro definitions are necessary because the symbol associated with dummy argument B (i.e., C) cannot be updated in the CON macro definition, because its numeric value has already been substituted for its symbolic name, i.e., the character 0 has replaced C in the argument string. In the CON macro definition, the number passed is treated as a string argument. (Where the value of the real argument is 0, only a single 0 character is passed to the macro expansion.)

Passing numeric values in this manner is useful in identifying source listings. For example, versions of programs created through conditional assemblies of a single source program can be identified through such coding as that shown below. Assume, for example, that the symbol ID in the macro call (IDT) has been equated elsewhere in the source program to the value 6.

.MACRO	IDT SYM	;ASSUME THAT THE SYMBOL ID TAKES
.IDENT	/V05A'SYM/	;ON A UNIQUE 2-DIGIT VALUE.
.ENDM		;WHERE V05A IS THE UPDATE
•		; VERSION OF THE PROGRAM.
•		

IDT \ID

The above macro call would then expand to:

.IDENT /V05A6/

where 6 is the numeric value of the symbol ID.

# 7.3.4 Number of Arguments in Macro Calls

If more arguments appear in the macro call than in the macro definition, an error code (Q) is generated in the assembly listing. If fewer arguments appear in the macro call than in the macro definition, missing arguments are assumed to be null values. The conditional directives .IF B and .IF NB (see Table 6-5) can be used within the macro to detect missing arguments. The number of arguments can also be specified using the .NARG directive (Section 7.4.1). Note that a macro can be defined with no arguments.

# 7.3.5 Creating Local Symbols Automatically

A label is often required in an expanded macro. In the conventional macro facilities thus far described, such a label must be explicitly

# MACRO DIRECTIVES

specified as an argument with each macro call. Be careful in issuing subsequent calls to the same macro, to avoid specifying a duplicate label as a real argument. This concern can be eliminated through a feature of MACRO-11 which creates a unique symbol where a label is required in an expanded macro.

As noted in Section 3.5, MACRO-11 can automatically create local symbols of the form n\$, where n is a decimal integer within the range 64 through 127, inclusive. Such local symbols are created by MACRO-11 in numerical order, as shown below:

64\$ 65\$ . 126\$ 127\$

This automatic facility is invoked on each call of a macro whose definition contains a dummy argument preceded by the question mark (?) character, as shown in the macro definition below:

.MACRO	ALPHA, A,?B	;CONTAINS;QUESTION	DUMMY MARK.	ARGUMENT	В	PRECEDED BY	
TST	A						
BEQ	B						
ADD	#5,A						

.ENDM

A local symbol is generated automatically by MACRO-11 only when a real argument of the macro call is either null or missing, as shown in Example 1 below, which reflects the expansion of the ALPHA macro defined above.

If the real argument is specified in the macro call, however, MACRO-11 inhibits the generation of a local symbol and normal argument replacement occurs, as shown in Example 2 below.

EXAMPLE 1: Generate a Local Symbol for the Missing Argument:

ALPHA	Rl	;SECOND ARGUMENT IS MISSING.
TST	Rl	
BEQ	64\$	;LOCAL SYMBOL IS GENERATED.
ADD	#5,R1	

64\$:

B:

EXAMPLE 2: Do Not Generate a Local Symbol:

ALPHA	R2,XYZ	;SECOND	ARGUMENT	XYZ	IS	SPECI	FIED.	
BEQ	XYZ	;NORMAL	ARGUMENT	REPI	LACE	SMENT	OCCURS	•
ADD	#5,R2							

XYZ:

Automatically-generated local symbols are restricted to the first 16(10) arguments of a macro definition.

Note that automatically-created local symbols resulting from the expansion of a macro, as described above, do not in any way influence local symbol block boundaries. In other words, such automatically-created local symbols do not establish a local symbol block in their own right.

# MACRO DIRECTIVES

However, when a macro has several arguments earmarked for automatic local symbol generation, substituting a specific label for one such argument introduces a risk that assembly errors will result. This is because MACRO-11 constructs its argument substitution list at the point of macro invocation. Therefore, the appearance of any label, the .ENABL LSB directive, or the .PSECT directive, in the macro expansion will create a new local symbol block. This could leave local symbol references in the previous block and the symbol definitions in the new one, resulting in error codes in the assembly listing (see Appendix D). Furthermore, a subsequent macro expansion that generates local symbols in the new block may duplicate one of the symbols in question, resulting in an additional error code (P) in the assembly listing.

# 7.3.6 Keyword Arguments

Macros may be defined with and/or invoked with keyword arguments. Α keyword argument has the following form:

name=string

where

name represents the dummy argument,

string represents the real symbolic argument.

The keyword argument may not contain embedded argument separators unless properly delimited as described in section 7.3.

When a keyword argument appears in the dummy argument list of a macro definition, the specified string becomes the default real argument at macro call.

When a keyword argument appears in the real argument list of a macro call, the specified string becomes the real argument for the dummy argument that exactly matches the specified name, whether or not the dummy argument was defined with a keyword. If a match fails, the entire argument specification is treated as the next positional real A keyword argument may be specified anywhere in the dummy argument. argument list of a macro definition and is part of the positional ordering of argument. On the other hand, a keyword argument may be specified anywhere in the real argument list of a macro call but does not affect the positional correspondence of the remaining arguments.

1	.LIST ME
2 3 4	; ; DEFINE A MACRO HAVING KEYWORDS IN DUMMY ARGUMENT LIST ;
5 6 7	.MACRO TEST CONTRL=1,BLOCK,ADDRES=TEMP .WORD CONTRL
8 9 10	.WORD BLOCK .WORD ADDRES .ENDM
11 12	
13	; ; NOW INVOKE SEVERAL TIMES
16	<ul> <li>If the second secon second second sec</li></ul>

17	000000		TEST	A,B,C
	000000	000000G	.WORD	
	000002	000000G	.WORD	<b>B</b> shakar a shekar a s
	000004	000000G	.WORD	C A C A C A C A C A C A C A C A C A C A
18			1, 1 10 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	
19	000006		TEST	ADDRES=20, BLOCK=30, CONTRL=40
	000006	000040	.WORD	40
	000010	000030	.WORD	30
	000012	000020	.WORD	20
20				
21	000014		TEST	BLOCK=5
	000014	000001	.WORD	1
	000016	000005	.WORD	5
	000020	000000G	.WORD	TEMP
22				
23	000022		TEST	CONTRL=5,ADDRES=VARIAB
	000022	000005	.WORD	5
	000024	000000	.WORD	
	000026	000000G	.WORD	VARIAB
24				
25	000030		TEST	
	000030	000001	.WORD	1
	000032	000000	.WORD	
	000034	000000G	.WORD	TEMP
26				
27	000036		TEST	ADDRES=JACK!JILL
	000036	000001	.WORD	1
	000040	000000	.WORD	
	000042	00000C	.WORD	JACK!JILL
28				
29				
30		000001	. END	

# 7.3.7 Concatenation of Macro Arguments

The apostrophe or single quote character (') operates as a legal delimiting character in macro definitions. A single quote that precedes and/or follows a dummy argument in a macro definition is removed, and the substitution of the real argument occurs at that point. For example, in the following statements:

	.MACRO	DEF A,B,C
A'B:	.ASCIZ	/C/
	.BYTE	''A,''B
	.ENDM	

when the macro DEF is called through the statement:

DEF X,Y,<MACRO-11>

it is expanded, as follows:

XY: .ASCIZ /MACRO-11/ .BYTE 'X,'Y

In expanding the first line, the scan for the first argument terminates upon finding the first 'character. Since A is a dummy argument, the ' is removed. The scan then resumes with B; B is also noted as another dummy argument. The two real arguments X and Y are then concatenated to form the label XY:. The third dummy argument is noted in the operand field of the .ASCIZ directive, causing the real argument MACRO-11 to be substituted in this field.
When evaluating the arguments to the .BYTE directive during expansion of the second line, the scan begins with the first ' character. Since it is neither preceded nor followed by a dummy argument, this ' character remains in the macro expansion. The scan then encounters the second ' character, which is followed by a dummy argument and is therefore discarded. The scan of argument A is terminated upon encountering the comma (,). The third ' character is neither preceded nor followed by a dummy argument and again remains in the macro expansion. The fourth (and last) ' character is followed by another dummy argument and is likewise discarded. (Note that four ' characters were necessary in the macro definition to generate two ' characters in the macro expansion.)

#### 7.4 MACRO ATTRIBUTE DIRECTIVES: .NARG, .NCHR, AND .NTYPE

Three directives are available in MACRO-11 which allow the user to determine certain attributes of macro arguments. The use of these directives permits selective modifications of a macro expansion, depending on the nature of the arguments being passed. These directives are described separately below.

### 7.4.1 .NARG Directive

The .NARG directive is used to determine the number of arguments in the macro call currently being expanded. Hence, the .NARG directive can appear only within a macro definition; if it does not, an error code (O) is generated in the assembly listing. This directive takes the form:

label: .NARG symbol

where:

label represents an optional statement label.

symbol represents any legal symbol. This symbol is equated to the number of arguments in the macro call currently being expanded. If a symbol is not specified, the .NARG directive is flagged with an error code (A) in the assembly listing.

An	examp	ple of t	he .NARG	directiv	ve foll	OWS:
	1				.TITLE	NARG
	2				MACRO	NOPP, NUM
	<u> </u>				NARG	SYM
	5				TF	FO.SYM
	6				MEXIT	
	7				TFF	
	8				REPT	NUM
	9				NOP	
	10				ENDM	
	11				ENDC	
	12				ENDM	
	13				••••	
	14					
	15	000000			NOPP	
	• •		000000		NARG	SYM
					TF	EQ. SYM
					MEXIT	
					TEE	
					REPT	
					NOP	
					FNDM	
					ENDO	
	16				• • • •	
	17					
	18	000000			NOPP	6
	• •	e de la della de	000001		NARG	SYM
					TF	EQ. SYM
					MEXTT	
					TFF	
			000006		REPT	6
					NOP	
					FNDM	
		aaaaaa	000243		NOP	
		0000002	000240		NOP	
		000004	000240		NOP	
		RARAAL	000240		NOP	
		000010	999240		NOP	
		000012	000240		NOP	
		and a second second second			ENDC	
	19					
	20					
	21		020001		END	

### 7.4.2 .NCHR Directive

The .NCHR directive, which can appear anywhere in a MACRO-ll program, is used to determine the number of characters in a specified character string. This directive, which is useful in calculating the length of macro arguments, takes the following form:

label: .NCHR symbol, <string>

where: label represents an optional statement label.

symbol represents any legal symbol. This symbol is equated to the number of characters in the specified character string. If a symbol is not specified, the .NCHR directive is flagged with an error code (A) in the assembly listing (see Appendix D). represents any legal separator (comma, space, and/or tab).

<string> represents a string of printable characters. The character string need be enclosed within angle brackets (<>) or up-arrows (^) only if the specified character string contains a legal separator (comma, space, and/or tab). If the delimiting characters do not match or if the ending delimiter cannot be detected because of a syntactical error in the character string (thus prematurely terminating its evaluation), the .NCHR directive is flagged with an error code (A) in the assembly listing.

An example of the .NCHR directive follows:

.

1				.TITLE	NCHR
3				. MACRO	CHAR, MESS
4				NCHR	SYM, MESS
5				WORD	SYM
6				ASCII	/MESS/
7				FVEN	
e 9				ENDM	
10					
ii:	ABBBBB		MSG1:	CHAR	SHELLO>
		000005		NCHR	SYM. HELLO
	000000	000005		WORD	SYM
	000002	110		ASCTT	AHELL DA
	000003	105			C. C. C. C. S.
	000004	114			
	000005	114			
	000006	117			
				EVEN	
12					
14		000001		FND	

#### 7.4.3 .NTYPE Directive

The .NTYPE directive is used to determine the addressing mode of a specified macro argument. Hence, the .NTYPE directive can appear only within a macro definition; if it appears elsewhere, it is flagged with an error code (O) in the assembly listing. This directive takes the form:

.NTYPE symbol, aexp label:

where: label

19

symbol

represents an optional statement label.

represents any legal symbol. This symbol is equated to the 6-bit addressing mode of the following argument. If a symbol is not specified, the .NTYPE directive is flagged with an error code (A) in the assembly listing.

represents any legal separator (comma, space, and/or tab).

aexp

represents any legal address expression, as used with an opcode. If no argument is specified, the result will be zero.

An example of the use of an .NTYPE directive in a macro definition is shown below:

1			.TITLE	NTYPE	
3			MACRO	SAVE ARG	
4			NTYPE	SYM. APC	
5			TE	EO. SYMETA	
			MOV		DECISION NODE
ž			···() V	ARG; -(3P)	THEGISTER MODE
á			• J F F		
Š			MOV	#ARG,=(SP)	INON-REGISTER MODE
			. ENDC		
10			. ENDM		
11					
12					
13	000000	AAAAAA TEMP	.WORD	P	
14					
15					
16	000002		SAVE	R1	
		000001	NTYPE	SYM, R1	
			TF	FQ. SYMETA	
	000002	010146	MOV	R1.=(SP)	.PECISTED NONE
		•••	TEF		TACOLOTER HOUE
			MOV	*P1 -(50)	NON-DECISIED MODE
			ENDO	#h1] = (ar)	INDUMERECISIER MUDE
17			• E MIL		
10					
10			<sup>-</sup>		
14	NNNN94		SAVE	TEMP	
		000007	•NTYPE	SYM, TEMP	
			.IF	EQ, SYMR70	
			MOV	$TEMP_{,-}(SP)$	FREGISTER MODE
			, IFF		
	000004	P12746	MOV	#TEMP,=(SP)	INON-REGISTER MODE
		000000			
			.ENDC		
20					
15					
22		000001	FND		

For additional information concerning addressing modes, refer to Chapter 5 and Appendix B, Section B.2.

### 7.5 .ERROR AND .PRINT DIRECTIVES

The .ERROR directive is used to output messages to the listing file during assembly pass 2. A common use of this directive is to provide a diagnostic announcement of a rejected or erroneous macro call or to alert the user to the existence of an illegal set of conditions specified in a conditional assembly. If the listing file is not specified, the .ERROR messages are output to the command output device. The .ERROR directive takes the form:

label: .ERROR expr ;text

where: label represents an optional statement label.

expr

;

represents an optional expression whose value is output when the .ERROR directive is encountered during assembly.

denotes the beginning of the text string.

text represents the specified message associated with the .ERROR directive.

Upon encountering an .ERROR directive anywhere in a source program, MACRO-11 outputs a single line containing:

1. An error code (P)

2. The sequence number of the .ERROR directive statement

3. The value of the current location counter

4. The value of the expression, if one is specified

5. The source line containing the .ERROR directive.

For example, the following directive:

.ERROR A ;INVALID MACRO ARGUMENT

causes a line in the following form to be output to the listing file:

Seq.	Loc.	Exp.	
No.	No.	Value	Text

P 512 005642 000076 .ERROR A ;INVALID MACRO ARGUMENT

The .PRINT directive is identical in function to the .ERROR directive, except that it is not flagged with the P error code.

7.6 INDEFINITE REPEAT BLOCK DIRECTIVES: .IRP AND .IRPC

An indefinite repeat block is a structure that is similar to a macro definition; essentially a macro definition that has only one dummy argument. At each expansion of the indefinite repeat range, this dummy argument is replaced with successive elements of a specified real argument list. An indefinite repeat block directive and its associated repeat range are coded in-line within the source program. This type of macro definition and expansion does not require calling the macro by name, as required in the expansion of conventional macros previously described in this section.

An indefinite repeat block can appear either within or outside another macro definition, indefinite repeat block, or repeat block (see Section 7.7). The rules for specifying indefinite repeat block arguments are the same as for specifying macro arguments (see Section 7.3).

#### 7.6.1 .IRP Directive

The .IRP directive is used to replace a dummy argument with successive real arguments specified in an argument string. This replacement process occurs during the expansion of an indefinite repeat block range. This directive takes the following form:

label: .IRP sym,<argument list>

(range of indefinite repeat block)

ENDM

•

where: label represents an optional statement label.

sym

represents a dummy argument that is successively replaced with the specified real arguments enclosed within the angle brackets. If no dummy argument is specified, the .IRP directive is flagged with an error code (A) in the assembly listing.

represents any legal separator (comma, space, and/or tab).

<argument list> represents a list of real arguments enclosed within angle brackets that is to be used in the expansion of the indefinite repeat range. A real argument may consist of one or more characters; multiple arguments must be separated by any legal separator (comma, space, and/or tab). If no real arguments are specified, no action is taken.

> range represents the block of code to be repeated once for each occurrence of a real argument in the list. The range may contain other macro definitions and repeat ranges. The .MEXIT directive (see Section 7.1.3) is legal within the range of an indefinite repeat block.

> .ENDM indicates the end of the indefinite repeat block range.

An example of the use of the .IRP directive is shown in Figure 7-1.

### 7.6.2 .IRPC Directive

The .IRPC directive is available to permit single character substitution, rather than argument substitution. On each iteration of the indefinite repeat range, the dummy argument is replaced with each successive character in the specified string. The .IRPC directive is specified as follows:

label: .IRPC sym, <string>

(range of indefinite repeat block)

. ENDM

where: label represents an optional statement label.

sym

represents a dummy argument that is successively replaced with the specified real arguments enclosed within the angle brackets. If no dummy argument is specified, the .IRPC directive is flagged with an error code (A) in the assembly listing.

represents any legal separator (comma, space, and/or tab).

<string> represents a list of characters enclosed within angle brackets to be used in the expansion of the indefinite repeat range. Although the angle brackets are required only when the string contains separating characters, their use is recommended for legibility.

range represents the block of code to be repeated once for each occurrence of a character in the list. The range may contain macro definitions and repeat ranges. The .MEXIT directive (see Section 7.1.3) is legal within the range of an indefinite repeat block.

### .ENDM

indicates the end of the indefinite repeat block range.

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An example of the use of the .IRPC directive is shown in Figure 7-1.

: - <u></u>				LTET	M.C.	
4					-1 <b>L</b>	
3						
4						
5					V	
6				, IRP	X, <aa, 00,="" td="" u<=""><td>., UU, CE, FF,</td></aa,>	., UU, CE, FF,
				MUV	A, (#0)*	
8				. ENDM		
	000000	016720	000000G	MOV	AA, (RØ)+	
	000004	016720	OPOPOOG	MOV	BB,(RØ)+	
	000010	016720	000000G	MOV	CC, (RØ)+	
	000014	016720	000000G	MOV	DD, (RØ)+	
	000020	016720	ØPDPDUG	MOV	EE, (R0)+	
	020024	016720	RABBBBG	MOV	FF. (R0)+	
0	00002-				545-54 - C C C C C C C C	
10						
1.				TRPC	X. «ABCDEE>	
:::				MOUR	HIV - (PI)	
12				MUVD	# · A J = ( A 1 J	
13				.ENDM		
	000030	112741	000000G	MOVB	$#A_{\mu} = (R1)$	
	000034	112741	000000G	MOVB	#8,=(R1)	
	000040	112741	00000GG	MOVB	#C,-(R1)	
	000044	112741	000000G	MOVB	#D,=(R1)	
	000050	112741	0000006	MOVB	#E,=(R1)	
	000054	112741	AAAAAAG	MOVB	#F.=(R1)	
14	000004			e dat de l'alter d'Alter d'Alter		
15						
16		0000011		. END		

Figure 7-1 Example of .IRP and .IRPC Directives

### 7.7 REPEAT BLOCK DIRECTIVE: .REPT, .ENDR

It is sometimes useful to duplicate a block of code a number of times in-line with other source code. This duplication of code is accomplished by creating a repeat block, using a directive in the form:

label: .REPT exp

(range of repeat block)

. ENDM

.

where: label represents an optional statement label.

exp represents any legal expression whose value controls the number of times the block of code is to be assembled within the program. When the expression value is less than or equal to zero (0), the repeat block is not assembled. If this expression is not an absolute value, the .REPT statement is flagged with an error code (A) in the assembly listing.

- range represents the block of code to be repeated the number of times determined by the specified expression value. The repeat block may contain macro definitions, indefinite repeat blocks, or other repeat blocks. The .MEXIT directive is legal within the range of a repeat block.
- ENDM indicates the end of the repeat block range. The or terminating statement in a repeat block can be
   ENDR either an .ENDM directive or an .ENDR directive.

#### 7.8 MACRO LIBRARY DIRECTIVE: .MCALL

The .MCALL directive allows you to indicate in advance those system and/or user-defined macro definitions that are required in the assembly of the source program. The .MCALL directive allows you to specify the names of all system or user macro definitions not defined within the source program but which are required to assemble the program. The .MCALL directive must appear before the first occurrence of a call to any externally-defined macro. The .MCALL directive is of the form:

.MCALL argl, arg2, ... argn

where: argl, represent the symbolic names of the macro arg2,... definitions required in the assembly of the source argn program. The symbolic macro names may be separated by any legal separator (comma, space, and/or tab).

The .MCALL directive thus provides the means to access both user-defined and system macro libraries during assembly.

The /ML switch under RSX-11 and the /LIBRARY qualifier under IAS and RT-11, specified in connection with an input file specification, indicate to MACRO-11 that the file is a macro library. When a macro call is encountered in the source program, MACRO-11 first searches the user macro library for the named macro definitions, and, if necessary, continues the search with the system macro library.

Any number of such user-supplied macro files may be designated. In cases of multiple library files, the search for the named macros begins with the last such file specified. The search continues in reverse order until the required macro definitions are found, terminating again, if necessary, with a search of the system macro library.

If any named macro is not found upon completion of the search, i.e., if the macro is not defined, the .MCALL statement is flagged with an error code (U) in the assembly listing. Furthermore, a statement elsewhere in the source program which attempts to expand such an undefined macro is flagged with an error code (O) in the assembly listing.

The command strings to MACRO-11, through which file specifications are supplied, are described in detail in the appropriate system manual (see Section 0.3 in the Preface).

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# APPENDIX A

# MACRO-11 CHARACTER SETS

A.1 ASCII CHARACTER SET

EVEN PARITY BIT	OCTAL CODE	CHARACTER	REMARKS
0 1	000 001	NUL SOH	Null, tape feed, CONTROL/SHIFT/P. Start of heading; also SOM, start of message, CONTROL/A.
1	002	STX	Start of text; also EOA, end of address, CONTROL/B.
0	003	ETX	End of text; also EOM, end of message, CONTROL/C.
1 and the	004	EOT	End of transmission (END); shuts off TWX machines, CONTROL/D.
0	005	ENQ	Enquiry (ENQRY); also WRU, CONTROL/E.
0	006	ACK	Acknowledge; also RU, CONTROL/F.
1	007	BEL	Rings the bell. CONTROL/G.
1	010	BS	Backspace; also FEO, format
•			effector. backspaces some machines, CONTROL/H.
0	011	HT	Horizontal tab. CONTROL/1.
0	012	LF	Line feed or Line space (new line); advances paper to next line, duplicated by CONTROL/J.
1	013	VT	Vertical tab (VTAB). CONTROL/K.
0	014	FF	Form Feed to top of next page (PAGE). CONTROL/L.
1	015	CR	Carriage return to beginning of line; duplicated by CONTROL/M.
1	016	SO	Shift out; changes ribbon color to red. CONTROL/N.
0	017	SI	Shift in; changes ribbon color to black. CONTROL/O.
1	020	DLE	Data link escape. CONTROL/P (DCU).
0	021	DC1	Device control 1; turns
			transmitter (READER) on, CONTROL/Q
			(X ON). U UZZ DCZ DEVICE CONCION
			CONTROL/R (TAPE, AUX ON).
1	023	DC 3	Device control 3; turns transmitter (READER) off, CONTROL/S (X OFF).
0	024	DC4	Device control 4; turns punch or auxiliary off. CONTROL/T (AUX OFF).

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EVEN PARITY BIT	7-BIT OCTAL CODE	CHARACTER	REMARKS
1	025	NAK	Negative acknowledge; also ERR, ERROR. CONTROL/U.
1	026	SYN	Synchronous file (SYNC). CONTROL/V.
0	027	ЕТВ	End of transmission block; also LEM, logical end of medium.
0 1 1 0 1 0 0 0 1 0 0 0 1 0 0 0 1 0 0 0 1 0 0 0 1 0 0 0 1 0 0 0 1 0 0 0 1 0 0 0 1 0 0 0 0 1 0 0 0 0 1 0 0 0 0 1 0 0 0 0 1 0 0 0 0 0 1 0 0 0 0 0 0 0 0 0 0 0 0 0	030 031 032 033 034 035 036 037 040 041 042 043 044 045 046 047 050 051 052 053 054 055 056 057 060 061 062 063 064 065 066 067 070 071	CAN EM SUB ESC FS GS RS US SP ! # \$ % &' () * + , /0 123456780	LEM, logical end of medium. CONTROL/W. Cancel (CANCL). CONTROL/X. End of medium. CONTROL/Y. Substitute. CONTROL/Z. Escape. CONTROL/SHIFT/K. File separator. CONTROL/SHIFT/M. Record separator. CONTROL/SHIFT/M. Necord separator. CONTROL/SHIFT/O. Space. Accent acute or apostrophe.
0 1 0 1 1 0 1 0	072 073 074 075 076 077 100 101	: ; = ? @ A	
0 1 0 1 1 0 0 0	102 103 104 105 106 107 110 111	B C D E F G H I	

EVEN Parity Bit	7-BIT OCTAL CODE	CHARACTER	REMARKS
1 0 1 0 1 0 1 1	112 113 114 115 116 117 120 121 122	J K L M N O P Q R	
0 1 0 1 1 0 0	123 124 125 126 127 130 131 132	S T U V W X Y Z	
1 0 1 1 0 0	133 134 135 136 137 140	1 1 ÷	SHIFT/K. SHIFT/L. SHIFT/M. * * Accent grave.
	141 142 143 144 145 146 147 150 151 152 153 154 155 156 157 160 161 162 163 164 165 166 167 170 171 172 173 174 175 176	a b c d e f g h i j j k l m n o P q r s t u v w x y z { l l } ~	This code generated by ALTMODE. This code generated by PREFIX key (if present). DELETE, RUBOUT.

A-3

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A.2 RADIX-50	CHARACTER	SET
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Character	ASCII Octal Equivalent	Radix-50 Equivalent
space	40	0
A-Z	101-132	1-32
\$	44	33
•	56	34
unused		35
0-9	60-71	36-47

The maximum Radix-50 value is, thus,

47\*50<sup>2</sup> +47\*50+47=174777

The following table provides a convenient means of translating between the ASCII character set and its Radix-50 equivalents. For example, given the ASCII string X2B, the Radix-50 equivalent is (arithmetic is performed in octal):

> X=113000 2=002400 B=000002 X2B=115402

SINGLE CHAR. OR FIRST CHAR.	SECOND CHARAC	fer	THIRD CHARACI	ER
Space 000000	Space	000000	Space	000000
A 003100	A	000050	A	000001
в 006200	В	000120	В	000002
C 011300	С	000170	С	000003
D 014400	D	000240	D	000004
E 017500	Е	000310	Е	000005
F 022600	F	000360	F	000006
G 025700	G	000430	G	000007
н 031000	Н	000500	н	000010
I 034100	I	000550	I	000011
J 037200	J	000620	J	000012
к 042300	K	000670	K	000013
L 045400	L	000740	L	000014
M 050500	M	001010	М	000015
N 053600	N	001060	N	000016
0 056700	0	001130	0	000017
P 062000	Р	001200	Р	000020
Q 065100	Q	001250	Q	000021
R 070200	R	001320	R	000022
S 073300	S	001370	S	000023
т 076400	$\mathbf{T} \cdot \mathbf{T}$	001440	Т	000024
U 101500	ា <b>ប</b> ដែននេះ	001510	U	000025
V 104600	- <b>V</b>	001560	V	000026
W 107700	W	001630	W	000027
X 113000	Х	001700	Х	000030
Y 116100	Y	001750	Y	000031
Z 121200	Z	002020	Z	000032
\$ 124300	\$	002070	Ş	000033

SINGLE CHA OR FIRST CHAR	R. SECONI CHARAC	SECOND CHARACTER		THIRD CHARACTER	
. 127 Unused 132 0 135 1 140 2 144 3 147 4 152 5 155 6 160 7 163 8 166	400 . 500 Unused 600 0 700 1 000 2 100 3 200 4 300 5 400 6 500 7 600 8	002140 002210 002260 002330 002400 002450 002520 002570 002640 002710 002760	Unused 0 1 2 3 4 5 6 7 8	$\begin{array}{c} 000034\\ 000035\\ 000036\\ 000037\\ 000040\\ 000041\\ 000042\\ 000043\\ 000044\\ 000045\\ 000046\\ 00004\\ 000004\\ 000004\\ 000004\\ 000004\\ 000000\\ 0000\\ 00000\\ 00000\\ 00000\\ 0000\\ 0000\\ 00000\\ 00000\\ 00000\\ 00$	
9 1/1	.700   9	003030 1	9	000047	

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#### 化二乙烯 后方的过去式和过去分词 医正常的 电脉冲

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### APPENDIX B

### MACRO-11 ASSEMBLY LANGUAGE AND ASSEMBLER DIRECTIVES

B.1 SPECIAL CHAR	ACTE	SRS
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Character	Function
:	Label terminator
=	Direct assignment operator
<b>9</b>	Register term indicator
tab	Item terminator or field terminator
space	Item terminator or field terminator
	Immediate expression indicator
0	Deferred addressing indicator
	Initial register indicator
() sei shuisgoga	Terminal register indicator
, (comma)	Operand field separator
💡 👘 संबद्ध तत्र अन्य संव	Comment field indicator
🕂 🕴 de marco de con De	Arithmetic addition operator or auto
	increment indicator
	Arithmetic subtraction operator or auto
	decrement indicator
★ 10 10 - 40 10 - 200 計算	Arithmetic multiplication operator
	Arithmetic division operator
&	Logical AND operator
	Logical OR operator
	Double ASCII character indicator
' (apostrophe)	Single ASCII character indicator or
	concatenation indicator
	Assembly location counter
	Initial argument indicator
	Terminal argument indicator
	indicator indicator
📘 🔪 — saeli de el como 🖉 el d	Macro call numeric argument indicator
vertical tab	Source line terminator

### B.2 SUMMARY OF ADDRESS MODE SYNTAX

Address mode syntax is expressed in the summary below using the following symbols: n is an integer between 0 and 7 representing a register number; R is a register expression; E is an expression; and ER is either a register expression or an expression in the range 0 to 7.

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Format	Address Mode Name	Address Mode Number	Meaning
R	Register	0n	Register R contains the operand.
@R or (ER)	Register deferred	ln	Register R contains the ad- dress of the operand.
(ER) +	Autoincrement	2n	The contents of the register specified as (ER) are incremented after being used as the address of the operand.
@(ER)+	Autoincrement Deferred	3n	The register specified as (ER) contains the pointer to the address of the operand; the register (ER) is incremented after use.
-(ER)	Autodecrement	<b>4n</b>	The contents of the register specified as (ER) are decremented before being used as the address of the operand.
@-(ER)	Autodecrement Deferred	station in the second sec	The contents of the register specified as (ER) are decremented before being used as the pointer to the address of the operand.
E(ER)	Index	бn	The expression E, plus the contents of the register specified as (ER), form the address of the operand.
@E(ER)	Index Deferred	7n	The expression E, plus the contents of the register specified as (ER), yield a pointer to the address of the operand.
#E	Immediate	27	The expression E is the operand itself.
@#E	Absolute	37	The expression E is the address of the operand.
E	Relative	67 67	The address of the operand E, relative to the instruction, follows the instruction.
<b>@ E</b>	Relative Deferred	77	The address of the operand is pointed to by E whose address, relative to the instruction, follows the instruction.

### **B.3** ASSEMBLER DIRECTIVES

The MACRO-ll assembler directives are summarized in the following table. For a detailed description of each directive, the table

в-2

contains references to the appropriate sections in the body of the manual.

Form	Section Reference	Operation
<ul> <li>(4) (4) (4) (4) (4) (4) (4) (4) (4) (4)</li></ul>	6.3.3 7.3.6	A single quote (apostrophe) followed by one ASCII character generates a word which contains the 7-bit ASCII representation of the
	1997日の中国の日本1997日 1997日の日本1997日の日本1997日の日本1997日 1997日の日本1997日の日本1997日 1997日の日本1997日の日本1997日の日本1997日の日本1997日の日本1997日の日本1997日の日本1997日の日本1997日の日本1997日の日本1997日の日本1	character in the low-order byte and zero in the high-order byte. This character is also used as a concatenation indicator in the
	inned bedi Sinner end y	expansion of macro arguments (see Section 7.3.6).
Hongan Agin daga Akin Akin Akina Akin Angin daga Akina Aking Agin daga Aking Aking Agin daga Aking Agin daga Aking daga Aking Agin daga Aking daga Aking daga Aking Agin daga Aking daga Aking daga Aking daga Aking Aking daga Aking Aking daga Aking dag	6.3.3 mars 6 (2.3.5 L) (2.0.4 (1.5.4 g) (2.6 L) (2.0.4 (1.5.4 g) (2.6 L) (2.6 L)	A double quote followed by two ASCII characters generates a word which contains the 7-bit ASCII representation of the two
na se in an ga eacada	ing appropria	characters. The first character is
요즘이 있는 상부에게 모네가 승규가지 그 모아 있는 것과 것 것 한 사람이 있었다. 같은 것 같은 것 같은 것 같은 것 같은 것 같이 있다.	nun parte ette Trisologie i 13 Turi di cinati	stored in the low-order byte; the second character is stored in the high-order byte.
<b>Bn</b> un britante de services Refin	6.4.1.2	Temporary radix control; causes the value n to be treated as a
· 가려에 감독 · · · · · · · · · · · · · · · · · ·	4877948342111447 1978年5月2日1日1月11日	binary number.
^Cexpr	6.4.2.2 99.90 1.00 1.00 1.00 1.00 1.00 1.00 1.0	Temporary numeric control; causes the expression's value to be ones- complemented.
n <b>`Dn</b> parent internet deserve and the distance of	6.4.1.2	Temporary radix control; causes the value n to be treated as a decimal number.
<b>Fn</b>	6.4.2.2	Temporary numeric control; causes
	(大学社) - 学校の時間第100 1015年4日 - 日本出版学校 - 1015年1日 - 1015年11月	the value n to be treated as a sixteen-bit floating-point number.
^On	6.4.1.2	Temporary radix control; causes
	1991年,北方陸南省南部 今日1月,陸軍和日子中北	the value n to be treated as an octal number.
Rccc	6.3.7	Convert ccc to Radix-50 form.
.ASCII /string/	6.3.4	Generates a block of data
	មេះ ១៣៩៩៩៩ ៩៩៩៩ ស្រុក ១៩៩៩ ស្រុកទី ស្រុក សន្តភាគតាល់អ្នក ស្រុក សន្តភាគតាល់អ្នក	containing the ASCII equivalent of the character string (enclosed in delimiting characters), one character per byte.

B-3

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Form	Section Reference	Operation
.ASCIZ /string/	6.3.5	Generates a block of data containing the ASCII equivalent of the character string (enclosed in delimiting characters), one character per byte, with a zero byte terminating the specified string.
ASECT	6.8.2	Begin or resume the absolute program section.
.BLKB exp	6.5.3	Reserves a block of storage space whose length in bytes is determined by the specified expression.
.BLKW exp	6.5.3	Reserves a block of storage space whose length in words is determined by the specified expression.
.BYTE expl,exp2,	6.3.1	Generates successive bytes of data; each byte contains the value of the corresponding specified expression.
.CSECT [name]	6.8.2	Begin or resume named or unnamed relocatable program section. This directive is provided for compatibility with other PDP-11 assemblers.
.DSABL arg	6.2	Disables the function specified by the argument.
.ENABL arg	6.2	Enables (invokes) the function specified by the argument.
.END [exp]	6.6.1	Indicates the logical end of the source program. The optional argument specifies the transfer address where program execution is to begin.
.ENDC	6.10.1	Indicates the end of a conditional assembly block.
.ENDM [name]	7.1.2	Indicates the end of the current repeat block, indefinite repeat block, or macro definition. The optional name, if used, must be identical to the name specified in the macro definition.
.ENDR	7.7	Indicates the end of the current repeat block. This directive is provided for compatibility with other PDP-ll assemblers.
.EOT	6.6.2	Ignored; indicates end-of-tape (which is detected automatically by the hardware). It is included for compatibility with earlier assemblers.

B-4

Form	Section Reference	Operation
.ERROR exp;text	7.5	User-invoked error directive; causes output to the listing file or the command output device containing the optional expression and the statement containing the directive.
.EVEN	6.5.1	Ensures that the current location counter contains an even address by adding l if it is odd.
.FLT2 argl,arg2,	6.4.2.1	Generates successive 2-word floating-point equivalents for the floating-point numbers specified as arguments.
.FLT4 argl,arg2,	6.4.2.1	Generates successive 4-word floating-point equivalents for the floating-point numbers specified as arguments.
.GLOBL syml,sym2,	6.9	Defines the symbol(s) specified as global symbol(s).
.IDENT /string/	6.1.5	Provides a means of labeling the object module with the program version number. The version number is the Radix-50 string appearing between the paired delimiting
.IF cond,argl	6.10.1	characters. Begins a conditional assembly block of source code which is included in the assembly only if the stated condition is met with respect to the argument(s) specified.
<pre>IFF control on the product of t</pre>	6.10.2	Appears only within a conditional assembly block, indicating the beginning of a section of code to be assembled if the condition upon entering the block tests false.
.IFT	6.10.2	Appears only within a conditional assembly block, indicating the beginning of a section of code to be assembled if the condition upon entering the block tests true.
.IFTF	6.10.2	Appears only within a conditional assembly block, indicating the beginning of a section of code to be assembled unconditionally.
.IIF cond,arg, statement	6.10.3	Acts as a 1-line conditional assembly block where the condition is tested for the argument specified. The statement is assembled only if the condition tests true.

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Form	Section Reference	Operation
.IRP sym, <argl,arg2,></argl,arg2,>	7.6.1	Indicates the beginning of an indefinite repeat block in which the symbol specified is replaced with successive elements of the real argument list enclosed within angle brackets.
.IRPC sym, <string></string>	7.6.2	Indicates the beginning of an indefinite repeat block in which the specified symbol takes on the value of successive characters, optionally enclosed within angle brackets.
.LIMIT	6.7	Reserves two words into which the Task Builder inserts the low and high addresses of the task image.
.LIST [arg]	6.1.1	Without an argument, the .LIST directive increments the listing level count by 1. With an argument, this directive does not alter the listing level count, but formats the assembly listing according to the argument specified.
.MACRO name,argl, arg2,	7.1.1	Indicates the start of a macro definition having the specified name and the following dummy arguments.
.MCALL argl,arg2,	7.8	Specifies the symbolic names of the user or system macro definitions required in the assembly of the current user program, but which are not defined within the program.
.MEXIT	7.1.3	Causes an exit from the current macro expansion or indefinite repeat block.
.NARG symbol	7.4.1	Can appear only within a macro definition; equates the specified symbol to the number of arguments in the macro call currently being expanded.
.NCHR symbol, <string></string>	7.4.2	Can appear anywhere in a source program; equates the symbol specified to the number of characters in the specified string.
.NLIST [arg]	6.1.1	Without an argument, the .NLIST directive decrements the listing level count by 1. With an argument, this directive suppresses that portion of the listing specified by the argument.

<u>.</u>

Form	Section Reference	Operation
.NTYPE symbol,aexp	7.4.3	Can appear only within a macro definition; equates the symbol to the 6-bit addressing mode of the specified address expression.
.ODD	6.5.2	Ensures that the current location counter contains an odd address by adding l if it is even.
.PAGE	6.1.6	Causes the assembly listing to skip to the top of the next page, and to increment the page count.
.PRINT exp;text	7.5	User-invoked message directive; causes output to the listing file or the command output device containing the optional expression and the statement containing the directive.
.PSECT name,attl, attn	6.8.1	Begin or resume a named or unnamed program section having the specified attributes.
.RADIX n	6.4.1.1	Alters the current program radix to n, where n is 2, 8, or 10.
.RAD50 /string/	6.3.6	Generates a block of data containing the Radix-50 equivalent of the character string enclosed within delimiting characters.
.REPT exp	7.7	Begins a repeat block; causes the section of code up to the next .ENDM or .ENDR directive to be repeated the number of times specified as exp.
.SBTTL string	6.1.4	Causes the specified string to be printed as part of the assembly listing page header. The string component of each .SBTTL directive is collected into a table of contents at the beginning of the assembly listing.
.TITLE string	6.1.3	Assigns the first six Radix-50 characters in the string as an object module name and causes the string to appear on each page of the assembly listing.
.WORD expl,exp2,	6.3.2	Generates successive words of data; each word contains the value of the corresponding specified expression.

B-7

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#### APPENDIX C

### PERMANENT SYMBOL TABLE (PST)

The permanent symbol table (PST) contains those symbols which are automatically recognized by MACRO-11. These symbols consist of both op codes and assembler directives. The op codes (i.e., the instruction set) are listed first, followed by the directives which cause specific actions during assembly.

For a detailed description of the instruction set, see the appropriate PDP-11 Processor Handbook.

С		1	OP	C	O	D	ES
~	•	-	~	~	~	-	

		4.98mm (1997)
MNEMONIC	OCTAL VALUE	FUNCTIONAL NAME
100	005500	
ADC	005500	Add Carry
ADCB	105500	Add Carry (Byte)
ADD	060000	Add Source To Destination
ASH	072000	Shift Arithmetically
ASHC	073000	Arithmetic Shift Combined
ASL	006300	Arithmetic Shift Left
ASLB	106300	Arithmetic Shift Left (Byte)
ASR	006200	Arithmetic Shift Right
ASRB	106200	Arithmetic Shift Right (Byte)
BCC	103000	Branch If Carry Is Clear
BCS	103400	Branch If Carry Is Set
BEQ	001400	Branch If Equal
BGE	002000	Branch If Greater Than Or Equal
BGT	003000	Branch If Greater Than
BHI	101000	Branch If Higher
BHIS	103000	Branch If Higher Or Same
BIC	040000	Bit Clear
BICB	140000	Bit Clear (Bvte)
BIS	050000	Bit Set
BISB	150000	Bit Set (Byte)
BIT	030000	Bit Test
BITB	130000	Bit Test (Byte)
BLE	003400	Branch If Less Than Or Equal
BLO	103400	Branch If Lower
BLOS	101400	Branch If Lower Or Same
BLT	002400	Branch If Less Than

C-1

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22

MNEMONIC	OCTAL VALUE	FUNCTIONAL NAME
BMI	100400	Branch If Minus
BNE	001000	Branch If Not Equal
BPL	100000	Branch If Plus
BPT	000003	Breakpoint Trap
BR	000400	Branch Unconditional
BVC	102000	Branch If Overflow Is Clear
BVS	102400	Branch If Overflow Is Set
CALL	004700	Jump To Subroutine (JSR PC,xxx)
ccc	000257	Clear All Condition Codes
CLC	000241	Clear C Condition Code Bit
CLN	000250	Clear N Condition Code Bit
CLR	005000	Clear Destination
CLRB	105000	Clear Destination (Byte)
CLV	000242	Clear V Condition Code Bit
CLZ	000244	Clear Z Condition Code Bit
CMP	020000	Compare Source To
		Destination
CMPB	120000	Compare Source To
		Destination (Byte)
COM	005100	Complement Destination
COMB	105100	Complement Destination
		(Byte)
DEC	005300	Decrement Destination
DECB	105300	Decrement Destination
	0.73.000	(Byte)
DIV	0/1000	Divide
EMT	104000	Emulator Trap
FADD	075000	Floating Add
FDIV	075030	Floating Divide
FMUL	075020	Floating Multiply
FSUB	075010	Floating Subtract
HALT	005200	Indit Ingroment Destination
	105200	Increment Destination
INCB	105200	(Buto)
TOT	000004	Input (Output Tran
	000100	
JSR	004000	Jump To Subroutine
MARK	006400	Mark
MFPT	006500	Move From Previous
		Instruction Space
MFPS	106700	Move from PS
		(LSI-11)
MOV	010000	Move Source To Destination
MOVB	110000	Move Source To Destination
		(Byte)
MTPI	006600	Move To Previous
		Instruction Space
MTPS	106400	Move to PS
		(LSI-11)
MUL	070000	Multiply
NEG	005400	Negate Destination
NEGB	105400	Negate Destination (Byte)
NOP	000240	No Operation
RESET	000005	Reset External Bus
RETURN	000207	Return From Subroutine (RTS PC)
ROL	006100	Rotate Left
ROLB	106100	Rotate Left (Byte)
ROR	006000	Rotate Right
1	1	

MNEMONIC	OCTAL VALUE	FUNCTIONAL NAME
RORB	106000	Rotate Right (Byte)
RTI	000002	Return From Interrupt (Permits a trace trap)
RTS	000200	Return From Subroutine
RTT	000006	Return From Interrupt
	a déné sa na saétar	(inhibits trace trap)
SBC	005600	Subtract Carry
SBCB	105600	Subtract Carry (Byte)
SCC	000277	Set All Condition Code Bits
SEC	000261	Set C Condition Code Bit
SEN	000270	Set N Condition Code Bit
SEV	000262	Set V Condition Code Bit
SEZ	000264	Set Z Condition Code Bit
SOB	077000	Subtract One And Branch
SUB	160000	Subtract Source From Destination
SWAB	000300	Swap Bytes
SXT	006700	Sign Extend
TRAP	104400	Trap
TST	005700	Test Destination
TSTB	105700	Test Destination (Byte)
WAIT	000001	Wait For Interrupt
XOR	074000	Exclusive OR

# OP CODES FLOATING POINT PROCESSOR ONLY

MNEMONIC	OCTAL VALUE	FUNCTIONAL NAME
ABSD	170600	Make Absolute Double
ABSF	170600	Make Absolute Floating
ADDD	172000	Add Double
ADDF	172000	Add Floating
CFCC	170000	Copy Floating Condition
		Codes
CLRD	170400	Clear Double
CLBE	170400	Clear Floating
CMPD	173400	Compare Double
CMPF	173400	Compare Floating
DIVD	174400	Divide Double
DIVF	174400	Divide Floating
LDCDF	177400	Load And Convert From
		Double To Floating
LDCFD	177400	Load And Convert From
	和 이 이 이 이 가 가 가 가 가 가 다.	Floating To Double
LDCID	177000	Load And Convert Integer To
	1	Double
LDCIF	177000	Floating
LDCLD	177000	Load And Convert Long
	177000	Integer To Double
LDCLF	177000	Integer To Floating
LDD	172400	Load Double
LDEXP	176400	Load Exponent
	1	

C-3

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22

MNEMONIC	OCTAL VALUE	FUNCTIONAL NAME
LDF	172400	Load Floating
LDFPS	170100	Load FPPs Program Status
MFPD	106500	Move From Previous Data Space
MODD	171400	Multiply And Integerize Double
MODF	171400	Multiply And Integerize Floating
MTPD	106600	Move To Previous Data Space
MULD	171000	Multiply Double
MULF	171000	Multiply Floating
NEGD	170700	Negate Double
NEGF	170700	Negate Floating
SETD	170011	Set Double Mode
SETF	170001	Set Floating Mode
SETI	170002	Set Integer Mode
SETL	170012	Set Long Integer Mode
SPL	000230	Set Priority Level
STCDF	176000	Store And Convert From Double To Floating
STCDI	175400	Store And Convert From Double To Integer
STCDL	175400	Store And Convert From Double To Long Integer
STCFD	176000	Store And Convert From Floating To Double
STCFI	175400	Store And Convert From Floating To Integer
STCFL	175400	Store And Convert From Floating To Long Integer
STD	174000	Store Double
STEXP	175000	Store Exponent
STF	174000	Store Floating
STFPS	170200	Store FPPs Program Status
STST	170300	Store FPPs Status
SUBD	173000	Subtract Double
SUBF	173000	Subtract Floating
TSTD	170500	Test Double
TSTF	170500	Test Floating

### C.2 MACRO-11 DIRECTIVES

DIRECTIVE	FUNCTIONAL SIGNIFICANCE
ASCII	Translates character string to ASCII equivalents
.ASCIZ	Translates character string to ASCII equivalents; inserts zero byte as last character.
ASECT	Begins absolute program section (provided for compatibility with other PDP-11 assembliers).
.BLKB	Reserves byte block in accordance with value of specified argument.
.BLKW	Reserves word block in accordance with value of specified argument.
.BYTE	Generates successive byte data in accordance with
.CSECT	Begins relocatable program section (provided for compatibility with other PDP-11 assemblers).

DIRECTIVE	FUNCTIONAL SIGNIFICANCE
.DSABL .ENABL	Disables specified function. Enables specified function.
.END	Defines logical end of source program.
.ENDC	Defines end of conditional assembly block.
.ENDM	Defines end of macro definition, repeat block, or
	indefinite repeat block.
.ENDR	Defines end of current repeat block (provided for
	compatibility with other PDP-11 assemblers).
.EOT	Define End of Tape condition (ignored).
.ERROR	Outputs diagnostic message to listing file or
	command output device.
.EVEN	word-aligns the current location counter.
• • • • • • •	each floating-point argument.
FLTA .	Causes four words of storage to be generated for
*****	each floating-point argument.
GLOBL	Declares global attribute for specified symbol(s).
IDENT	Labels object module with specified program
	version number.
.IF	Begins conditional assembly block.
.IFF	Begins subconditional assembly block (if
	conditional assembly block test is false).
.IFT	Begins subconditional assembly block (if
-	conditional assembly block test is true).
.IFTF	Begins subconditional assembly block (whether
	conditional assembly block test is true of faise).
• 1 1 F	(if aposified condition is satisfied)
TDD	Begins indefinite repeat block: replaces
• 1 1 1	specified symbol with specified successive real
	arguments.
.IRPC	Begins indefinite repeat block; replaces
	specified symbol with value of successive
	characters in specified string.
.LIMIT	Reserves two words of storage for high and low
	addresses of task image.
LIST	Controls listing level count and lormat of
	definition
MCALL	Identifies required macro definition(s) for
·HCALL	assembly.
.MEXIT	Exit from current macro definition or indefinite
	repeat block.
.NARG	Equates specified symbol to the number of
	arguments in the macro expansion.
.NCHR	Equates specified symbol to the number of
	characters in the specified character string.
.NLIST	Controls listing level count and suppresses
NUTYDE	Equator specified symbols to the addressing mode
.NIIPE	of the specified argument.
ממס	Byte-aligns the current location counter.
PAGE	Advances form to top of next page.
.PRINT	Prints specified message on command output device.
. PSECT	Begins specified program section having specified
	attributes.
.RADIX	Changes current program radix to specified radix.
.RAD50	Generates data block having Radix-50 equivalents
	of specified character string.
.REPT	Begins repeat block and replicates it according to
	the value of the specified expression.

DIRECTIVE	FUNCTIONAL SIGNIFICANCE
.SBTTL	Prints specified subtitle text as the second line of the assembly listing page header.
.TITLE	Prints specified title text as object module name in the first line of the assembly listing page header.
.WORD	Generates successive word data in accordance with specified arguments.

The MACRO-11 directives listed above are summarized in greater detail in Appendix B.

### APPENDIX D

### DIAGNOSTIC ERROR MESSAGE SUMMARY

#### D.1 MACRO-11 ERROR CODES

A diagnostic error code is printed as the first character in a source line which contains an error detected by MACRO-11. This error code identifies a syntactical problem or some other type of error condition detected during the processing of a source line. An example of such a source line is shown below:

Q 26 000236 010102 MOV R1,R2,A

The extraneous argument A in the MOV instruction above causes the line to be flagged with a Q (syntax) error.

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Error Code	Meaning
A	Assembly error. Because many different types of error conditions produce this diagnostic message, all the possible directives which may yield a general assembly error have been categorized below to reflect specific classes of error conditions:
	.RADIX A value other than 2, 8, or 10 is specified as a new radix.
	.LIST/.NLIST Other than a legally defined argument (see Table 6-1) is specified with the directive.
	.ENABL/.DSABL Other than a legally defined argument (see Table 6-2) is specified with the directive.
	.PSECT Other than a legally-defined argument (see Table 6-3) is specified with the directive.
	.IF/.IIF Other than a legally defined conditional test (see Table 6-5) or an illegal argument expression value is specified with the directive.
	.MACRO An illegal or duplicate symbol found in dummy argument list.

### DIAGNOSTIC ERROR MESSAGE SUMMARY

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Error Code	Meaning
A (Cont'd)	CATEGORY 2: NULL ARGUMENT OR SYMBOL SPECIFIED.
	.TITLE Program name is not specified in the directive, or first non-blank character following the directive is a non-Radix-50 character.
	.IRP/.IRPC No dummy argument is specified in the directive.
	.NARG/.NCHAR/.NTYPE No symbol is specified in the directive.
	.IF/.IIF No conditional argument is specified in the directive.
	CATEGORY 3: UNMATCHED DELIMITER/ILLEGAL ARGUMENT CONSTRUCTION.
	.ASCII/.ASCIZ/.RAD50/.IDENT Character string or argument string delimiters do not match, or an illegal character is used as a delimiter, or an illegal argument construction is used in the directive.
	.NCHAR Character string delimiters do not match, or an illegal character is used as a delimiter in the directive.
	CATEGORY 4: GENERAL ADDRESSING ERRORS.
an a	This type of error results from one of several possible conditions:
	<ol> <li>Permissible range of a branch instruction, i.e., from -128(10) to +127(10) words, has been exceeded.</li> </ol>
	<ol> <li>A statement makes invalid use of the current location counter, e.g., a ".=expression" statement attempts to force the current location counter to cross program section (.PSECT) boundaries.</li> </ol>
	3. A statement contains an invalid address expression. In cases where an absolute address expression is required, specifying a global symbol, a relocatable value, or a complex relocatable value (see Section 3.9) results in an invalid address expression. Similarly, in cases where a relocatable address expression is required oither a
	relocatable or absolute value is permissible, but a global symbol or a complex relocatable value in the statement likewise results in an invalid address expression. Specific cases of this type of error are those which follow:

### DIAGNOSTIC ERROR MESSAGE SUMMARY

Error Code	Meaning
	.BLKB/.BLKW/.REPT Other than an absolute value or an expression which reduces to an absolute value has been specified with the directive.
	<ol> <li>Multiple expressions are not separated by a comma. This condition causes the next symbol to be evaluated as part of the current expression.</li> </ol>
	CATEGORY 5: ILLEGAL FORWARD REFERENCE.
	This type of error results from either of two possible conditions:
	1. A global assignment statement (symbol==expression) contains a forward reference to another symbol.
	<ol> <li>An expression defining the value of the current location counter contains a forward reference.</li> </ol>
В	Bounding error. Instructions or word data are being assembled at an odd address. The location counter is incremented by 1.
	Doubly-defined symbol referenced. Reference was made to a symbol which is defined more than once.
	End directive not found. When the end-of-file is reached during source input and the .END directive has not yet been encountered, MACRO-11 generates this error code, ends assembly pass 1, and proceeds with assembly pass 2.
I	Illegal character detected. Illegal characters which are also non-printable are replaced by a question mark (?) on the listing. The character is then ignored.
<b>L</b> 22.2	Input line is greater than 132(10) characters in length. Currently, this error condition is caused only through excessive substitution of real arguments for dummy arguments during the expansion of a macro.
М	Multiple definition of a label. A label was encountered which was equivalent (in the first six characters) to a label previously encountered.
N	A number contains a digit that is not in the current program radix. The number is evaluated as a decimal value.
O	Opcode error. Directive out of context. Permissible nesting level depth for conditional assemblies has been exceeded. Attempt to expand a macro which was unidentified after .MCALL search.

### DIAGNOSTIC ERROR MESSAGE SUMMARY

3.

Error Code	Meaning
P	Phase error. A label's definition of value varies from one assembly pass to another or a multiple definition of a local symbol has occurred within a local symbol block. Also, when in a local symbol block defined by the .ENABL LSB directive, an attempt has occurred to define a local symbol in a program section other than that which was in effect when the block was entered. A P error code also appears if an .ERROR directive is assembled.
Q	Questionable syntax. Arguments are missing, too many arguments are specified, or the instruction scan was not completed.
R	Register-type error. An invalid use of or reference to a register has been made, or an attempt has been made to redefine a standard register symbol without first issuing the .DSABL REG directive.
T T	Truncation error. A number generated more than 16 bits in a word, or an expression generated more than 8 significant bits during the use of the .BYTE directive or trap (EMT or TRAP) instruction.
	Undefined symbol. An undefined symbol was encountered during the evaluation of an expression; such an undefined symbol is assigned a value of zero. Other possible conditions which result in this error code include unsatisfied macro names in the list of .MCALL arguments and a direct assignent (symbol=expression) statement which contains a forward reference to a symbol whose definition also contains a forward reference; also, a local symbol may have been referenced that does not exist in the current local symbol block.
Z	Instruction error. The instruction so flagged is not compatible among all members of the PDP-11 family. See Section 5.3 for details.

D-4

#### APPENDIX E

#### SAMPLE CODING STANDARD

#### E.1 INTRODUCTION

Standards eliminate variability and the requirement to make a decision. Much of the difficulty in establishing standards stems from the notion that they should be optimal. However, to be successfully applied, standards must represent an agreement on certain aspects of the programming process.

This Appendix contains DIGITAL'S PDP-11 Program Coding Standard. It is suggested that this be used as a model to assist users in preparing standards for their own installations.

#### E.2 LINE FORMAT

All source lines shall consist of from one to a maximum of eighty characters (not including the audit trail added by SLIPR (SLP in RSX-11M) editor. This program is described in the applicable <u>RSX-11M</u> or <u>RSX-11D Utilities Manual</u> or in the <u>IAS Editing Utilities Reference</u> Manual (see Section 0.3 in the Preface).

Assembly language code lines shall have the following format:

- Label Field if present, the label shall start at tab stop 0 (column 1).
- Operation field the operation field shall start at tab stop 1 (column 9).
- Operand field the operand field shall start at tab stop 2 (column 17).
- Comments field the comments field shall start at tab stop 4 (column 33) and may continue to column 80.

Comment lines that are included in the code body shall be delimited by a line containing only a leading semicolon. The comment itself contains a leading semicolon and starts in column 3. Indents shall be 1 tab.

If the operand field extends beyond tab stop 4 (column 33) simply leave a space and start the comment. Comments which apply to an instruction but require continuation should always line up with the character position which started the comment.

E-1

#### E.3 COMMENTS

Comment all coding to convey the global role of an instruction, rather than simply a literal translation of the instruction into English. In general this will consist of a comment per line of code. If a particularly difficult, obscure, or elegant instruction sequence is used, a paragraph of comments must immediately precede that section of code.

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Preface text, which describes formats, algorithms, program-local variables, etc., will be delimited by the character sequence ;+ at the start of the text and ;- at the end; these delimiters facilitate automated extraction of narrative commentary. The comment itself will start in column 3.

For example:

;+ ; THE INVERT ROUTINE ACCEPTS ; A LIST OF RANDOM NUMBERS AND ; APPLIES THE KOLMOGOROV ALGORITHM ; TO ALPHABETIZE THEM. ;-

E.4 NAMING STANDARDS

E.4.1 Register Standards

E.4.1.1 General Purpose Registers - Only the following names are permitted as register names; and may not be used for any other purpose:

R0=%0	;REG 0
R1=%1	;REG 1
R2=%2	; REG 2
R3=%3	REG 3
R4=%4	REG 4
R5=%5	REG 5
SP=%6	STACK POINTER (REG 6)
PC=%7	; PROGRAM COUNTER (REG 7)

E.4.1.2 Hardware Registers - These registers must be named identically to the hardware definition. For example, PS and SWR.

E.4.1.3 Device Registers - These are symbolically named identically to the hardware notation. For example, the control status register for the RK disk is RKCS. Only this symbolic name may be used to refer to this register.
### SAMPLE CODING STANDARD

## E.4.2 Processor Priority

Testing or altering the processor priority is done using the symbols

PR0, PR1, PR2, .....PR7

which are equated to their corresponding priority bit pattern.

## E.4.3 Other Symbols

Frequently-used bit patterns such as CR and LF will be made conventional symbolics on an as-needed basis.

# E.4.4 Using the Standard Symbolics

The register standards will be defined within the assembler. All other standard symbols will appear in a file and will be linked prior to program execution.

### E.4.5 Symbols\*

E.4.5.1 Global Symbols - Global symbols should be easily recognized by their format. The following standards apply and completely define symbol standards for PDP-11 Medium/Large software products.

symbol	pos-1	pos-2	pos-3	pos-4	pos-5	po <b>s-</b> 6	length
non-glbl-sym	letter	a-num/ null	a-num/ null	a-num/ null	a-num/ null	a-num/ null	>=1
glbl-sym	\$/.	a-num/ null	a-num/ null	a-num/ null	a-num/ null	a-num/ null	>=1
glbl-offset	letter	\$/.	a-num	a-num/ null	a-num/ null	a-num/ null	>=3
glbl-bit-ptrn	letter	a-num	\$/.	a-num/	a-num/ null	a-num/ null	>=4
local-sym	number *	Ś					>=2

\* Symbols that are branch targets are also called labels, but we will always use the term "symbol".

\*\* Number is in the range 0<number<65535.

\*\*\* The use of \$ or . for global names is reserved for DEC-supplied software.

where:

a-num non-glbl-sym	is an alphanumeric character.
	least weeks la state of the least
local-sym	local sympols, as defined by
	MACRO-11.
qlbl-sym	are global symbols (addresses).
albl-offact	are global offering (abgelute
gibi-orrsec	are grobal orrsets (absolute
	quantities).
glbl-bit-ptrn	are global bit patterns.

A program never contains a .GLOBL statement without showing cause.

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E.4.5.2 Symbol Examples

Non-Global Symbols

AlB

ZXCJ1

INSRT

Global Address Symbols

ŞJIM

.VECTR

**\$SEC** 

Global Absolute Offset Symbols

AŞJIM

AŞXT

A.ENT

Global Bit Pattern Symbols

A1\$20

B3.6

JI.M

Local Symbols

37\$

271\$

6\$

E.4.5.3 **Program-Local Symbols** - Self-relative address arithmetic (.+n) is absolutely forbidden in branch instructions; its use in other contexts must be avoided if at all possible and practical.

Target symbols for branches that exist solely for positional reference will use local symbols of the form

### <num>\$:

Use of non-local symbols is restricted, within reason, to those cases where reference to the code occurs external to the code. Local-symbols are formatted such that the numbers proceed sequentially down the page and from page to page.

E.4.5.4 Macro Names - The last two characters (with the last character possibly being null) have special significance. The next to last character is a \$, the last, a character specifying the mode of the macro.

For example, in the three macro forms in-line, stack, and p-section, the in-line form has no suffix, the stack has an  $\langle S \rangle$ , and the p-section a  $\langle C \rangle$ . Thus the RSX Queue I/O macro can be written as any of

QIOŞ

QIO\$S

QIO\$C

depending on the form required. These are not reserved letters. Only the form of the name is standard.

### E.5 PROGRAM MODULES

### E.5.1 General Comments on Programs

In our software, a program provides a single distinct function. No limits exist on size, but the single function limitation should make modules larger than 1K a rarity. Since any software may eventually exploit the virtual memory capacity of the 11/40 and 11/45, programs should make every attempt to maintain a dense reference locus (do not promiscuously branch over page boundaries or over a large absolute address distance).

All code is read-only. Code and data areas are distinct and each contains explanatory text. Read-only data should be segregated from read-write data.

### E.5.2 The Module Preface

Each program module in the system shall exist as a separate file. The file name will reflect the name of the module and the file type shall be of the form 'NNN'. The 'NNN' signifies the edit number or the version number. The version number shall be changed only when a new base level is created. Furthermore, if no corrections are made to a file from one base level to the next, the version number will not be changed. The availability of File Control Services and File Control Primitives will greatly simplify version number maintenance. Program modules adhere to a strict format. This format adds to the readability and understandability of the module. The following sections are included in each module: For the Code Section:

- A .TITLE statement that specifies the name of the module. If a module contains more than one routine, subtitles may be used.
- 2. An .IDENT statement specifying the version number. The PDP-11 version number standard appears in section E.10.

3

12

- 3. A .PSECT statement that defines the program section in which the module resides.
- 4. A copyright statement, and the disclaimer.

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5. The version number of the file.

The PDP-11 version number standard is described in section E.10.

- 6. The name of the principal author and the date on which the module was first created.
- 7. The name of each modifying author and the date of modification. Names and modification dates appear one per line and in chronological order.
- 8. A brief statement of the function of the module.

Note: Items 1-8 should appear on the same page.

- 9. A list of the definitions of all equated local symbols used in the module. These definitions appear one per line and in alphabetical order.
- 10. All local macro definitions, preferably in alphabetical order by name.
- 11. All local data. The data should indicate

a. Description of each element (type, size, etc.)b. Organization (functional, alpha, adjacent, etc.)c. Adjacency requirements

## SAMPLE CODING STANDARD

- 12. A more detailed definition of the function of the module.
- 13. A list of the inputs expected by the module. This includes the calling sequence if non-standard, condition code settings, and global data settings.
- 14. A list of the outputs produced as a result of entering this module. These include delivered results, condition code settings, but not side effects. (All these outputs are visible to the caller.)
- 15. A list of all effects (including side effects) produced as a result of entering this module. Effects include alterations in the state of the system not explicitly expected in the calling sequence, or those not visible to the caller.
- 16. The module code.

### E.5.3 Formatting the Module Preface

Rules:

- The first eight items appear on the same page and will not have explicit headings. Item 3 may be omitted if the blank p-section is being used.
- Headings start at the left margin\*; descriptive text is indented 1 tab position.
- 3. Items 7-14 will have headings which start at the left margin, preceded and followed by lines containing only a leading <;>. Items which do not apply may be omitted.

A template for the module preface follows.

FILE-EXAMPL.SO1

.TITLE	EXAMPLE
.IDENT	/01/
.PSECT	KERNEL

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\*The left margin consists of a <;> a <space> then the heading, so the text of the heading begins in column 3.

# SAMPLE CODING STANDARD

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```
; DEC ASSUMES NO RESPONSIBILITY FOR THE USE OR RELIABILITY OF
; ITS SOFTWARE ON EQUIPMENT WHICH IS NOT SUPPLIED BY DEC.
; VERSION 01
;
; JOE PASCUSNIK 1-JAN-72
;
; MODIFIED BY:
        RICHARD DOE 21-JAN-73
;
        SPENCER THOMAS 12-JUN-73
:
;
     Brief statement of the module's function
;
;
; EQUATED SYMBOLS
;
        List equated symbols
;
; LOCAL MACROS
        Local Macros
;
 LOCAL DATA
;
;
        Local data
;+
; Module function-details
;
; INPUTS:
;
        Description of inputs
;
; OUTPUTS:
;
        Description of outputs
;
;
; EFFECTS:
;
        Description of effects
;
;-
        Begin Module Code
```

# E.5.4 Modularity

No other characteristic has more impact on the ultimate engineering success of a system than does modularity. Modularity for PDP-11 Software Engineering's products consists of the application of the single-function philosophy described in section E.5.1, and adherence to a set of calling and return conventions.

E.5.4.1 Calling Conventions (Inter-Module) - The following calling conventions must be observed.

Transfer of Control

Macros will exist for call and return. The actual transfer will be via a JSR PC instruction. For register save routines, a JSR Rn,SAVE will be permitted.

The CALL macro is:

CALL subr-name

The RETURN macro is:

RETURN

Register Conventions

On entry, a subroutine minimally saves all registers it intends to alter except result registers. On exit it restores these registers. (State preservation is assumed across calls.)

Argument Passing

Any registers may be used, but their use should follow a coherent pattern. For example, if passing three arguments, pass them in R0, R1 and R2 rather than R0, R2, R5. Saving and restoring occurs in one place.

E.5.4.2 Exiting - All subroutine exits occur through a single RETURN macro.

E.5.4.3 Intra-Module Calling Conventions - Designer optional, but consistency favors a calling sequence identical to that of the inter-module sequence.

E.5.4.4 Success/Failure Indication - The C bit will be used to return the success/failure indicator, where success equals 0, and failure equals 1. The argument registers can be used to return values or additional success/failure data.

E.5.4.5 Module Checking Routines - Modules are responsible for verifying the validity of arguments passed to them. The design of a module's calling sequence should aim at minimizing the validity checks by minimizing invalid combinations. Programmers may add test code to perform additional checks during checkout. All code should aim at discovering an error as close (in terms of instruction executions) to its occurrence as possible.

E.6 FORMATTING STANDARDS

E.6.1 Program Flow

Programs will be organized on the listing such that they flow down the page, even at the cost of an extra branch or jump.

For example:



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shall appear on the listing as:

	TST	
	BNE	BBB
AAA:	• • • •	• • • •
	• • • •	• • • •
	••••	
	BR	CMN
BBB:	••••	• • • •
	• • • •	• • • •
	• • • •	• • • •
CMN:	• • • •	• • • •
	• • • •	• • • •
	••••	••••

Rather than:

	TST BNE	BBB
AAA:	••••	••••
	••••	••••
CMN:		••••
0		
	••••	••••
BBB:	••••	• • • •
	••••	• • • • • • • •
	BR	CMN

# E.6.2 Common Exits

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A common exit appears as the last code sequence on the listing. Thus the flow chart:



# will appear on the listing as:

PR1:	••••	
	••••	
	BR	EXIT
PR2:		사람은 사람은 동네는 것은 사람은 것이 있는 것이 있는 것이 있다. 가지 않는 것이 있는 것이 있다. 이 같은 바빠에 바빠 이 있는 것이 있다.
		n de 1997 - Constantino de la companya de la constantina de la constantina de la constantina de la constantina La constantina de la c
	• • • •	
	BR	EXIT
<b>DD</b> 2 .		and the state of the second state of the
PR3:	••••	
	• • • •	
	BR	EXIT
PR4:		
	••••	• • • •
	• • • •	••••
EXIT:		
And not	as:	
PR1:		
		literet and the second se
	• • • •	
EVIT.		
EXII:	• • • •	••••
PR2:	• • • •	
	• • • •	
	••••	•••• EVTM
	BR	EVII
PR3:		
		••••
	BR	EXIT
PR4:	•••	••••
	· · · · · · · · ·	
	BR	EXIT
	٩	E-11

# E.6.3 Code with Interrupts Inhibited

Code that is executed with interrupts inhibited, shall be flagged by a three semicolon (;;;) comment delimiter. For example:

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ERTZ	:		;ENABLE BY RETURNING ;BY SYSTEM SUBROUTINES,
	BIS BIT BEQ RTT	#PR7,PS #PR7,+2(SP) 10\$	;;; INHIBIT INTERRUPTS ;;; C ;;; O ;;; M
10\$:	• • • •	• • • •	;;; E
	• • • •	• • • •	;;; N
	••••	• • • •	;;; T
		• • • •	;;; S

### E.7 PROGRAM SOURCE FILES

Source creation and maintenance shall be done in base levels. A base level is defined as a point at which the program source files have been frozen. From the freeze point to the next base level, corrections will not be made directly to the base level itself. Rather a file of corrections shall be accumulated for each file in the base level. Whenever an updated source file is desired, the correction file will be applied to the base file.

The accumulation of corrections shall proceed until a logical breaking point has occurred (i.e. a milestone or significant implementation point has been reached). At this time all accumulated corrections shall be applied to the previous base level to create a new base level. Correction files will then be started for the new base level.

## E.8 FORBIDDEN INSTRUCTION USAGE

- 1. The use of instructions or index words as literals of the previous instruction. For example:
  - MOV @PC,Register
  - BIC Src,Dst

uses the bit clear instruction as a literal. This may seem to be a very "neat" way to save a word but what about maintaining a program using this trick? To compound the problem, it will not execute properly if I/D space is enabled on the 11/45. In this case @PC is a D bank reference.

 The use of the MOV instruction instead of a JMP instruction to transfer program control to another location. For example:

MOV #ALPHA, PC

transfers control to location ALPHA. Besides taking longer to execute (2.3 microseconds for MOV vs. 1.2 for JMP) the use of MOV instead of JMP makes it nearly impossible to pick up someone else's program and tell where transfers of control

### SAMPLE CODING STANDARD

take place. What if one would like to get a jump trace of the execution of a program (a move trace is unheard of)? As a more general issue, perhaps even other operations such as ADD and SUB from PC should be discouraged. Possibly one or two words can be saved by using these operations but how many such occurrences are there?

3. The seemingly "neat" use of all single word instructions where one double-word instruction could be used and would execute faster and would not consume additional memory. Consider the following instruction sequence:

CMP - (R1), (-R1)

CMP - (R1), - (R1)

The intent of this instruction sequence is to subtract 8 from register Rl (not to set condition codes). This can be accomplished in approximately 1/3 the time via a SUB instruction (9.4 vs. 3.8 microseconds) at no additional cost in memory space. Another question here is also, what if Rl is odd? SUB always wins since it will always execute properly and is always faster!

# E.9 RECOMMENDED CODING PRACTICE

### E.9.1 Conditional Branches

When using the PDP-11 conditional branch instructions, it is imperative that the correct choice be made between the signed and the unsigned branches.

SIGNED	UNSIGNED		
BGE	BHIS	(BCC)	
BLT	BLO		
BGT	BHI		
BLE	BLOS	(BCS)	

A common pitfall is to use a signed branch (e.g. BGT) when comparing two memory addresses. All goes well until the two addresses have opposite signs; that is, one of them goes across the 16K (100000(8)) bound. This type of coding error usually shows itself as a result of re-linking at different addresses and/or a change in size of the program.

## E.10 PDP-11 VERSION NUMBER STANDARD

The PDP-11 Version Number Standard applies to all modules, parameter files, complete programs, and libraries which are written or caused to be written, as part of the PDP-11 Software Development effort. It is used to provide unique identification of all released, pre-released, and in-house software.

It is limited in that, as currently specified, only six characters of identification are used. Future implementations of the Macro Assembler, linker, and librarian should provide for at least nine

characters, and possibly twelve. It is expected that this standard will be enhanced as the need arises.

Version Identifier = <form> <version> <edit> <patch>

<form> Used to identify a particular form of a module or program, where applicable, as in the case of LINK-11. One alphabetic character, if used, and null (i.e., a binary 0) if not used.

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<version> Used to identify the release, or generation, of a program. Two decimal digits, starting at 00, and incremented at the discretion of the project in order to reflect what, in their opinion, is a major change.

- <edit> Used to identify the level to which a particular release, or generation, of a program or module has been edited. An edit is defined to be an alteration to the source form. Two decimal digits, beginning at 01, and incremented with each edit; null if no edits.
- <patch> Used to identify the level to which a particular release, or generation, of a program or module has been patched. A patch is defined as an alteration to a binary form. One alphabetic character, starting at B, and running sequentially toward Z, each time a set of patches is released; null if no patches.

These fields are interrelated. When <version> is changed, then <patch> and <edit> must be reset to nulls. It is intended that when <edit> is incremented, then <patch> will be re-set to null, because the various bugs have been fixed.

# E.10.1 Displaying the Version Identifier

The visible output of the version identifier should appear as:

<key-letter> <form> <version> - <edit> <patch>,

where the following Key Letters have been identified:

- V released or frozen version
- X in-house experimental version
- Y field test, pre-release, or in-house release version

Note that 'X' corresponds roughly to individual support, 'Y' to group support, and 'V' to company support.

The dash which separates <version> from <edit> is used only if <edit> and/or <patch> is not null. When a version identifier is displayed as part of program identification, then the format is:

Program

<space><key-letter><form><version>-<edit><patch>
Name

Examples:

PIP X03 LINK VB04-C MACRO Y05-01

### E.10.2 Use of the Version Number in the Program

All sources must contain the version number in an .IDENT directive. For programs (or libraries) which consist of more than one module, each individual module will follow this version number standard. The version number of the program or library is not necessarily related to the version numbers of the constituent modules; it is perfectly reasonable, for example, that the first version of a new FORTRAN library, V00, contain an existing SIN routine, say V05-01.

Parameter files are also required to contain the version number in an .IDENT directive. Because the assembler records the last .IDENT seen, parameter files must precede the program.

Entities which consist of a collection of modules or programs, e.g., the FORTRAN Library, will have an identification module in the first position. An identification module exists solely to provide identification, and normally consists of something like:

> ;OTS IDENTIFICATION .TITLE FTNLIB .IDENT /003010/ .END

an Arean Araba Araba Araba Araba Araba Araba Araba Araba

### APPENDIX F

# ALLOCATING VIRTUAL MEMORY

This appendix is intended for the MACRO-11 user who wants to avoid the problem of thrashing, by optimizing the allocation of virtual memory. Users of smaller systems, particularly those with the 8K subset version of MACRO-11, should become thoroughly familiar with the conventions discussed herein. In this regard, Appendix F addresses the following topics:

- 1. General hints and space-saving guidelines
- 2. Macro definitions and expansions
- 3. Operational techniques.

The user is assumed to have pursued a policy of modular programming, as advised in Appendix E. In addition to the obvious advantages accruing from small, distinct, highly-functional bodies of code, one can usually avoid the problem of insufficient dynamic memory during assembly by practicing such a policy. Other suggestions as to how available memory can be best utilized are discussed in the following sections.

### F.1 GENERAL HINTS AND SPACE-SAVING GUIDELINES

Work-file memory is shared by a number of MACRO-ll's tables, each of which is allocated space on demand (64K words of dynamically pageable storage are available to the assembler). The tables and their corresponding entry sizes are as follows:

User-defined symbols - five words.

- 2. Local symbols four words.
- 3. Program sections six words.
- 4. Macro names four words.
- 5. Macro text nine words.
- 6. Source files six words.

In addition, several scratch pad tables are used during the assembly process, as follows:

- 1. Expression analysis five words.
- 2. Object code generation five words.

- 3. Macro argument processing three words.
- 4. .MCALL argument processing five words.

The above information can serve as a guide for estimating dynamic storage requirements and for determining ways to reduce such requirements.

For example, the use of local symbols whenever possible is highly encouraged, since their internal representation requires 25% less dynamic storage than that required for regular user-defined symbols. The usage of local symbols can often be maximized by extending the scope of local symbol blocks through the .ENABL LSB/.DSABL LSB MACRO-11 directives (see Sections 3.5 and 6.2).

Since MACRO-11 does not support a purge function, once a symbol is defined, it permanently occupies its dynamic memory allocation. Numerous instances occur during conditional assemblies and repeat loops when a temporarily assigned symbol is used as a count or offset indicator. If possible, the symbols so used should be re-used.

In keeping with the same principle, special treatment should be given to the definition of commonly-used symbols. Instead of simply appending a prefix file which defines all possibly-used symbols for each assembly, users are encouraged to group symbols into logical classes. Each class so grouped can then become a shortened prefix file or a macro in a library (see Section F.2 below). In either case, selective definition of symbolic assignments is achieved, resulting in fewer defined (but unreferenced) symbols.

An appropriate example of this idea is seen in the definition of standard symbols. The system macro library, for example, supplies several macros used to define distinct classes of symbols. These groupings and associated macro names are, as follows:

DRERR\$ - Directive return status codes IOERR\$ - I/O return status codes FILIO\$ - File-related I/O function codes SPCIO\$ - Special I/O function codes

# F.2 MACRO DEFINITIONS AND EXPANSIONS

By far, dynamic storage is used most heavily for the storage of macro text. Upon macro definition or the issuance of an .MCALL directive, the entire macro body is stored, including all comments appearing in the macro definition. For this reason, comments should not be included as part of the macro text. An RSX-11 utility program (called SQZ for RSX-11D only) and a Librarian function switch (/SZ) are available to compress macro source text by removing all trailing blanks and tabs, blank lines, and comments. The system macro library (RSXMAC.SML) has already been compressed. User-supplied macro libraries (.MLB) and macro definition prefix files should also be compressed. For additional information regarding these two utility tasks, consult the applicable <u>RSX-11M</u> or <u>RSX-11D</u> Utilities Manual (see Section 0.3 in the Preface).

It often seems expedient to append a macro definition prefix file to each assembly to provide commonly-used macros. This practice, however, may produce the undesirable allocation of valuable dynamic

### ALLOCATING VIRTUAL MEMORY

storage for unnecessary macros. This side effect can be avoided by specifying that the prefix file containing the macros is a user-supplied macro library file (see Table 8-1). This action imposes the stipulation that the names of all desired macros must be listed as arguments in the .MCALL directive (see Section 7.8).

Storage for macro text can be re-used effectively by redefining certain types of macros to null after they have been invoked. This practice releases their dynamic memory for the storage of later macro text and also eliminates the overhead and the need for dynamic memory which would otherwise be required during the subsequent invocation and expansion of such non-redefined macros. The practice of redefining macros to null applies mainly to those that only define symbolic assignments, as shown in the example below. The redefinition process may be accomplished as follows:

SYM1 SYM2	=	.MACRO I VAL1 VAL2	DEFIN	;DEFINE	SYMBOLIC	ASSIGNMENTS.
		•				
OFF1 OFF2 OFF3	=	SYMBOL OFF1+SIZ1 OFF2+SIZ2		;DEFINE	SYMBOLIC	OFFSETS.
OFFN	, 1	Offm+SIZM				

.MACRO DEFIN .ENDM

;MACRO NULL REDEFINITION.

.ENDM DEFIN

Macros exhibiting this redefinition property should be defined (or read via the .MCALL directive) and invoked before all other macro definition and/or .MCALL processing. So doing ensures more efficient use of dynamic memory.

The following system macros have the automatic null redefinition property after once being invoked:

DRERR\$ - Directive return status codes

IOERR\$ - I/O return status codes

FILIO\$ - File-related I/O function codes

SPCIO\$ - Special I/O function codes

CSI\$ - Command String Interpreter codes and offsets

GCMLD\$ - Get Command Line codes and offsets

BDOFF\$ - FCS buffer descriptor offsets

FCSBT\$ - FCS bit value codes

FDOFF\$ - FCS file descriptor block offsets

FSROF\$ - FCS file storage region (FSR) offsets

NBOFF\$ - FCS filename block offsets

F-3

## F.3 OPERATIONAL TECHNIQUES

When, despite adhering to the guidelines discussed above, performance still falls below expectations, several additional measures may be taken to improve performance.

The first measure involves shifting the burden of symbol definition from MACRO-11 to the linker. In most cases, the definition of system I/O and FCS symbols (and user-defined symbols of the same nature) is not necessary during the assembly process, since such symbols are defaulted to global references (see Section 3.9 and Section D.1, category 4 of error code A). The linker attempts to resolve all global references from user-specified default libraries and/or the system object library (SYSLIB). Furthermore, by applying the selective search option for object modules consisting only of global symbol definitions, the actual additional burden to the linker is minimal.

A second way of making more dynamic memory available is to produce only one output file (either object or listing), as opposed to two. The additional file descriptor block (FDB) and file storage region (FSR) required to support the second output file are allocated from available dynamic memory at the start of each assembly. Furthermore, the size of the file storage region allocated is the minimum required for the second (listing) output file. For disk files, this is 264(10) words, and for direct line printer output, it is 74(10) words.

The final way of increasing available dynamic memory is related only to the operating environment. Under RSX-11M, MACRO-11 allocates all storage between its highest address and the end of its partition as dynamic memory. Consequently, the amount of working storage can be increased by installing and running MACRO-11 in a larger partition.

In IAS and RSX-llD, the assembler's dynamic memory is fixed at link time. If a larger assembler is not available, you may build one by increasing the size of the task's stack. This is accomplished by altering the STACK= option in the command file to build MACRO-ll.

### APPENDIX G

### WRITING POSITION INDEPENDENT CODE

# G.1 INTRODUCTION TO POSITION INDEPENDENT CODE

The output of a MACRO-11 assembly is a relocatable object module. The Task Builder binds one or more modules together to create an executable task image. Once built, a task can generally be loaded and executed only at the virtual address specified at link time. This is because the linker has had to modify some instructions to reflect the memory locations in which the program is to run. Such a body of code is considered position-dependent (i.e., dependent on the virtual addresses to which it was bound).

All PDP-11 processors offer addressing modes that make it possible to write instructions that are not dependent on the virtual addresses to which they are bound. A body of such code is termed position-independent and can be loaded and executed at any virtual address. Position-independent code can improve system efficiency, both in use of virtual address space and in conservation of physical memory.

In multiprogramming systems like IAS, RSX-11D and RSX-11M, it is important that many tasks be able to share a single physical copy of common code; for example a library routine. To make the optimum use of a task's virtual address space, shared code should be position-independent. Code that is not position-independent can also be shared, but it must appear in the same virtual locations in every task using it. This restricts the placement of such code by the Task Builder and can result in the loss of virtual addressing space.

The construction of position-independent code is closely linked to the proper usage of PDP-11 addressing modes. The remainder of this Appendix assumes you are familiar with the addressing modes described in Chapter 5.

All addressing modes involving only register references are position-independent. These modes are as follows:

R	register mode
(R)	deferred register mode
(R)+	autoincrement mode
@(R)+	deferred autoincrement mode
- (R)	autodecrement mode
@-(R)	deferred autodecrement mode

When using these addressing modes, you are guaranteed position-independence, providing the contents of the registers have been supplied such that they are not dependent upon a particular virtual memory location.

G-1

The relative addressing modes are position-independent when a relocatable address is referenced from a relocatable instruction. These modes are as follows:

A	relative	mode	
@A	relative	deferred	mode

Relative modes are not position-independent when an absolute address (that is a non-relocatable address) is referenced from a relocatable instruction. In this case, absolute addressing (i.e., @#A) may be employed to make the reference position-independent.

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Index modes can be either position-independent or position-dependent, according to their use in the program. These modes are as follows:

X(R) index mode @X(R) index deferred mode

If the base, X, is an absolute value (e.g., a control block offset), the reference is position-independent. For example:

N = 4	MOV	2(SP),R0	; POSITION-INDEPENDENT
14	MOV	N(SP),RO	; POSITION-INDEPENDENT

If, however, X is a relocatable address, the reference is position-dependent. For example:

CLR ADDR(R1) ; POSITION-DEPENDENT

Immediate mode can be either position-independent or not, according to its usage. Immediate mode references are formatted as follows:

#N immediate mode

When an absolute expression defines the value of N, the code is position-independent. When a relocatable expression defines N, the code is position-dependent. That is, immediate mode references are position-independent only when N is an absolute value.

Absolute mode addressing is position-independent only in those cases where an absolute virtual location is being referenced. Absolute mode addressing references are formatted as follows:

@#A absolute mode

An example of a position-independent absolute reference is a reference to the directive status word (\$DSW) from a relocatable instruction. For example:

MOV @#\$DSW,R0 ;RETRIEVE DIRECTIVE STATUS

### G.2 EXAMPLES

The RSX-11 library routine, PWRUP, is a FORTRAN callable subroutine to establish or remove a user power failure AST entry point address. Imbedded within the routine is the actual AST entry point which saves all registers, effects a call to the user-specified entry point, restores all registers on return, and executes an AST exit directive. The following examples are excerpts from this routine. The first example has been modified to illustrate position-dependent references, (see Figure G-1). The second example, Figure G-2, is the position-independent version.

# WRITING POSITION INDEPENDENT CODE

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PWRUP::			
	CLR CALL	-(SP) .X.PAA	;ASSUME SUCCESS ;PUSH (SAVE) ARGUMENT ADDRESSES ONTO .STACK
a Brian Ha Na Angela Sa Angela Tang A	.WORD MOV MOV BNE CLR BR	1.,\$DSW \$OTSV,R4 (SP)+,R2 10\$ -(SP) 20\$	CLEAR DSW, AND SET R1=R2=SP ;GET OTS IMPURE AREA POINTER ;GET AST ENTRY POINT ADDRESS ;IF NONE SPECIFIED, SPECIFY NO POWER ;RECOVERY AST SERVICE ;
10\$:			CEM ACM ENTRY DOINT
205	MOV MOV	#BA,-(SP)	; PUSH AST SERVICE ADDRESS
205:	CALL .BYTE	.X.EXT 109.,2.	;ISSUE DIRECTIVE, EXIT.
	•		
BA:	MOV MOV MOV	R0,-(SP) R1,-(SP) R2,-(SP)	;PUSH (SAVE) R0 ;PUSH (SAVE) R1 ;PUSH (SAVE) R2
		Figure G-1 Pos	sition-Dependent Code
PWRUP::	CLR	- (SP)	ASSUME SUCCESS
lagt við Max Gaby	CALL .WORD MOV MOV BNE CLR BR	.X.PAA 1.,\$DSW @#\$OTSV,R4 (SP)+,R2 10\$ -(SP) 20\$	;PUSH ARGUMENT ADDRESSES ONTO STACK ;CLEAR DSW, AND SET R1=R2=SP. ;GET OTS IMPURE AREA POINTER ;GET AST ENTRY POINT ADDRESS ;IF NONE SPECIFIED, SPECIFY NO POWER ;RECOVERY AST SERVICE
105:	MOV MOV	R2,F.PF(R4) PC,-(SP)	SET AST ENTRY POINT PUSH CURRENT LOCATION
20\$:	ADD CALL	#BA,(SP) .X.EXT	;COMPUTE ACTUAL LOCATION OF AST ;ISSUE DIRECTIVE, EXIT.
; : ACTUA	L AST SE	RVICE ROUTINE:	
; ACTON ; ; ;	1) SAVE 2) EFFE 3) REST 4) ISSU	REGISTERS CT A CALL TO SPE ORE REGISTERS E AST EXIT DIREC	CIFIED SUBROUTINE TIVE
; BA:	MOV MOV MOV	R0,-(SP) R1,-(SP) R2,-(SP)	;PUSH (SAVE) R0 ;PUSH (SAVE) R1 ;PUSH (SAVE) R2
		Figure G-2 Pos	ition-Independent Code
		이 물건 것 같은 것 같은 것 같이 있는 것 같이 많이 있다.	그 그는 것 같아요. 그렇게 한 것 같이 그는 것 것 같은 것 같아요. 이 것 것 같은 것 같아요.

The position-dependent version of the subroutine contains a relative reference to an absolute symbol (\$OTSV) and a literal reference to a relocatable symbol (BA). Both references are bound by the Task Builder to fixed memory locations. Therefore, the routine will not execute properly as part of a resident library if its location in virtual memory is not the same as the location specified at link time.

### WRITING POSITION INDEPENDENT CODE

In the position-independent version, the reference to \$OTSV has been changed to an absolute reference. In addition, the necessary code has been added to compute the virtual location of BA based upon the value of the program counter. In this case, the value is obtained by adding the value of the program counter to the fixed displacement between the current location and the specified symbol. Thus, execution of the modified routine is not affected by its location in the image's virtual address space.

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The MACRO-11 Assembler provides a way of checking the position-independence of code. In an assembly listing, MACRO-11 inserts a ' character following the contents of any word which requires the linker to perform a relocation operation. In some cases this character indicates a position-dependent instruction; in other cases, it merely draws the user's attention to the use of a symbol which may or may not be position-independent. The cases which cause a ' character to be inserted in the assembly listing are as follows:

 Absolute mode references are flagged with a ' character when the reference is relocatable. References are not flagged when they are position-independent (i.e., absolute). For example:

MOV @#ADDR,R1 ;PIC ONLY IF ADDR IS ABSOLUTE.

 Index and index deferred mode references are flagged with a ' character when the offset is relocatable. For example:

MOV	ADDR(R1),R5	;NON-PIC	IF	ADDR	IS	RELOCATABLE.
MOV	@ADDR(R1),R5	;NON-PIC	IF	ADDR	IS	RELOCATABLE.

3. Relative and relative deferred mode references are flagged with a ' character when the address specified is relocatable with respect to another program section. For example:

MOV	ADDR1,R1	;NON-PIC WHEN ADDR1 IS BOUND
MOV	@ADDR1,R1	;TO ANOTHER PROGRAM SECTION

4. Immediate mode references to relocatable addresses are always flagged with a ' character.

MOV	#3,R0	;ALWAYS POSITION-INDEPENDENT.
MOV	#ADDR,R1	;NON-PIC WHEN ADDR IS RELOCATABLE.

There is one case in which the MACRO-11 assembler does not flag a potential position-dependent reference. This occurs where a relative reference is made to an absolute virtual location from a relocatable instruction (i.e., MOV \$OTSV,R4 in Figure F-1).

Those references requiring more than simple relocation at link time are also indicated in the assembly listing. Simple global references are flagged with the letter G. Those which contain multiple global references or complex relocation, are flagged with the letter C (see Section 3.9 and Chapter 4). In such cases, it is difficult to positively state which are or are not position-independent. However, in general, it is safe to apply the guidelines discussed earlier in this Appendix to the resulting address value produced at link time.

### APPENDIX H

SAMPLE ASSEMBLY AND CROSS REFERENCE LISTING

T	I	T	L	E	P3PRE	
1	D	E	N	T	/05/	

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1 COPYRIGHT 1976, DIGITAL EQUIPMENT CORP., MAYNARD, MASS. I THIS SOFTWARE IS FURNISHED TO PURCHASER UNDER A LICENSE FOR USE ON A SINGLE COMPUTER SYSTEM AND CAN BE COPIED (WITH INCLUSION OF DEC'S COPYRIGHT NOTICE) ONLY FOR USE IN SUCH SYSTEM, EXCEPT AS MAY OTHERWISE BE PROVIDED IN WRITING BY DEC. THE INFORMATION IN THIS DOCUMENT IS SUBJECT TO CHANGE WITHOUT NOTICE AND SHOULD NOT BE CONSTRUED AS A COMMITMENT BY DIGITAL Equipment corporation. DEC ASSUMES NO RESPONSIBILITY FOR THE USE OR RELIABILITY OF ITS SOFTWARE ON EQUIPMENT WHICH IS NOT SUPPLIED BY DEC. VERSION 05 1 C. MONIA 05-FEB-76 PERFORM PRELIMINARY SYMBOL TABLE PROCESSING I LOCAL MACROS DEFINE CANNED SECTION TABLE ENTRIES SECTN NAME, FLAGS, SEG I WHEREI NAME=SECTION NAME OFFSET=OFFSET IN SEGMENT DESCRIPTOR TO RECEIVE SECTION ADDRESS. BLANK IF NONE. FLAGS=SECTION FLAGS BYTE CONTENTS SEG='ROOT' IF SECTION IS DEFINED IN ROOT SEGMENT ONLY 1 . 1 MACRO SECTN NAME, FLAGS, OFFSET, SEG .=.-2 \$\$0=. .WORD 0 , WORD 0 \$55=. . . . 4 .RADSØ /NAME/ .=\$\$\$ .BYTE FLAGS IF 8 SEG BYTE Ø IFF BYTE =1 ENDC IF NB OFFSET WORD OFFSET IFF

. WORD & ENOC Sctlgh <b>e. = 55</b> 0 . Word Endm . Endm	FLOCAL DATA	I COMMON STORAGE SHAPED WITH MODULE "PRCLM"	THE CONTENTS OF THE SECTION DEFINED BELOW MUST MATCH The section contents specified in module "prclm".	PSECT PRCLSD, D, OVR	CRSCTI BLKW 10° I CUPRENT SECTION DATA CrvSCI Blkw 1 I VIRTUAL ADDRESS OF CURRENT LOADNMI RADSO /SLOAD/ I NAME OF MANUAL LOAD ROUTIN	MODN: BLKW 1 5 FIRST HALF OF MODULE NAME LCLSYM: BLKW 3. 1 LOCAL SYMBOL STORAGE	SKCPF: BLKW 1 5 EARCH PATH FLAG SRVSG: BLKW 1 1 1 VIRTUAL ADDRESS OF SEARCH SVDEF: BLKW 1 5 YMBOL DEFINITION FLAG	SYENT BLAW 1 SYMBOL ENTRY ADDRESS Symbol Blaw 1 J CURRENT SYMBOL FLAGS	. PSECT	1 6SD DISPATCH VECTOR TABLE	THIS TABLE CONTAINS THE ADDRESS OF SUBROUTINES WITHIN MODU PRCLM THAT PROCESS SPECIFIED GSD RECORD TYPES, PRIOR TO ISSUING A SUBROUTINE CALL TO ANY OF THESE ENTRY POINTS, TH FOLLOWING MUST BE SETUP.	INPUTS.	RGBREAL ADDRESS OF SEGMENT DESCRIPTOR RIEADDRESS OF LOCAL SYMBOL DESCRIPTOR (LCLSYM) R4=ADDRESS OF DUMMY ELEMENT DECRIPTOR CRSCT=CURRENT SEGMENT DESCRIPTOR CRSCT=CURRENT SECHENT DESCRIPTOR CRSCT=CURRENT SECHENT DETA (SETUP BY PSTNME OR SCTNMI MODN=ME OF SEGMENT BEING PROCESSED SCRELM=ADDRESS OF DUMMY ELEMENT DESCRIPTOR SCRELM=ADDRESS OF SEGMENT DESCRIPTOR SCRVEL=0 (VITTUAL ADDRESS OF SEGMENT DESCRIPTOR SCRVEL=0 (VITTUAL ADDRESS OF SEGMENT DESCRIPTOR
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293	000542	012700		NOM MO	#LCLSYM+2.RG	POINT TO LOCAL DESCRIPTOR
		000036	-			
294	000546	012320		>OM	(R3)+,(R0)+	I COPY SECTION NAME
295	000550	001434		BEQ	305	IF EG AT END OF LIST
296	000552	012320		>0¥	(R3)+.(R0)+	COPY SECOND HALF OF NAME
297	00554	012320		YOY	(R3)+.(R0)+	COPY FLAGS
860	000200	010100		CLR		
000	000200	105740		TSTB		· RONT SEGMENT ON Y?
100	000000	00100		BED		
200	192000	11140			EPUSC EPTER	
•		0000000				
		0000000				
302	000572	001022		BNE	205	· TF NF ND
6	999574					
105	899574	195010		CL RB	(Ra)	· FIEAR HTCH FLACS RVTF
202	000576	032767		BIT	TOXOL TN9304	BUILDING RESIDENT LIBRARY?
		00000000				
		177316				
306	000604	001402		BEG	158	IF EQ NO
105	900606	112710		MOVB	#CSSRES, (R0)	J FLAG SECTION AS TASK-RESIDENT
10.8	000412		. 5.5 .			-
	000010	010500		20		. FODV DEAL ANNDERS OF RECHENT
	010000	101010				· COLL ADDREA DE LOFAL DESENTAT
		000034				
311	000620			CALL		FUTER SECTION IN TARIES
312	000624	010500		NOV	R5. R0	I GET ADDRESS OF SEGMENT DESCRIPTOR
313	000626	061300		ADD	(R3) . R0	OFFSET TO PROPER SLOT IN DESCRIPTOR
314	000630	020005		Q M D	R0, R5	NO OFFSET?
5	000632	001402		BEQ		IF EQ NO OFFSET TO BE SETUP
316	000634	016710		NOM	CRVSC, (RØ)	I STORE SECTION ADDRESS
		000024				
317	000040		2051			
318	000640	005723		<b>TST</b>	(R3)+	I CLEAR C. ADVANCE TO NEXT ENTRY
319	000642		3051			
320	800642			RETURN		
321						
322			-			
323			INSER	T A GLOB	AL SYMBOL IN SEG	MENT.
326						
S			1 LUCAL	SYMBUL	DESCRIPTOR AND C	URRENT SECTION MUST BE SETUP
979						
	1170000	013701	I WASNI			
		000014		>	1 × 4 × 4 × 4 × 1 × 1	THE AUTRESS OF LUCAL DESCALFION
330	000650	010500		NOM	R5, R0	I COPY ADDRESS OF SEGMENT DESCRIPTOR
122	000652			CALLR	•SYMDCL	J ENTER SYMBOL IN TABLE
333		000001		END.		

# SAMPLE ASSEMBLY AND CROSS REFERENCE LISTING

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SYMBOL	ŝ	ISS REFERE	NCE				CREF VØ1	
SYMBOL	5	LUE		REFERENCES				
80		000015		#1-197				
CRSCT		800000 R		#2-77				
CRVSC		000054 R		#2-78	2-316			
CSSALO	•	*****	×	2=166	2-167	2-169		
CSSGBL	•	*****	×	2-166				
CSSREL		*****	×	2-164	2-165	2-166	2-167	201-2
repre	•		>	2-107 2-107	a/1=>			
COST COST	• •			2-144	2-165	2-167	2-168	2-169
	)			2-170	•		•	•
VMMUQ		900000 R		#2-135	2-263	2-284		
ESLCMT		******	×	+2-264				
ESLNUM		*****	X	+2-265				
ESLOWT		*****	×	982-2				
		1 1 0 0 0 0 0		101-13				
					100-04			
- JOLT				2-281	#2=72B			
TNTSVM		000000		#2-122				
LCLSVM		000034 8		#2=81	2-276	2-293	2-310	2+329
		000012		#1-197		1		
LOADNH		000026 R		#2-79				
MARKS		000022 R		#2-149	2-228	2-245		
NOOM		000032 R		#2-80	2-266			
MODNME		000000 B		#2-120				
DOT NME		800012 R		#2-125	2-311			
		1 1/1000		102-2	012-24			
	1	-			212-2			
					# 3= 1 A 5	#2=144	#2-167	#2-168
	•	a10000		691-04	-140 -140	2-272		
SCT NW		000010 8		12-162	2-268			
SCTNME		000002 R		#2-121				
SGSGBL		*****	×	2-274				
SPA		900049		*1-197				
SRCPF		000054 R		#2=82				
<b>JEVSG</b>		0000020 4		50-28				
			4	C 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2	CAC-2			
SVFLG		000007 B		#2-86				
SYMDEL		000010 R		#2-124	2-331			
SYSDEF		******	XS	2-279				
SYSGBL		******	X	2-230	2-279			
SYSREL	•	*****	×	2-279				
XX9X9		*****	X	2-237	2-250			
986AUT	•	*****	X	2-164				
JENDS0	•	*****	X	2-201	292-2	2-277	2-210	
		****	× >	Ce1e2				
			× > 9 C	100100	2-254			
			ر بر 2 د	2-244	0-3-3			
リーリラウウ			()	1-1-1-1				

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No.     No.     No.     No.       No.     No.     No.	JUDAT	KEPEKENC	D U			
N - 2     N - 2     N - 2     N - 2       N - 2     N - 2 <th>2</th> <th>+2-205</th> <th>+2-208</th> <th>*2-210</th> <th>*2-212</th> <th>+2-23</th>	2	+2-205	+2-208	*2-210	*2-212	+2-23
		*2*222	+2=240	+2-243	*2-257	*2=25
No.     No.     No.     No.     No.       No.     No.		+2-270	+2-281	*2*311	+2-320	
No.2     No.2     No.2     No.2     No.2       No.2     No.2 <t< td=""><td>8</td><td>2-206</td><td>+2-209</td><td>+2-226</td><td>*2-227</td><td>2-24</td></t<>	8	2-206	+2-209	+2-226	*2-227	2-24
2     2     2       2     2       2 <td></td> <td>+2=546</td> <td>+2=247</td> <td>+2-248</td> <td>2-249</td> <td>2-25</td>		+2=546	+2=247	+2-248	2-249	2-25
		+2-255	+2-256	+2-260	*2-261	2-26
2-200     2-200       2-200		+2-276	*2-277	*2=278	+2-279	2-28
2-304     2-304       2-316     2-316       2-316		+2-284	*2-293	+2-24	*2-296	+2-24
1     2-314     2-314     2-314     2-314       2-204     2-204     2-224     2-224       2-203     2-233     2-234     2-224       2-204     2-234     2-234     2-224       2-204     2-235     2-234     2-224       2-204     2-234     2-234     2-224       2-204     2-234     2-234     2-224       2-204     2-234     2-234     2-224       2-214     2-234     2-234     2-234       2-214     2-244     2-244     2-244       2-214     2-244     2-245     2-24       2-215     2-234     2-245     2-24       2-215     2-245     2-245     2-24       2-215     2-245     2-245     2-24       2-215     2-245     2-245     2-24       2-215     2-245     2-245     2-24       2-215     2-245     2-245     2-245       2-215     2-245     2-245     2-245       2-215     2-245     2-245     2-245       2-216     2-245     2-245     2-245       2-217     2-234     2-234     2-245       2-216     2-245     2-245     2-245       2-217 <td< td=""><td></td><td>2-298</td><td>*2-299</td><td>2-304</td><td>2-307</td><td>+2-30</td></td<>		2-298	*2-299	2-304	2-307	+2-30
1       *2*200		*2-312	*2-313	2-314	2-316	*2-33
2<:33	-	+2-584	+2-207	*2-221	2-224	+2-22
2     2 <td></td> <td>2.229</td> <td>2-233</td> <td>2-239</td> <td>*2-242</td> <td>+2-31</td>		2.229	2-233	2-239	*2-242	+2-31
2-239     2-239     2-239     2-249     2-249       2-313     2-240     2-294     2-294       2-313     2-294     2-294       2-313     2-294     2-294       2-219     2-294     2-294       2-219     2-294     2-294       2-219     2-294     2-294       2-219     2-294     2-294       2-219     2-294     2-294       2-219     2-294     2-294       2-219     2-294     2-294       2-219     2-294     2-294       2-312     2-294     2-294       2-312     2-294     2-294       2-312     2-294     2-294       2-312     2-294     2-294       2-312     2-294     2-294       2-312     2-294     2-294       2-312     2-294     2-294       2-312     2-314     2-313       2-312     2-314     2-314       2-314     2-314     2-314       2-315     2-314     2-314       2-315     2-314     2-314       2-315     2-314     2-314       2-315     2-314     2-314		+2-329				
		852-24	*2=237	2-239	+2-245	*2-27
			24244	405-54	100-00	00 - 0 -
		2-112	+2=118			
<ul> <li>3.2</li> <li>4.2</li> <li>5.338</li> <li>5.338</li> <li>5.424</li> <li>5.424</li> <li>5.424</li> <li>5.424</li> <li>5.424</li> <li>5.424</li> <li>5.425</li> <li>5</li></ul>	4		2-264	24245	2-266	
	-	+2-219	2-24	2=226	2.50	40-0
		2-262	2-274	2-277	2-278	2-30
		2.312	2-314	2=330		
			•			
構成長期 一番 医素を含くていた。 「「「」」」」 「「」」」」 「「」」」」 「」」」」 「」」」」 「」」」 「」」」 「」」」 「」」」 「」」」 「」」」 「」」」 「」」」 「」」」」 「」」」」 「」」」」 「」」」」 「」」」 「」」」」 「」」 「」」」 「」」 「」」」 「」」 「」」 「」」 「」」 「」」 「」」」 「」 「						

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# SAMPLE ASSEMBLY AND CROSS REFERENCE LISTING

44CRO CROSS REFERENCE 44CRO NAME REFERENCES C4LL #1-446 2=205 2=208 2=210 2=220 C4LL #1-446 2=205 2=243 2=220 C4LL #1-75 #2=205 2=243 2=257 2=220 C4LLR #1=75 #2=205 #2=331 C4LLR #1=197 #1=197 #1=197 C4LLR #1=105 #2=208 #2=150 C4LLR #1=145 #2=208 =2=210 C4LLR #1=197 #2=208 =2=210 C4LLR #2=100 #2=107 #2=150 #2=150 C4LLR #2=160 #2=150 #2=150 #2=150 C4LLR #2=160 #2=150 #2=150 #2=150 C4LLR #2=160 #2=150 #2=150 #2=150 #2=150 C4LLR #1=175 #2=160 #2=170 #2=150 #2=150 #2=150 C4LLR #1=175 #2=160 #2=170 #2=150							
44CRO NAME REFERENCES CALL #1-446 2-205 2-208 2-219 2-220 CALL #1-446 2-243 2-27 2-220 CALLR #1-75 #2-201 2-219 2-250 CALLR #1-75 #2-201 2-311 CALLR #1-197 #1-197 2-250 #1-197 #1-197 CALLR #1-145 #2-212 #2-320 MAVRG #1-145 #2-212 #2-320 MAVRG #1-156 2-219 #2-320 MAVRG #1-175 2-210 #2-320 MAVRG #1-175 2-210 #2-320 MAVRG #1-195 2-210 #2-320 MAVRG #1-195 2-210 #2-100 #2-100 #2-1007 MAVRG #1-175 2-230 MAVRG #1-175 2-230 MAVRG #1-195 2-230 MAVRG #1-195 2-230 MAVRG #1-195 2-230 MAVRG #1-195 2-230 MAVRG #2-104 #2-106 #2-106 #2-106 #2-107 MAVRG #2-106 #2-106 #2-106 #2-106 #2-107 MAVRG #1-175 2-320 MAVRG #2-106 #2-1	ACRO CROSS	REFERENCE				CREF VØ	-
CALL #1-46 2-205 2-208 2-210 2-220 2-222 2-249 2-210 2-250 2-21LR #1-75 #2-281 2-311 2-257 2-250 2-21LR #1-75 #2-285 #2-31 25F 2-281 2-31 26F 200 2-281 2-31 26F 200 2-281 2-29 27805 #1-197 2-219 20L 201 201 2-219 20L 201 201 2-219 20L 201 201 2-219 20L 201 2-200 2-219 20L 201 201 2-219 20L 201 2-200 2-219 201 2-200 2-219 201 2-200 2-200 2-219 201 2-200 2-2	ACRO NAME	REFERENCES					
2=222     2=240     2=243     2=257     2=250       2=10     2=281     2=311     2=251     2=250       201     #1=75     #2=285     #2=331     2=257     2=250       201     #1=197     #1=197     #1=197     #1=197       201     #1=145     #2=212     #2=320     #2=250       201     #1=145     #2=212     #2=320     #2=160       201     #1=156     #2=219     #2=320     #2=160       201     #1=156     #2=116     #2=160     #2=160       201     #1=156     #2=160     #2=160     #2=160       201     #1=175     #2=160     #2=170     #2=160       201     #1=175     #2=160     #2=160     #2=160	SALL	#1-46	2-205	2-208	2-210	2-22	6
ALLR     2-270     2-201     2-311       SEF     #1-75     #2-205     #2-311       SERORS     #1-107     #1-107     #1-107       SRAORS     #1-143     #2-212     #2-320       SETURN     #1-145     #2-212     #2-320       SETURN     #1-145     #2-212     #2-320       AVRG     #1-146     #2-210     #2-104       AVVR     #1-166     #2-104     #2-106       ALVR     #1-166     #2-104     #2-106       ALVR     #1-175     #2-109     #2-106       ALVR     #1-175     #2-108     #2-106       ALVR     #1-175     #2-108     #2-106	•	2-22	2-240	2-243	2-257	2-25	80
.ALLR     #1-75     #2-285     #2-331       DEF     #1-197     #1-197       DEF     #1-197     #1-197       DEF     #1-123     #1-197       DECT     #1-145     #2-212       AVVR     #1-166     #2-164       AVVR     #1-166     #2-164       AVVR     #1-166     #2-169       AVVR     #1-175     2-219       AVVR     #1-166     #2-169       AVVR     #1-175     #2-169       AVVR     #1-175     #2-169       AVVR     #1-175     #2-169       AVVR     #1-175     #2-169		2-270	2-281	2-311			
SEF $\#1=03$ $\#1=197$ SRORs $\#1=123$ $\#1=197$ SUL $\#1=123$ $\#1=123$ ALVR $\#1=145$ $\#2=212$ $\#2=320$ SLVRG $\#1=156$ $2=219$ $\#2=320$ SLVRG $\#1=156$ $2=219$ $\#2=320$ SLVRG $\#1=156$ $2=219$ $\#2=160$ SLVRG $\#1=156$ $#2=160$ $\#2=160$ SECTN $#2=160$ $#2=160$ $#2=160$ SOB $#1=90$ $#2=160$ $#2=170$	SALLR	#1-75	#2-285	#2-331			
ZMRORS #1-123 DCL #1-145 TCL #1-145 RETURN #1-167 #2-212 #2-320 AAVRG #1-156 2-219 AAVRG #1-156 #2-164 #2-166 #2-167 RECTN #2-168 #2-169 #2-170 #2-166 #2-167 DSECT #1-98	55	#1-83	#1-197	#1-197			
JOL #1-145 TETURN #1-145 TETURN #1-156 2-212 #2-320 34VVR #1-156 2-219 34VVR #1-156 2-219 34VVR #2-166 #2-166 #2-166 #2-169 #2-169 #2-166 #2-169 #2-169 #2-167 305 #1-175 305 #1-175	ERORS	#1-123					
RETURN #1-167 #2-212 #2-329 SAVRG #1-156 2-219 #2-329 SAVVR #1-166 2-219 #2-165 #2-166 #2-167 BECTN #2-168 #2-169 #2-178 #2-166 #2-167 \$08 #1-175 #2-169 #2-178	JOL	#1-145					
SAVRG #1-156 2-219 SAVVR #1-166 2-219 SAVVR #1-166 #2-165 #2-166 #2-167 BECTN #2-168 #2-169 #2-170 #2-166 #2-167 \$08 #1-175 #2-169 #2-170 \$08 #1-175	RETURN	#1-187	#2-212	#2-320			
34VR #1=166 BECTN #2-40 #2-164 #2-165 #2-166 #2-167 #2=169 #2-169 #2-170 #2-170 #1=175 \$0B #1-98	BAVRG	#1-156	2-219				
SECTN #2-40 #2-164 #2-165 #2-166 #2-167 #2-168 #2-169 #2-170 #2-170 \$05 #1-175 \$08ECT #1-98	SAVR	#1-166					
#2=168 #2=169 #2=170 808 #1=175 #2=169 #2=170 05ECT #1=98	JECTN	#2-40	#2-164	#2-165	#2-166	#2-16	~
808 #1=175 05ECT #1=98		#2-168	#2-169	#2-170			
.09ECT #1-98	308	#1-175					
	, DSECT	86-1#					

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2-241
2-203 2-215 2-235 2-255
2-232 2-25
2-282 2-207 2-26 2-304
2-202 2-22 2-252 2-25 2-265 2-33
2-205 2-206 2-240 2-243 2-311
2-204 2-206
2-247 2-248
2-262 2-263
2-296 2-297
2-329 2-338 2-238 2-239
2-212 2-320 2-292
2-246 2-318 2-299
2-151 2-77 2-7A
2-84 2-85
2-123 2-124
2-152 2-164
2-166 2-167 2-169 2-178
2-333
2=169 2=169 1=3 2=2
1-174 2-164
2=166 2=167
2-220 2-22
2-222 2-22
2-240 2-24

SAMPLE ASSEMBLY AND CROSS REFERENCE LISTING

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H-8

INDEX

Absolute addresses, 6-14 Absolute binary output, 6-14 Absolute expression, 3-16 Absolute mode, 5-5, 5-7 Absolute module, 6-34 Absolute program section, 6-37 Address boundaries, 6-31 Address mode syntax, B-1 Address modes, 5-1 Addressing forms, summary, 5-7 Allocating byte data, 6-17 Allocating dynamic memory, F-1 Allocating word data, 6-18 Allocation requirements, 6-35 Alternate radix, 6-25 Ampersand, 3-1 Angle brackets, 3-3, 3-15, 6-4, 6-25, 6-28, 7-4 to 7-5, 7-16 to 7-17 Apostrophe, 7-10 Argument substitution, 7-16 Arithmetic addition operator or autoincrement indicator 3 - 1Arithmetic division operator 3 - 1Arithmetic multiplication operator, 3-1 Arithmetic subtraction operator or autodecrement indicator, 3-1 ASCII character set, A-1 ASCII conversion, 3-14 ASCII conversion characters 6-19 .ASCII directive, 6-20 .ASCIZ directive, 6-21 .ASECT directive, 6-38 Assembler directives, 6-1, B-1, B-2 Assembler version, 6-8 Assembly language, B-1 Assembly listing, 2-6 Assembly pass 1, 1-1 Asterisk, 3-1 At sign, 3-1 Attribute of the current location counter, 3-12 Autodecrement deferred mode, 5-3, 5-7 Autodecrement mode, 5-7 Autoincrement deferred mode, 5-3, 5-7 Autoincrement mode, 5-2, 5-7

B operator, 6-25
Backslash, 3-1
Binary operators, 3-15
Blank lines, 2-2
Blocks of storage, reserving,
 6-30
.BLKB directive, 6-30
BLKW directive, 6-30
Branch instruction addressing,
 5-8
.BYTE directive, 6-17

<sup>°</sup>C operator, 6-27 Calling conventions, E-8 Calling macros, 7-3 Changing default radix, 3-13 Changing value of location counter, 3-12 Character set, 3-1 Character substitution, 7-16 Code and data separation, 6-38 Code or data sharing, 6-38 Coding standard, E-1 Colon, 3-1 Comma, 3-1 Comment, 6-14, E-2 Comment field, 2-5 Comment field indicator, 3-1 Complementing an argument, 6-27 Complex relocatable expression, 3-16 Complex relocation, 4-1 Concatenated, 6-35 Concatenation of macro arguments, 7-10 Conditional assembly block, 6-41 Conditional assembly directive, 6-41, 6-42 Conditional branches, E-13 Continuation lines, 2-2 Creating local symbols automatically, 7-7 Creating program sections, 6-36 Cross-reference listing (CREF), H-5 .CSECT directive, 6-17 Current location counter, 2-2, 3-11, 3-14, 5-6, 6-29

<sup>^</sup>D operator, 6-25 Data storage directives, 6-17 Date, 6-8 Default object module name, 6-11 Default register definitions, 6-15 Deferred addressing indicator, 3-1 Defining macros, 7-1 Device registers, E-2 Diagnostic, 7-14 Diagnostic error message summary, D-1 Direct assignment operator, 3-1 Direct assignment statements, 3-7 Directives, 2-5, 5-9, 6-1 Double ASCII character indicator, 3-1 Double colon, 3-1, 3-7 Double equal sign, 3-1, 3-7 Double quote, 3-1, 3-14, 6-19 .DSABL directive, 3-7, 3-9 6-13 to 6-15, 6-27 Duplication of code, 7-17

EMT, 5-8 .ENABL directive, 5-7, 6-13, 6-15, 6-27 .END directive, 6-31 .ENDC directive, 6-41 .ENDM directive, 7-2 End of the source input, 6-31 .ENDR directive, 7-17 Entry-point instructions, 6-33 .EOT directive, 6-31 Equal sign, 3-1 Error codes, D-1 .ERROR directive, 7-14 Evaluation of expressions, 3-15 .EVEN, 6-29 Exclamation point, 3-1 Executable task image, 4-1 Exiting, E-9 Expressions, 3-14, 3-15 External expression, 3-15, 3-16 External symbols, 6-40 Externally-defined macro, 7-18

^F operator, 3-14, 6-27 Finding address mode of macro arguments, 7-13 Finding number of characters in strings, 7-12 Floating-point data, 6-26 number, 6-28 number specification, 6-27 rounding, 6-15, 6-27 storage directives, 6-27 truncation, 6-14, 6-27.FLT2 directive, 6-27 .FLT4 directive 6-27 Forbidden instruction usage, E-12 Form-feed, 6-13, 7-3 Format control, 2-6 Formatting standards, E-9 Forward referencing, 3-8 Function control switches, 8-6 Function directives, 6-13

Å.

ł

ų

General purpose registers, E-2
General registers, 3-9
Global
 label, 6-40
 references, 6-15
 symbol, 2-3, 6-40
 symbol directory, 1-2
.GLOBL directive, 3-7, 6-39
GSD, 1-2

Hardware registers, E-2 Horizontal formatting, 2-6

.IF directive, 6-41 .IFF directive, 6-43, 6-44 .IFT directive, 6-43 .IFTF directive, 6-43 .IIF directive, 6-46 Illegal characters, 3-3 Immediate conditional assembly, 6-46 Immediate expression indicator, 3-1 Immediate mode, 5-4, 5-7 Immediate mode deferred, 5-5 Implicit .WORD directive, 2-5, 6-18 Indefinite repeat block directives, 7-15 Index deferred mode, 5-4, 5-7 Index mode, 5-4, 5-7 Initial argument or expression indicator, 3-1 Initial register indicator, 3-1

Instruction set, C-1 .IRP directive, 7-15 .IRPC directive, 7-15, 7-16 Item or field terminator, 3-1

Keyword arguments, 7-4, 7-9

Label field, 2-2 Label terminator, 3-1 Left angle bracket, 3-1 Left parenthesis, 3-1 .LIMIT directive, 6-31 Line format, E-1 Linker, 2-2, 4-1 Linking, 4-1, 4-39 .LIST directive, 6-1 Listing conditional assemblies, 6-4 Listing control directives, 6-1 Listing level count, 6-2 Listing of binary extensions, 6-3 Listing of comments, 6-4 Listing of generated binary code, 6-3 Listing of macro class, 6-4 Listing of macro definitions, 6-4 Listing of macro expansion binary code, 6-4 Listing of repeat range expansions, 6-4 Listing of source lines, 6-3 Listing of the current location counter, 6-3 Listing of the symbol table, 6 - 4Local symbol block, 6-14 Local symbol block delimiters, 3-10 Local symbols, 3-6, 3-10, 3-11 Location counter, 6-36 Location counter control directives, 6-29 Logical AND operator, 3-1, 6-42 Logical inclusive OR operator, 3-1, 6-42 Lower-case ASCII, 6-14

Macro arguments, 7-6 Macro attritube directives, 7-11 Macro call, 2-5, 7-3, 7-5 Macro call arguments, 7-4 Macro call numeric argument indicator, 3-1 Macro definition, 7-1, 7-15 Macro definition arguments, 7-4 Macro definition formatting, 7-3 Macro definition termination, 7 - 2Macro definitions and expansions, F-2 Macro directives, 7-1 Macro expansion termination, 7-3 Macro library directive, 7-18 Macro name, 7-1, 7-4 Macro names, E-5 Macro nesting, 7-5 Macro symbol table, 3-5 MACRO-11 character sets, A-1 MACRO-11 directives, 5-9, C-4 MACRO-11 symbols, 3-5 .MCALL directive, 7-18 Memory allocation, 6-32, 6-33, 6-38 Memory allocation and mapping, 6-32 .MEXIT directive, 7-18 Minus sign, 3-1 Modularity, E-8 Module checking routines, E-9 Module preface, E-5 Multi-defined label, 2-4 Multiple definitions of local symbols, 3-11 Multiple labels, 2-4

Naming standards, E-2 .NARG directive, 7-11 .NCHR directive, 7-11, 7-12 Negative numbers, 3-14 Nested conditional directives, 6-43 Nested macros, 7-3, 7-5 .NLIST directive, 6-1, 6-11 .NTYPE directive, 7-11, 7-13 Number of arguments in macro calls, 7-7, 7-11 Number sign, 3-1, 5-4 Numbers, 3-13 Numeric control, 6-24 Numeric control operators, 6-26, 6-27 Numeric directives, 6-26

^O operator, 6-25 Object module, 4-1 Object module name, 6-11

Octal radix, 3-13 .ODD directive, 6-29 Op codes, 2-4, C-1 Operand field, 2-4 Operator field, 2-4 Order of symbol table search, 3-6 Other symbols, E-3 Overlaid, 6-35 Overlays, 6-33

.PAGE directive, 6-13 Page eject, 7-3 Page ejection, 6-13 Page formatting, 2-6 Page headings, 6-8 Page number, 6-8 PAL-11R conditional assembly, 6-46 Passing numeric arguments as symbols, 6-46 Percent sign, 3-1 Permanent symbol table, 3-5, C-1 Plus sign, 3-1 Position independent code, G-1 .PRINT directive, 7-14 Processor priority, E-3 Program boundaries directive, 6-31 Program counter, 3-9, 5-1 Program modules, E-5 Program section access, 6-33 Program section name, 6-33 Program sections, 3-12, 6-32 Program source files, E-12 Program-local symbols, E-4 Programming standards and conventions, 2-1 .PSECT directive, 3-12, 6-32, 6-35

R operator, 6-23 .RAD50 directive, 3-13, 6-24 Radix control, 6-24 Radix control operators, 6-25 Radix-50 character set, A-4 Radix-50 control operator, 6-23 Radix-50 data, 6-22 .RADIX directive, 3-13, 6-24 Read-only access, 6-33 Read/write access, 6-33 Register deferred mode, 5-2 Register expression, 5-1

Register, mode, 5-1, 5-7 Register standards, E-2 Register symbols, 3-9 Register term indicator, 3-1 Relative addresses, 6-14 Relative addressing mode, 5-6 Relative deferred mode, 5-6, 5-7 Relative mode, 5-5, 5-7 Relocatability, 6-34 Relocatable expressions, 3-16, 4 - 1Relocatable module, 6-34 Relocatable program sections, 6-37 Relocation, 4-1 Relocation bias, 2-2, 6-34 Repeat block directive, 7-17 .REPT directive, 7-17 Reserving storage, 6-30 Reserving storage space, 3-13, 6-29 Right parenthesis, 3-1

.SBTTL directive, 6-8, 6-11 Scope of the program section, 6-33 Semicolon, 3-1 Sending messages to listing file, 7-14 Separating and delimiting characters, 3-2 Single ASCII character indicator, 3-1 Single quote, 3-1, 3-14, 6-19, 7-10 Slash, 3-1 Source line sequence numbers, 6-3 Space, 3-1 Special characters, B-1 Special characters in macro arguments, 7-6 Stack pointer, 3-9 Statement format, 2-1 Storing Radix-50 data, 6-23 Subconditional assembly, 6-43 subtitle, 6-8 Success/failure indication, E-9 Symbol control directive, 6-39 Symbol examples, E-4 Symbol table listing, 1-2 Symbolic arguments of listing control directives, 6-3, 6-4 Symbols, E-3 Symbols and expressions, 3-1 System macro libraries, 7-18

Tab, 3-1 Tab character, 2-2 Table of contents, 6-4, 6-11 Teleprinter mode, 6-4 Terminal argument or expression indicator, 3-1 Terminal register indicator, 3-1 Terminating directives, 6-31 Terms, 3-14 Time-of-day, 6-8 .TITLE directive, 6-11 Title of the object module, 6-8 Translating to ASCII, 6-20, 6-21 Translating to Radix-50, 6-22 Trap instructions, 5-8

1

1

Unary operator ordering, 6-28 Unary operators, 3-15 Unconditional assembly, 6-43 Undefined symbols, 3-7, 3-14 Universal unary operator or argument indicator, 3-1 Up arrow or circumflex, 3-1 Up-arrow, 3-3 Up-arrow (°) construction, 7-5 User symbol table, 3-5 User-defined and macro symbols, 3-5 User-defined macro libraries, 7-18 Using the standard symbolics, E-3

Version number, 6-12 Version number standard, E-13

Unary and binary operators, 3-4 Unary control, 6-24

.WORD directive, 3-11, 6-18

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